

Welcome to CS103!

Documents:

- Course Syllabus

Today:

- Course Overview
- Introduction to Set Theory
- The Limits of Computation

Zoom Etiquette

Are there “laws of physics”
in computer science?

Introduction to Set Theory

Key Questions in CS103

What problems can you solve with a computer?

- ***Computability Theory***

Why are some problems harder to solve than others?

- ***Complexity Theory***

How can we be certain in our answers to these questions?

- ***Discrete Mathematics***

Instructor

Ryan Smith ([*rsmith20@cs.Stanford.edu*](mailto:rsmith20@cs.Stanford.edu))

TAs

Sofía Dudas

Woody Wang

Staff Email List: cs103-sum1920-staff@lists.stanford.edu

Course Website

<https://cs103.stanford.edu>

Prerequisite / Corequisite

CS 106B

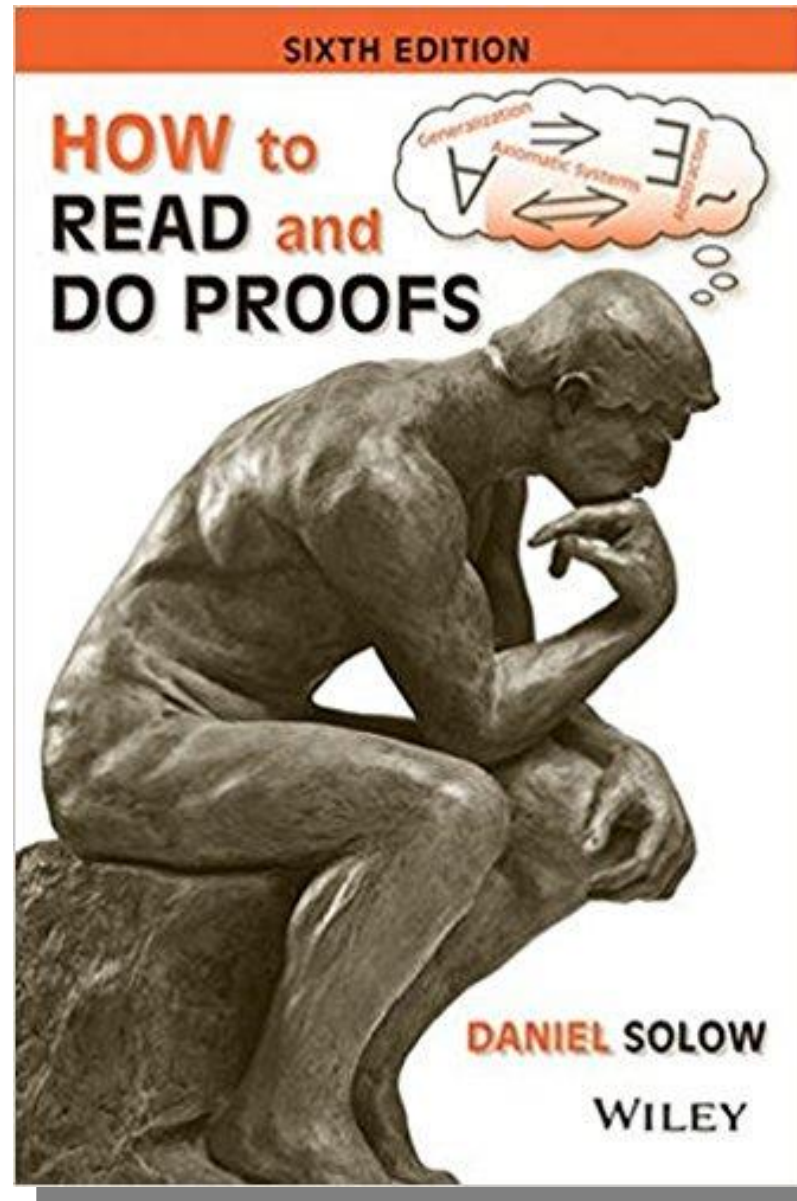
The problem sets throughout the quarter will have some programming assignments. We'll also reference some concepts from CS106B/X, particularly recursion, throughout the quarter.

There aren't any math prerequisites for this course – high-school algebra should be enough!

Problem Set 0

- Your first assignment, Problem Set 0, goes out today. It's due Thursday at 11:59PM.
- You'll need to get your development environment set up, though there's no actual coding involved.
- It covers a few bits of adminstrivia that are important but easily covered offline.

Recommended Reading



Online Course Notes


CS103

Handouts ▾

Assignments ▾

Exams ▾

HW Tools ▾

 Lecture Schedule



CS103: Mathematical Foundations of Computing


Summer 2020


Monday/Wednesday/Friday 9:30am to 11:20am PDT on [Zoom](#) with recordings found on [Canvas](#).


RESOURCES


 [Office Hours and Staff](#)


[Bios](#)

 [Lecture Schedule](#)


 [Lecture Videos \(Canvas\)](#)

 [QT Creator](#)

 [Submit \(Gradescope\)](#)

 [Q&A \(Campuswire\)](#)

TEXTBOOKS

 [Course Reader](#)

 [How to Read & Do Proofs](#)

ASSIGNMENTS

[0 - Warmup](#)

ANNOUNCEMENTS

Welcome to Summer Quarter!

3 hours ago

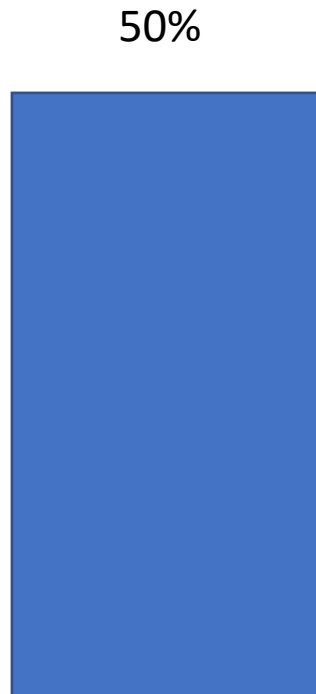
Welcome from instructor Ryan Smith and the rest of the teaching team!

Here is the basic information for getting started in the course.

- The website is still under construction so some links may be dead or out of date. All information should be correct by the start of class.
- Lectures are live on [Zoom](#) from 9:30am to 11:20am PDT every Monday, Wednesday, Friday. You'll need to make a Zoom account with your Stanford email address. Recordings can be found on [Canvas](#).
- Assignment 0 has been released and can be found [here](#). It is due Thursday at 11:59pm PDT.
- CS 106B is a corequisite, so you need to have already taken it or be currently taking it.
- If any of you anticipate that any of the adjustments we've made to move lecture, office hours, exams, etc online this quarter will not work well for you for whatever reason, please email the staff list to let us know so we can find a workaround: cs103-sum1920-staff@lists.stanford.edu.

Grading

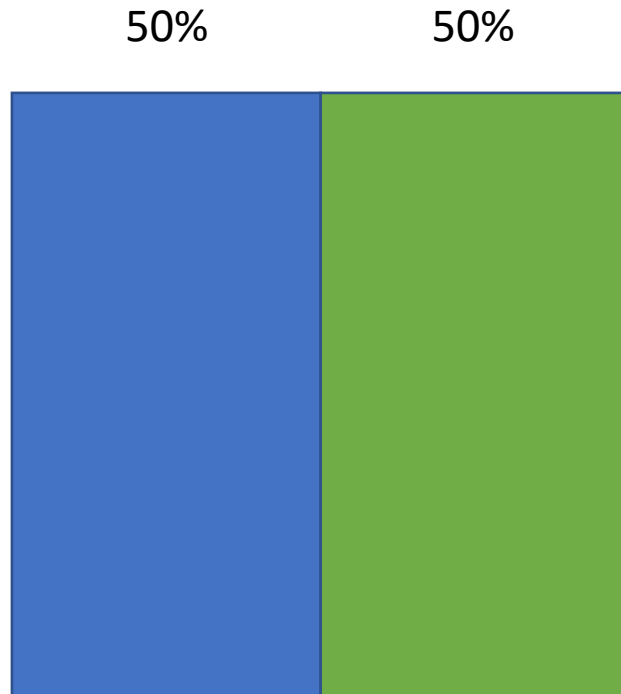
Grading



Seven Problem Sets

Problem sets may be completed individually or in pairs.

Grading



Midterm and Final
Each worth 25%*
Thursday, July 23rd and Friday,
August 14th.

Seven Problem Sets
Problem sets may be completed
individually or in pairs.

Current Events

Life is stressful right now.

Extraordinary events are happening.

If you need assistance, come and talk to us.

How to Succeed in CS103

Proof-Based Mathematics

- Most high-school math classes – with the exception of geometry – focus on *calculation*.
- CS103 focuses on *argumentation*.
- Your goal is to *see why things are true*, not *check that they work in a few cases*.
- Be curious! Ask questions. Try things out on your own. You'll learn this material best if you engage with it and refuse to settle for a “good enough” understanding.

Mental Traps to Avoid

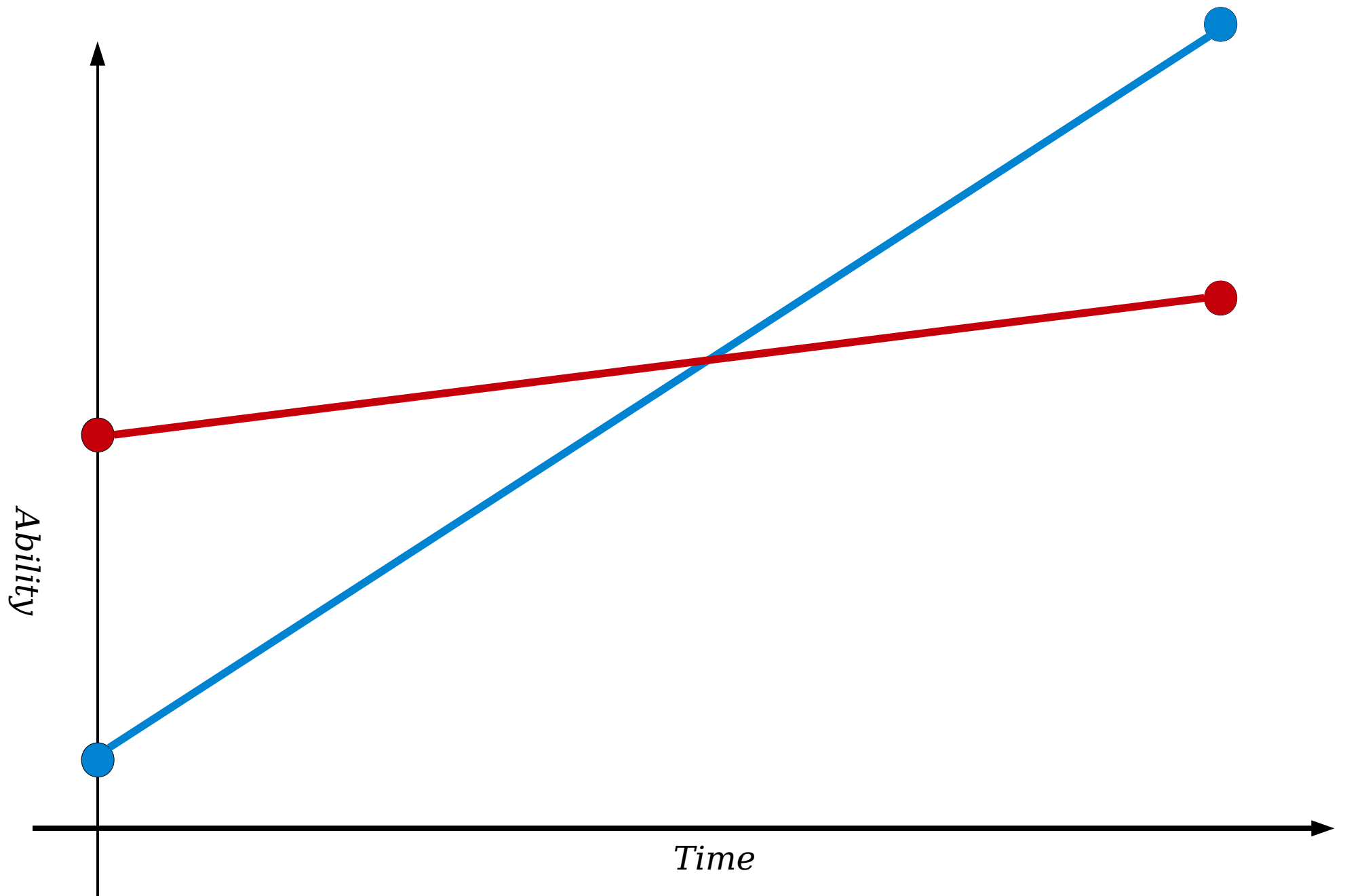
- “Everyone else has been doing math since before they were born and there is no way I'll ever be as good as them.”
- “A small minority of people are math geniuses and everyone else has no chance at being good at math.”
- “Being good at math means being able to instantly solve any math problem thrown at you.”

Mental Traps to Avoid

- “Everyone else has been doing math since before they were born and there is no way I'll ever be as good as them.”

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“Being good at math means being able to instantly solve any math problem thrown at you.”



“A little slope makes up for a lot of y-intercept.”
- John Ousterhout

Mental Traps to Avoid

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Pro Tip #1:

Never Confuse Experience for Talent

Pro Tip #2:

Have a Growth Mindset

Mental Traps to Avoid

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- “Being good at math means being able to instantly solve any math problem thrown at you.”

My Advice

- Question everything!
- Attend lecture.
- Study strategically and intentionally.
- Stick with it, but know when to get help.

We've got a big journey ahead of us.

Let's get started!

“CS103 students”

“All the computers on the
Stanford network”

“Cool people”

“The chemical elements”

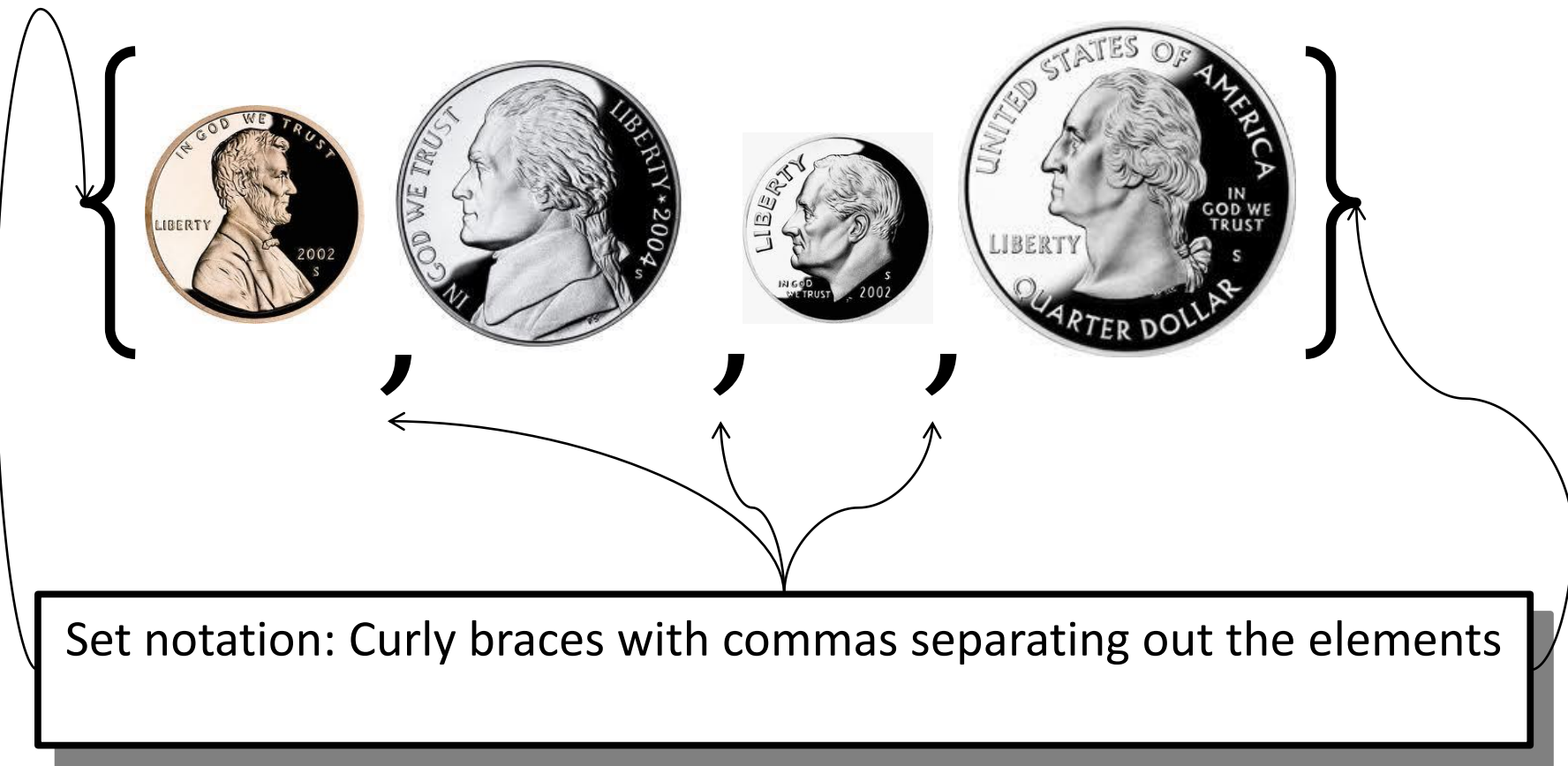
“Cute animals”

“US coins”

A ***set*** is an unordered collection of distinct objects, which may be anything (including other sets).



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Two sets are equal when they have exactly the same contents, ignoring order.



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Two sets are equal when they have exactly the same contents, ignoring order.

Repeated elements in a set are ignored.



Repeated elements in a set are ignored.



Repeated elements in a set are ignored.



These are the same set!



Repeated elements in a set are ignored.

$\{\}$ $=$ \emptyset 

The **empty set** contains no elements.

We use this symbol to denote the empty set.

This is a number.



1

?
≡

This is a set. It contains a number.



{ 1 }

Are these objects equal to one another?

This is a number.



1

≠

This is a set. It contains a number.



{ 1 }

Are these objects equal to one another?

This set contains
nothing at all.

\emptyset

This set has one element, which
happens to be the empty set.

$\stackrel{?}{=}$

$\{\emptyset\}$

Are these objects equal to one another?

This set contains nothing at all.

\emptyset

\neq

This set has one element, which happens to be the empty set.

$\{\emptyset\}$

Are these objects equal to one another?

Membership

Membership



Membership



Is



Membership



Is



Membership



Is



Membership



Is



Set Membership

Given a set S and an object x , we write

$$x \in S$$

if x is contained in S , and

$$x \notin S$$

otherwise.

If $x \in S$, we say that x is an ***element*** of S .

Given any object x and any set S , either $x \in S$ or $x \notin S$.

Infinite Sets

- Some sets contain *infinitely many* elements!
- The set $\mathbb{N} = \{ 0, 1, 2, 3, \dots \}$ is the set of all the ***natural numbers***.
- Some mathematicians don't include zero; in this class, assume that 0 is a natural number.
- The set $\mathbb{Z} = \{ \dots, -2, -1, 0, 1, 2, \dots \}$ is the set of all the ***integers***.
- Z is from German "Zahlen."
- The set \mathbb{R} is the set of all ***real numbers***.
- $e \in \mathbb{R}$, $\pi \in \mathbb{R}$, $4 \in \mathbb{R}$, etc.

Describing Complex Sets

Here are some English descriptions of infinite sets:

“The set of all even natural numbers.”

“The set of all real numbers less than 137.”

“The set of all negative integers.”

To describe complex sets like these mathematically, we'll use ***set-builder notation***.

Even Natural Numbers

$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$

Even Natural Numbers

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Even Natural Numbers

$$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$



The set of all n

Even Natural Numbers

$$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

The set of all n

where

Even Natural Numbers

$$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

The set of all n

where

n is a natural number

Even Natural Numbers

$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$

The set of all n

where

n is a natural number

and n is even

Even Natural Numbers

$\{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$

The set of all n

where

n is a natural number

and n is even

$\{ 0, 2, 4, 6, 8, 10, 12, 14, 16, \dots \}$

Set Builder Notation

A set may be specified in ***set-builder notation***:

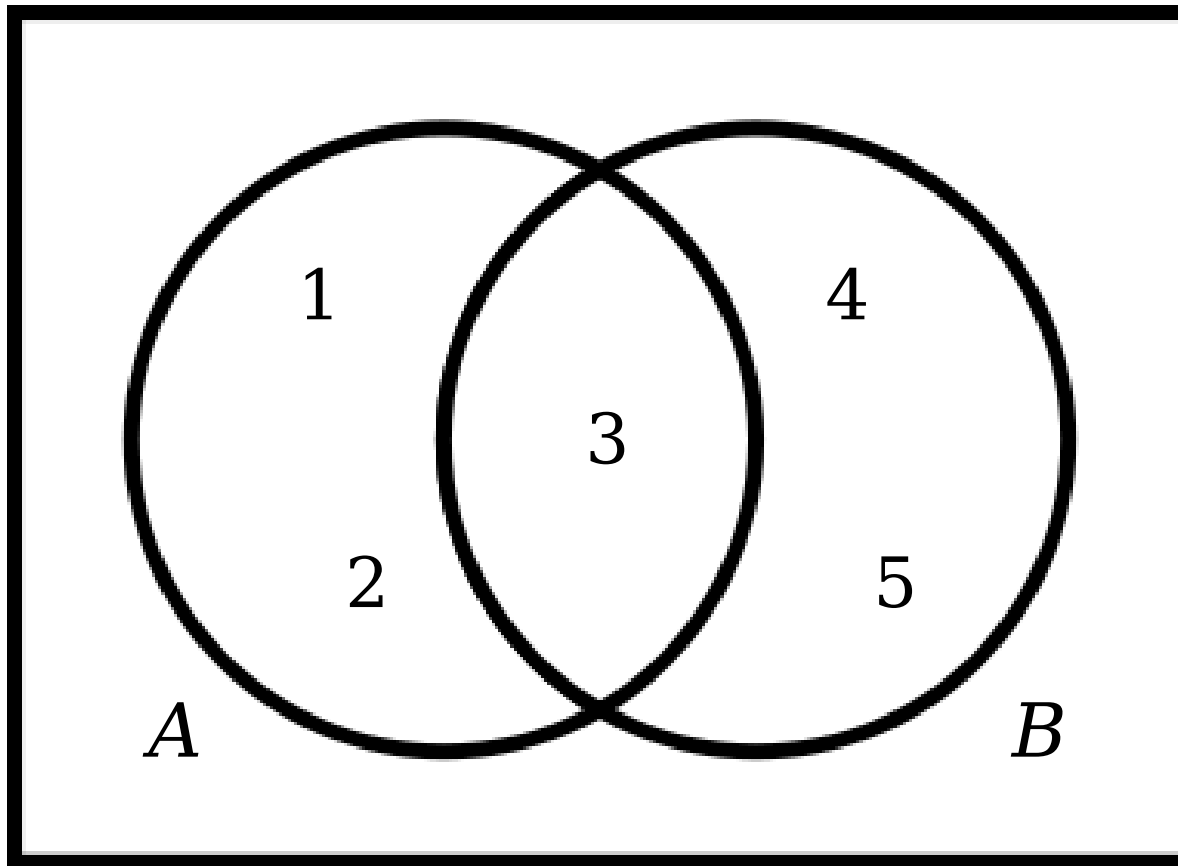
$$\{ x \mid \text{some property } x \text{ satisfies} \}$$

For example:

- $\{ r \mid r \in \mathbb{R} \text{ and } r < 137 \}$
- $\{ n \mid n \text{ is an even natural number} \}$
- $\{ S \mid S \text{ is a set of US currency} \}$
- $\{ a \mid a \text{ is cute animal} \}$
- $\{ r \in \mathbb{R} \mid r < 137 \}$
- $\{ n \in \mathbb{N} \mid n \text{ is odd} \}$

Combining Sets

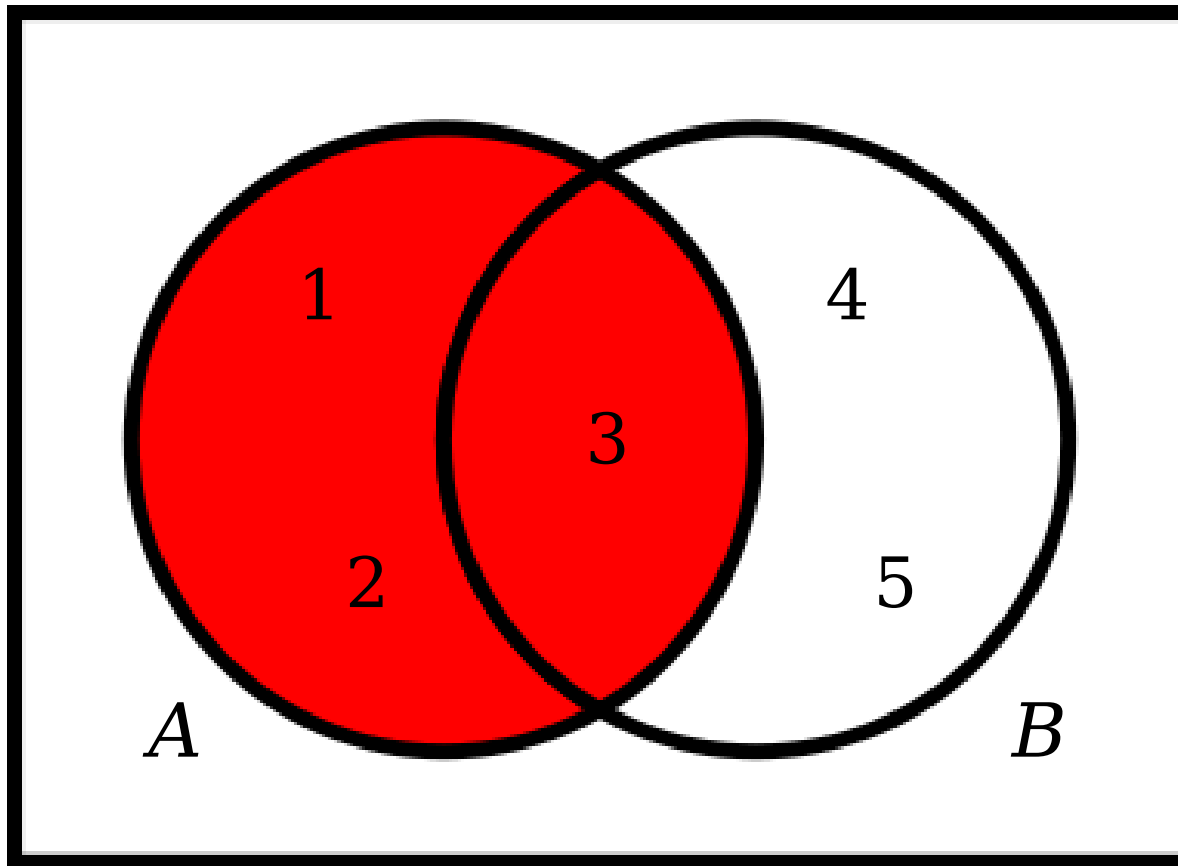
Venn Diagrams



$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

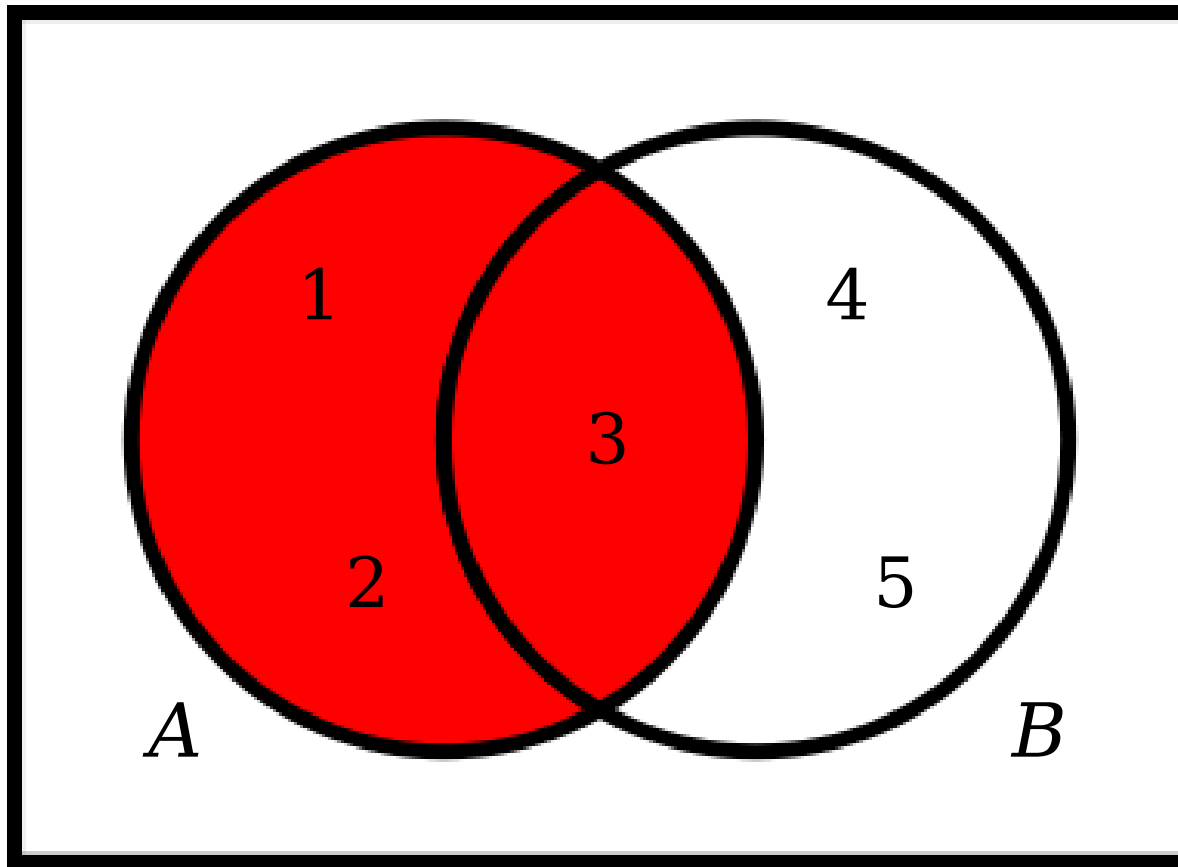
Venn Diagrams



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Venn Diagrams

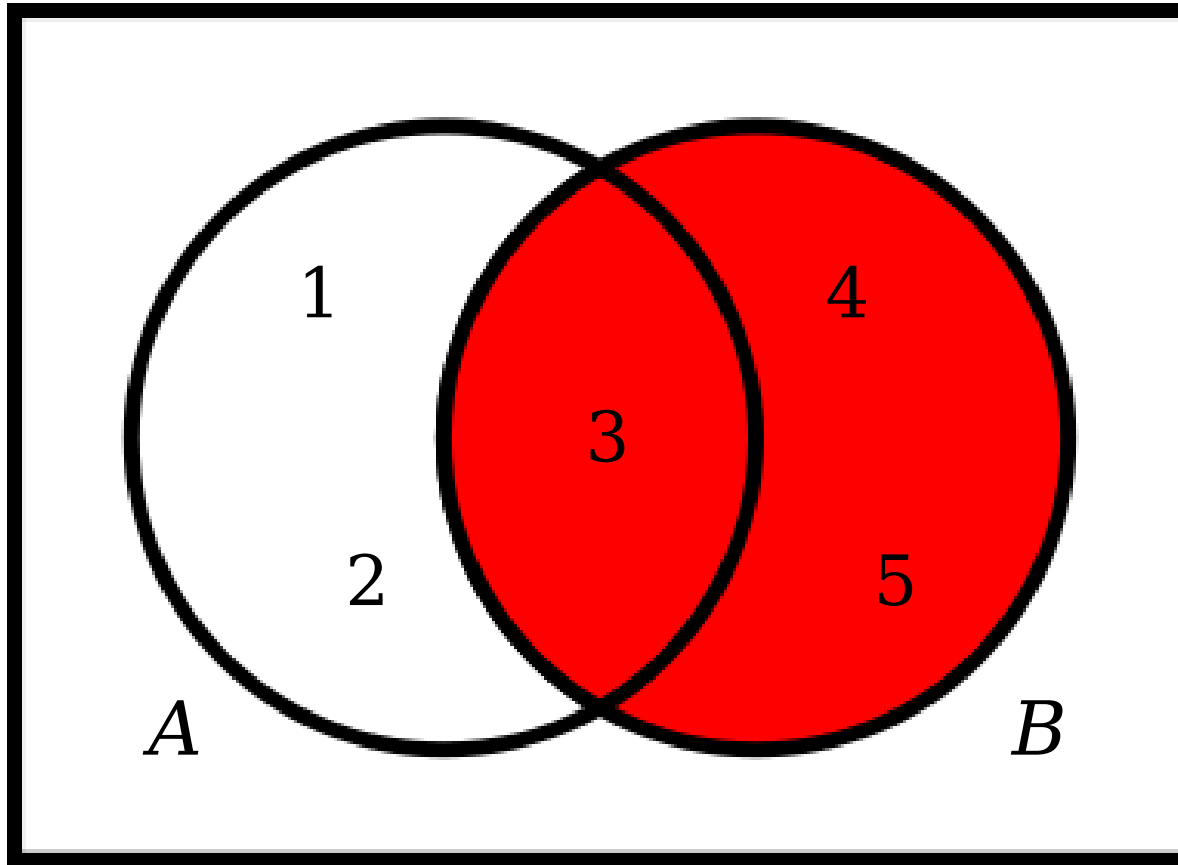


A

$$A = \{ 1, 2, 3 \}$$

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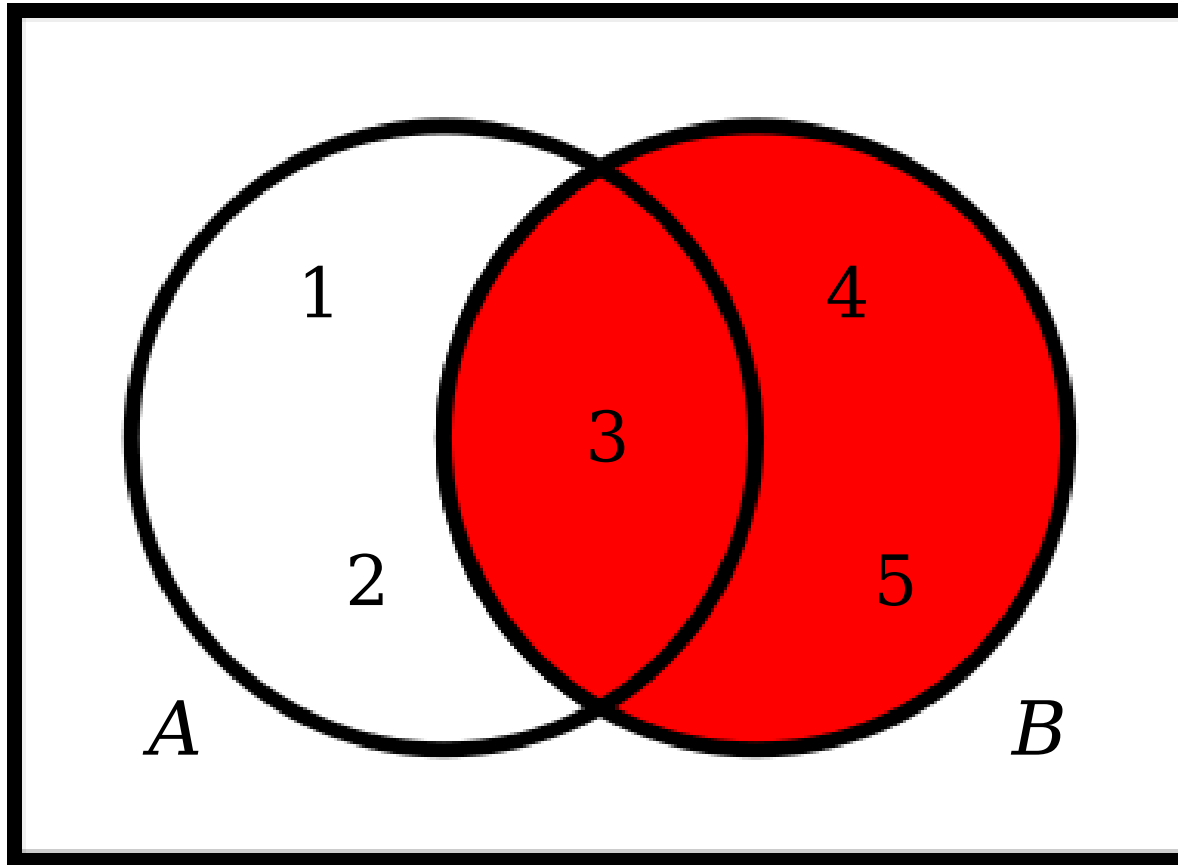
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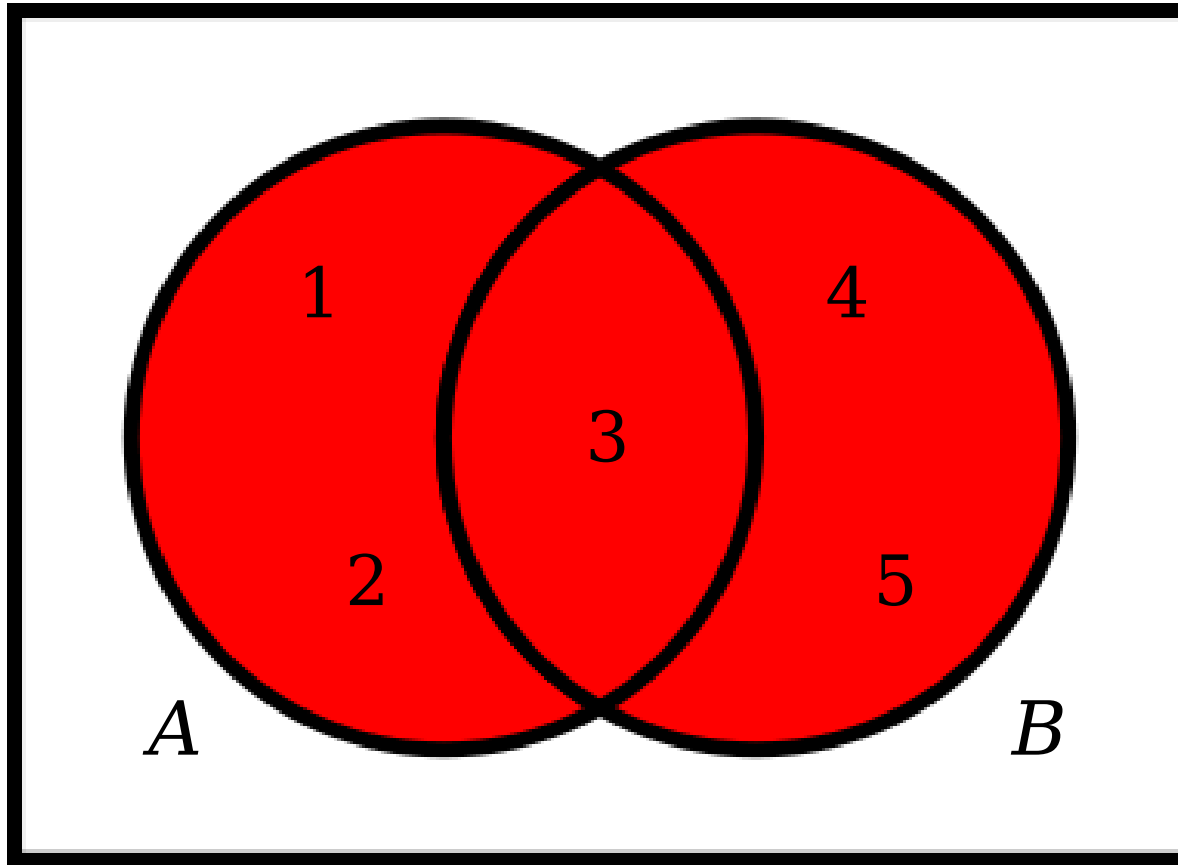
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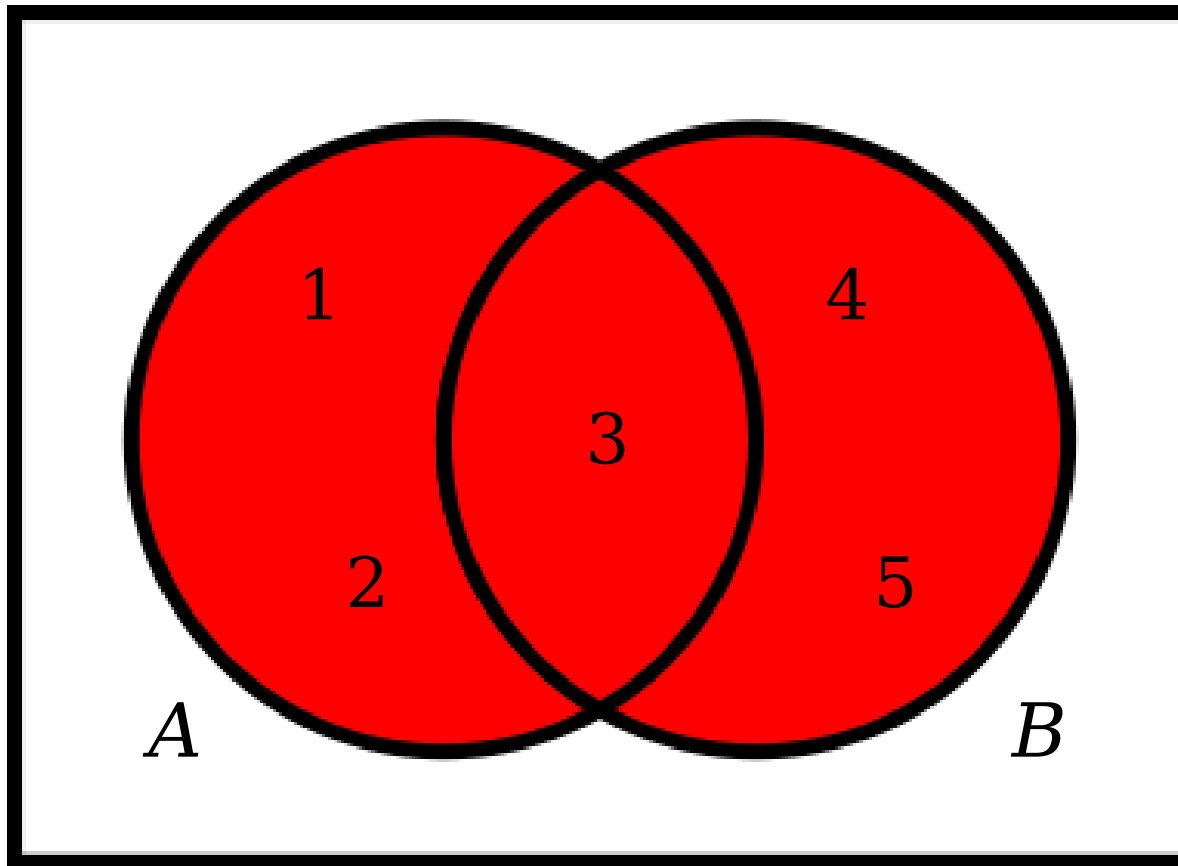
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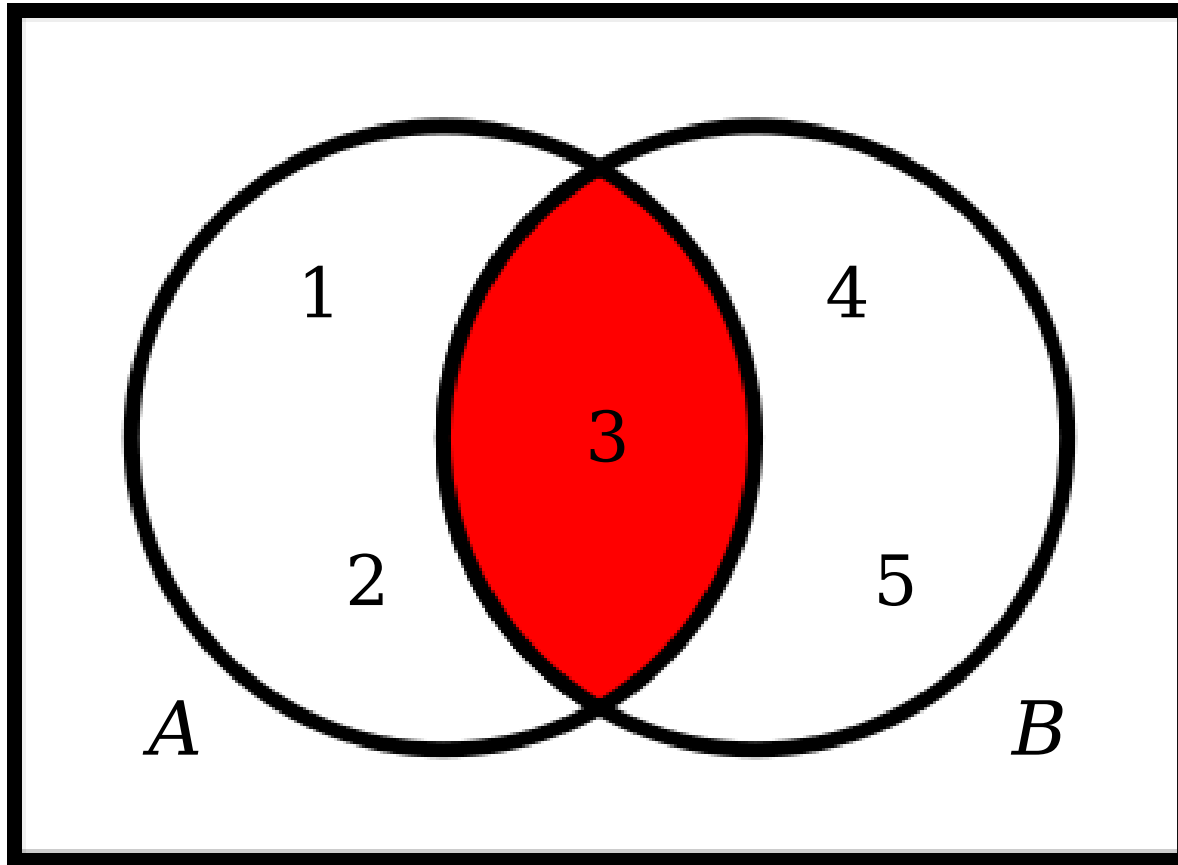


Union
 $A \cup B$
 $\{ 1, 2, 3, 4, 5 \}$

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

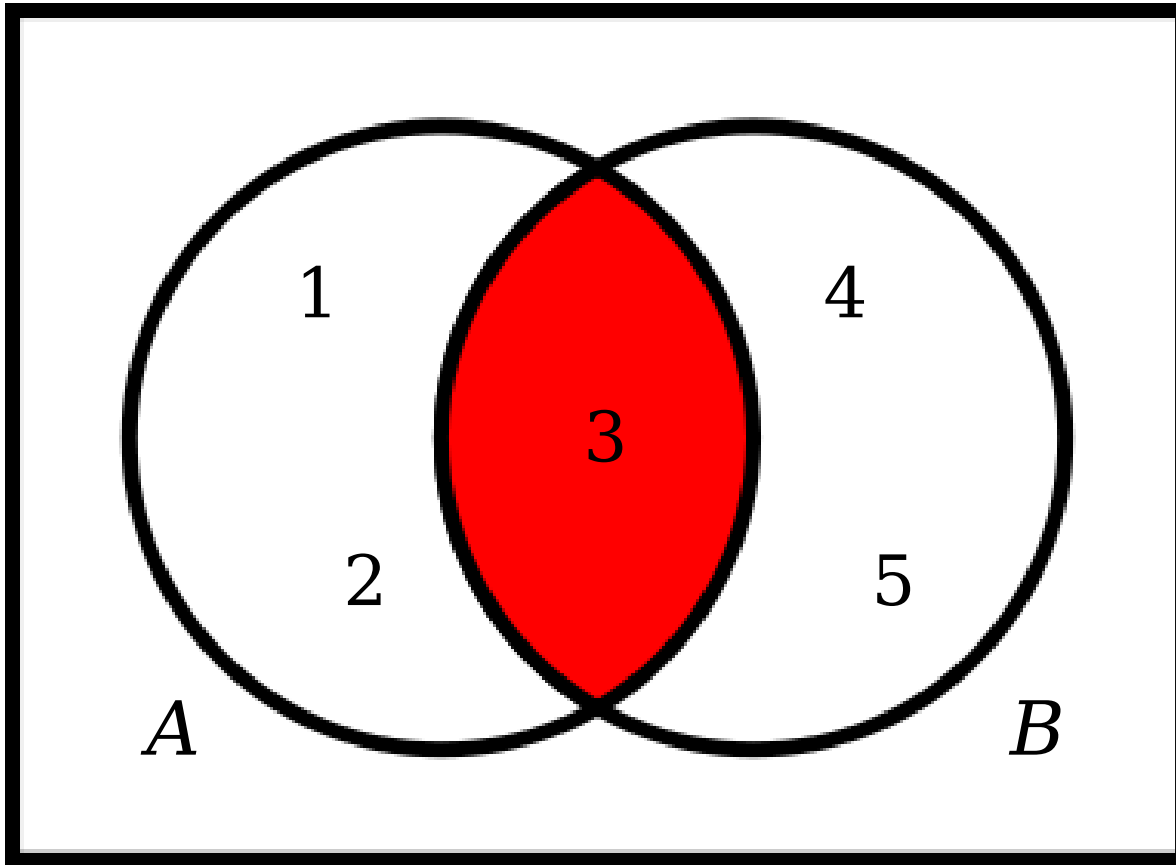
Venn Diagrams



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$$B = \{ 3, 4, 5 \}$$

Venn Diagrams



Intersection

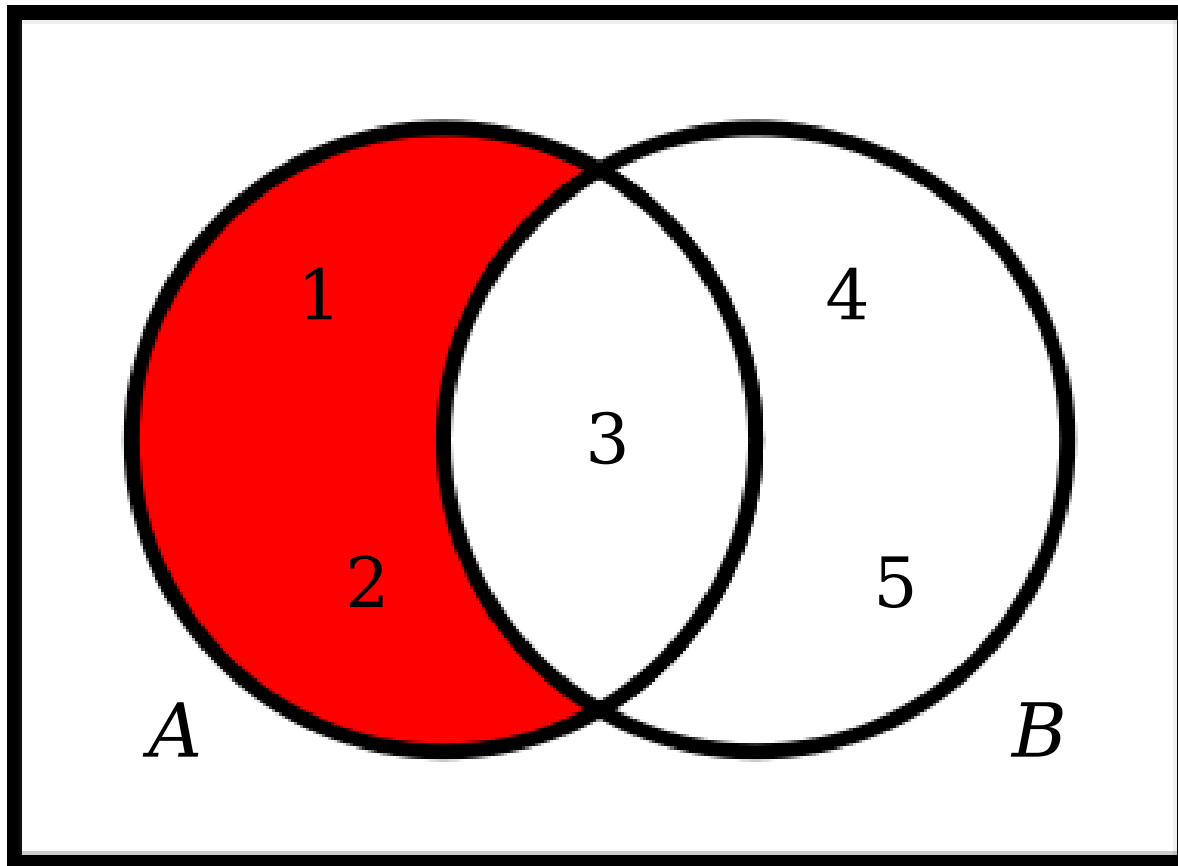
$$A \cap B$$

$$\{ 3 \}$$

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

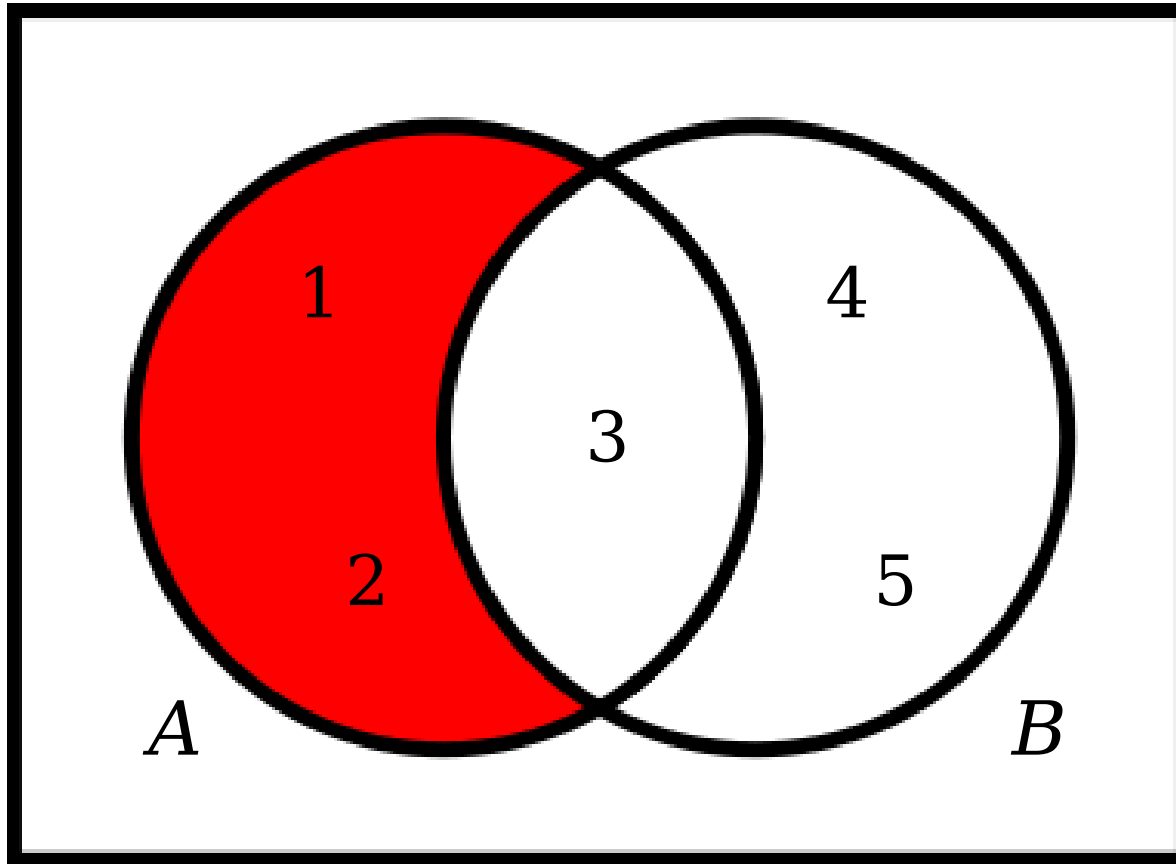
Venn Diagrams



$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

Venn Diagrams



Difference

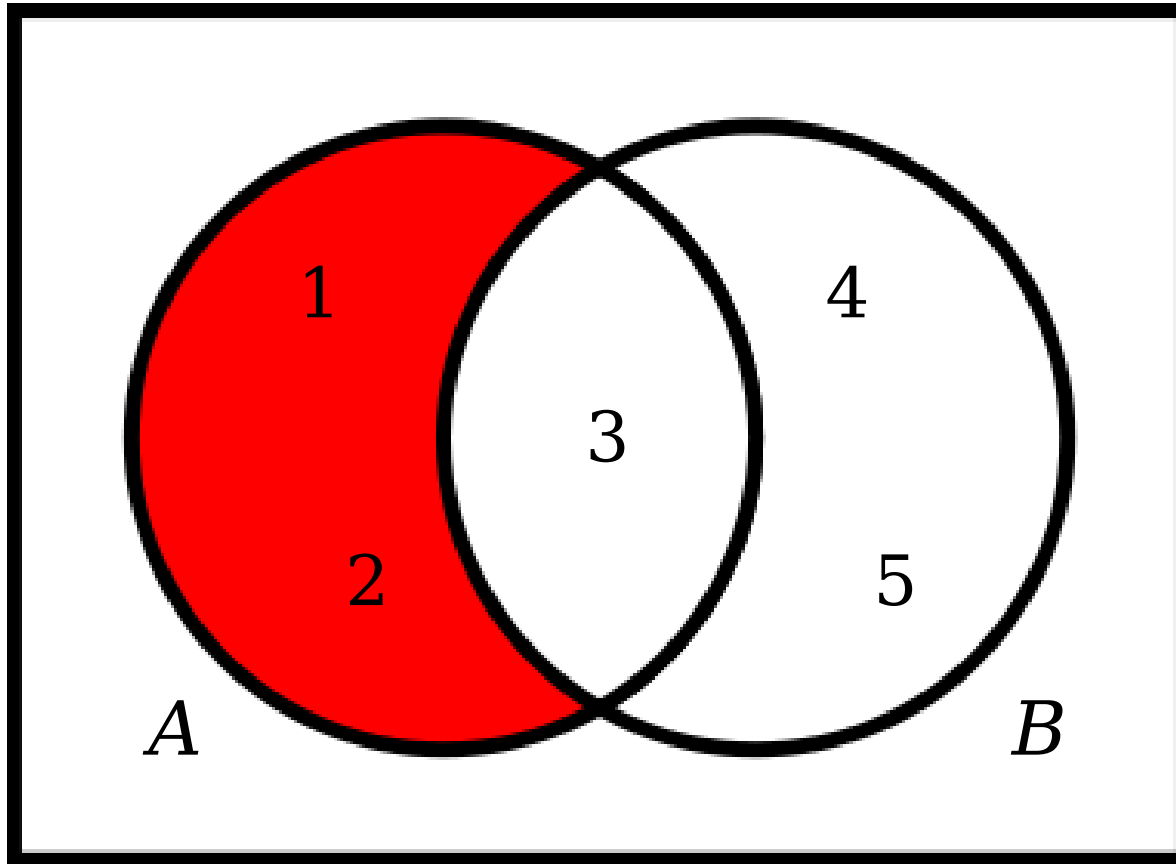
$$A - B$$

$$\{ 1, 2 \}$$

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

Venn Diagrams



Difference

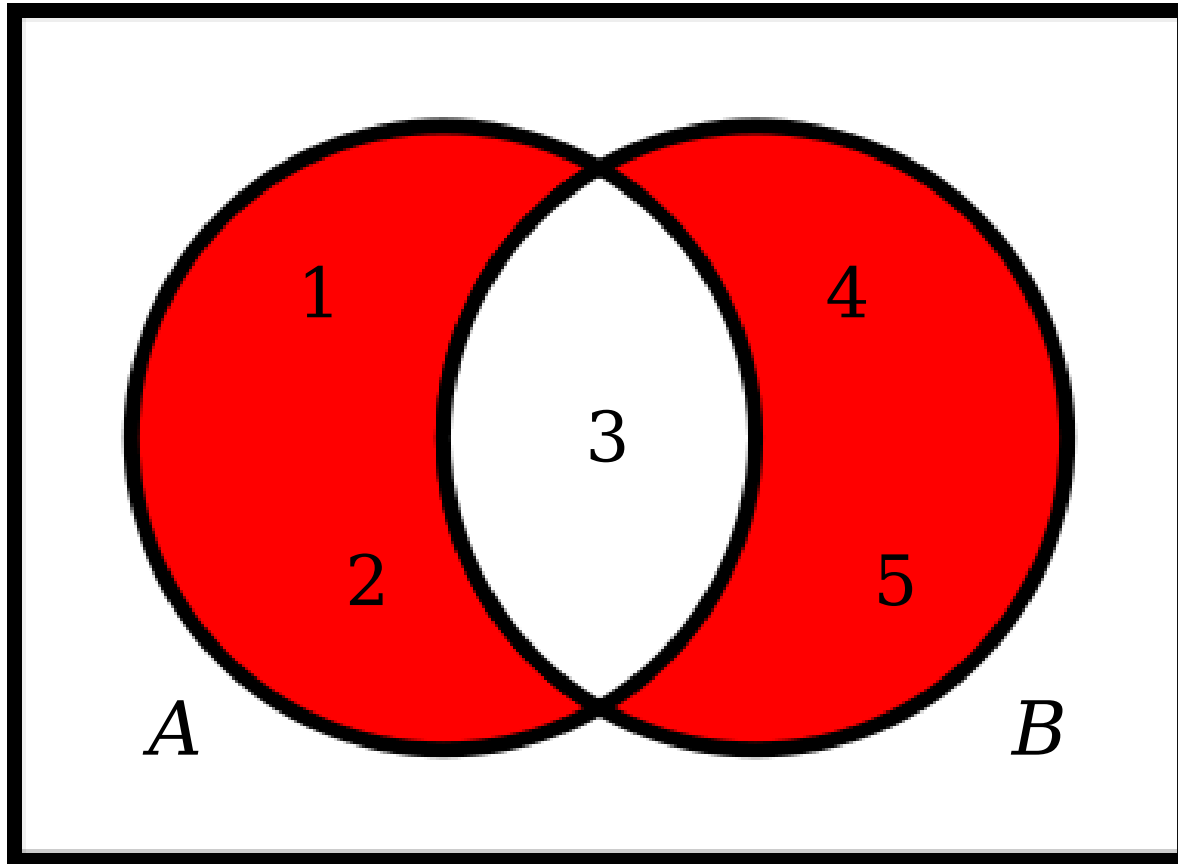
$$A \setminus B$$

$$\{ 1, 2 \}$$

$$A = \{ 1, 2, 3 \}$$

$$B = \{ 3, 4, 5 \}$$

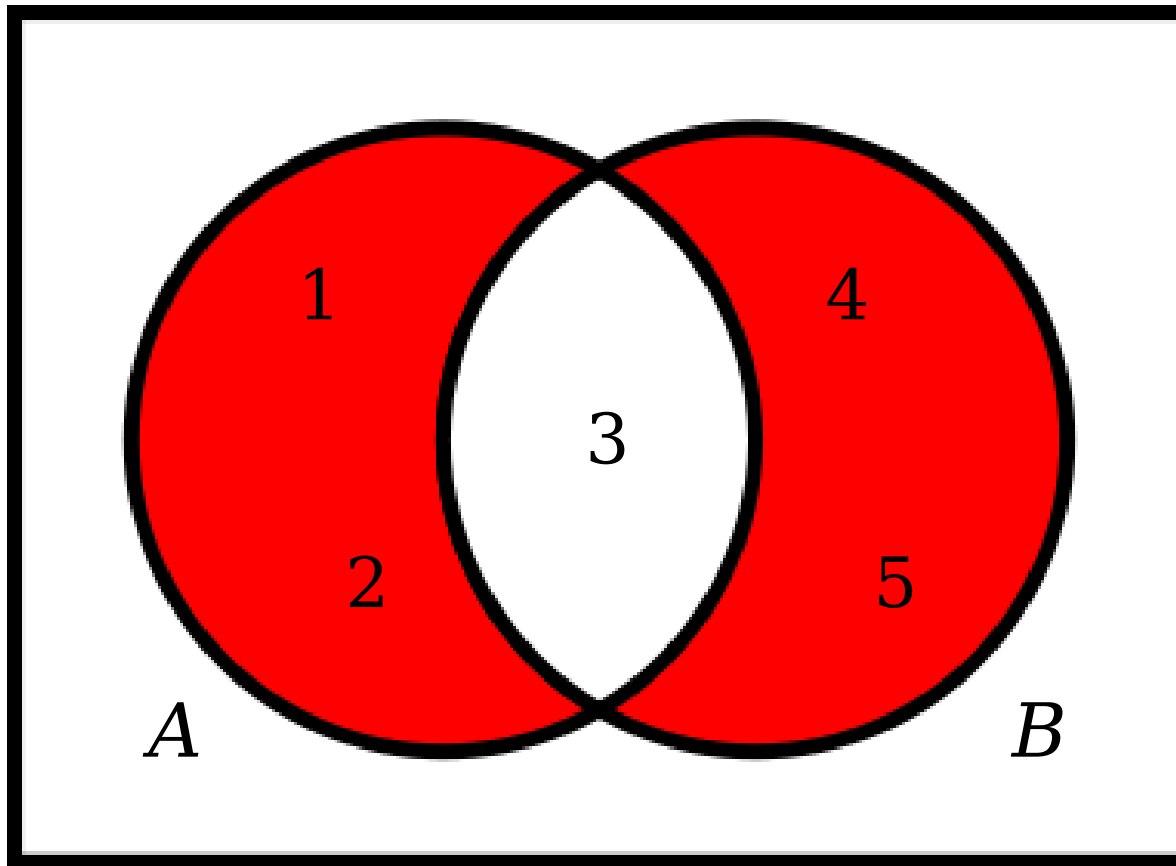
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Venn Diagrams

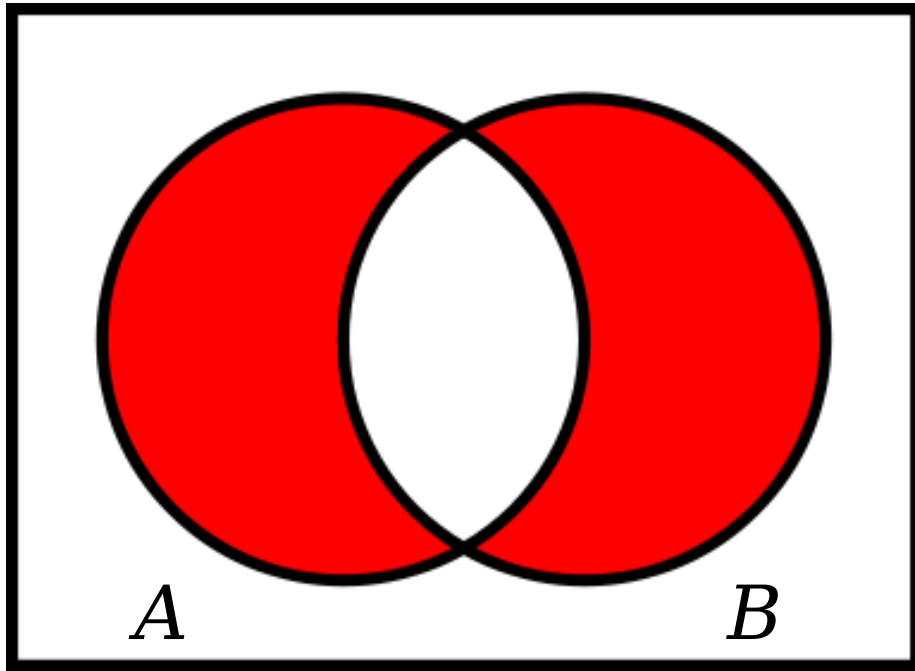


Symmetric
Difference
 $A \Delta B$
 $\{ 1, 2, 4, 5 \}$

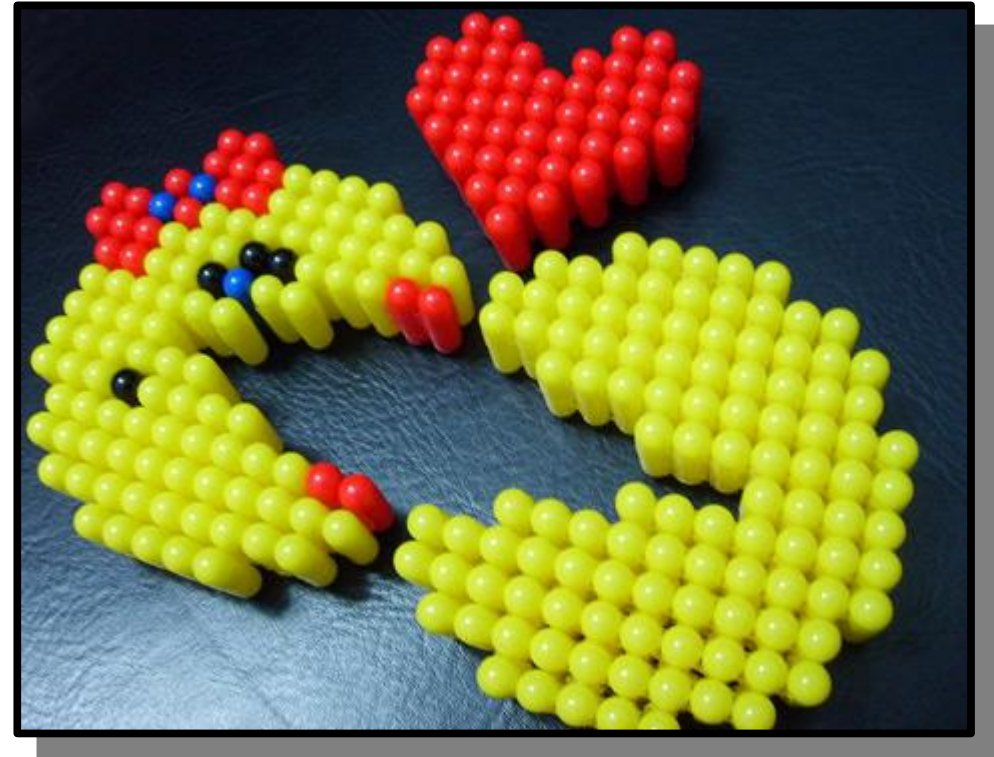
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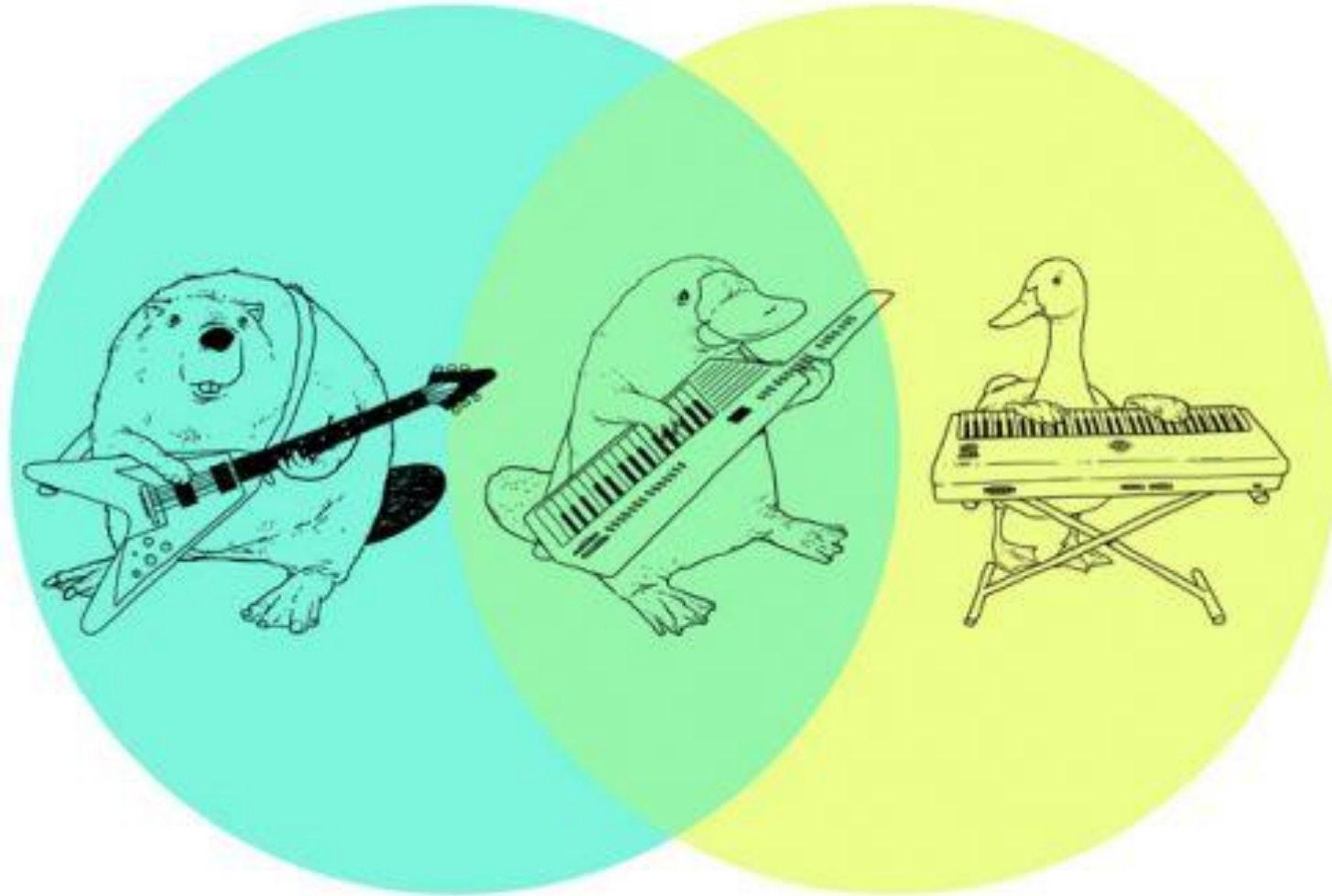
Venn Diagrams



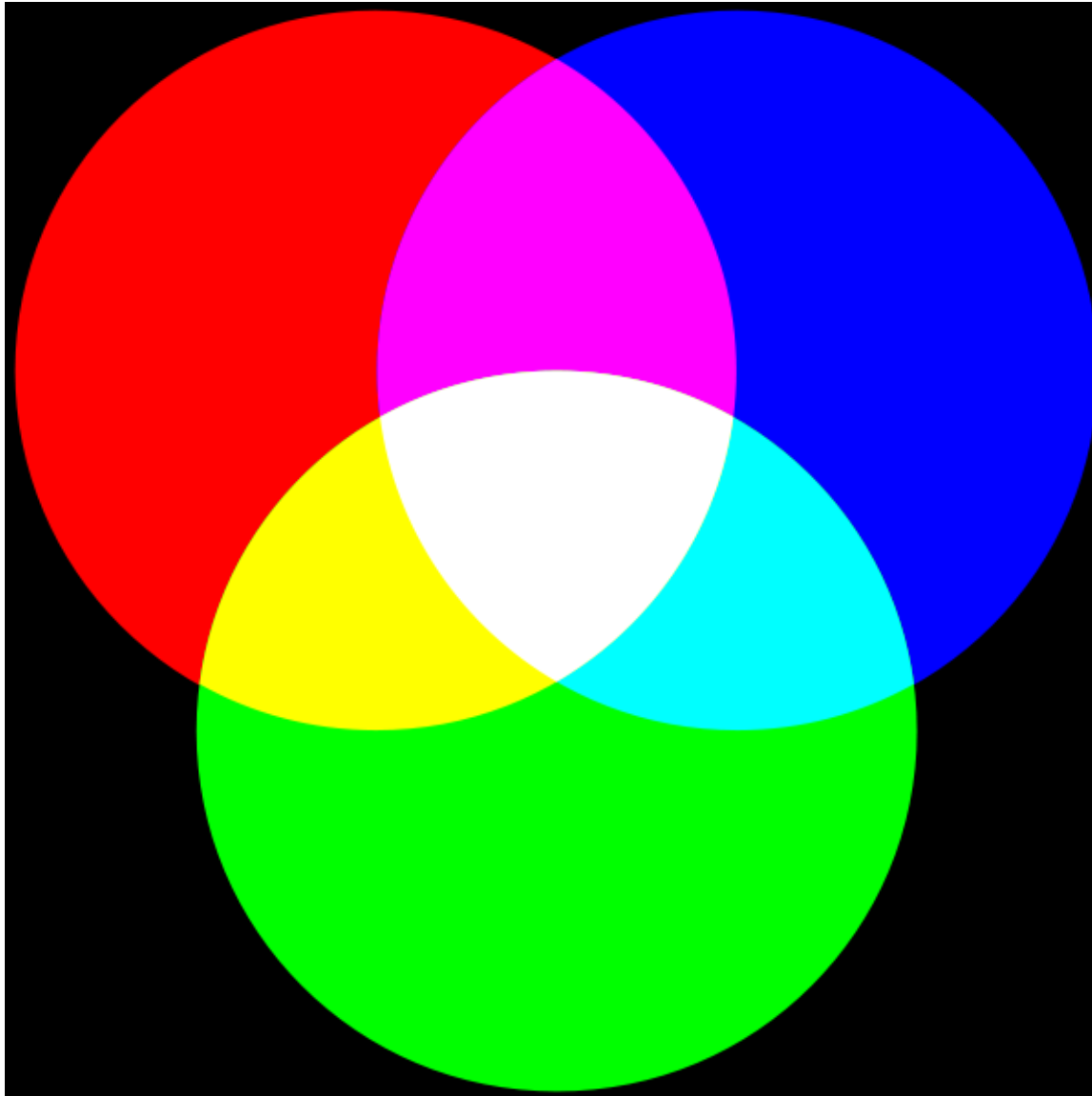
$$A \Delta B$$



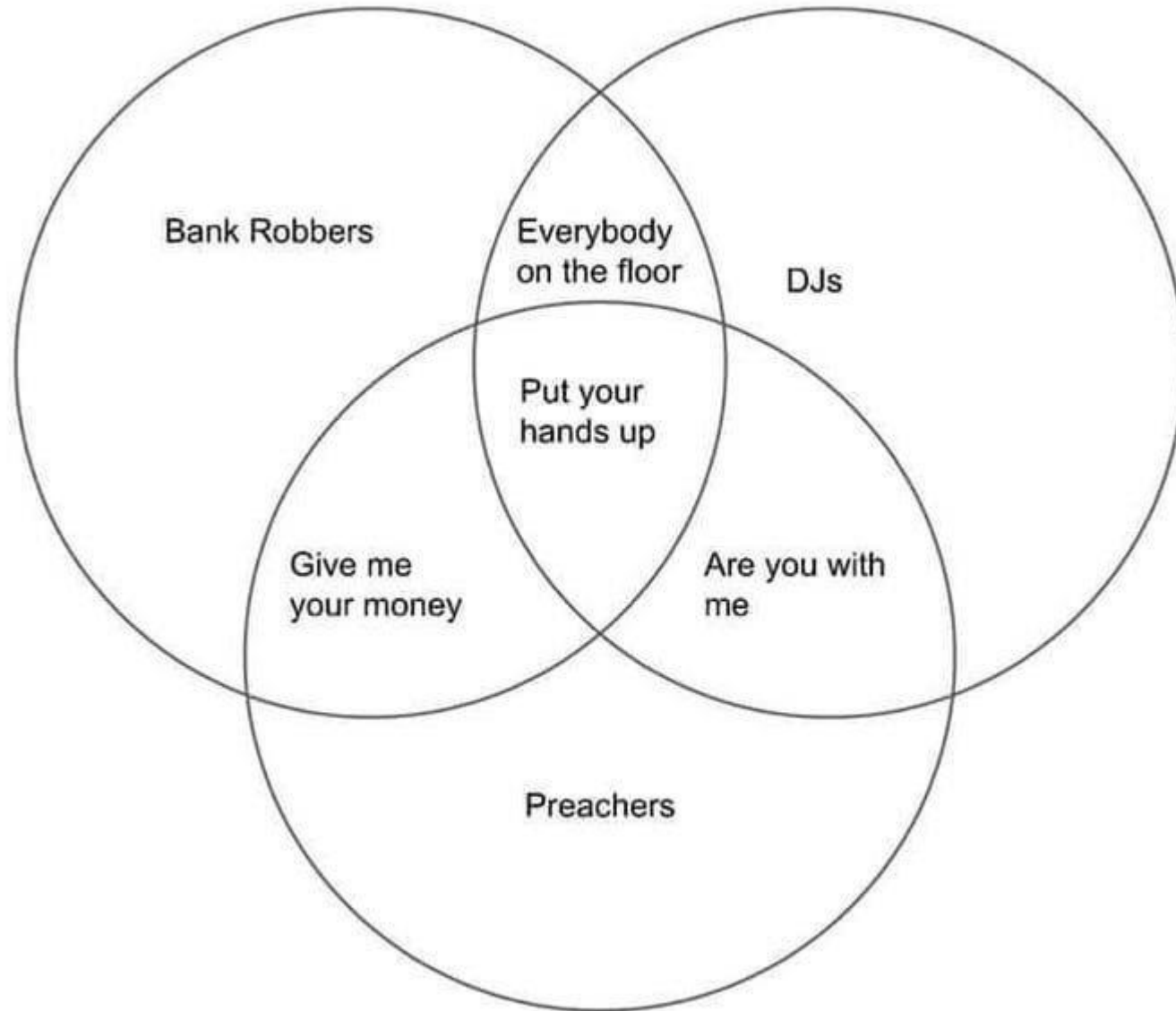
Venn Diagrams



Venn Diagrams for Three Sets



Venn Diagrams for Three Sets

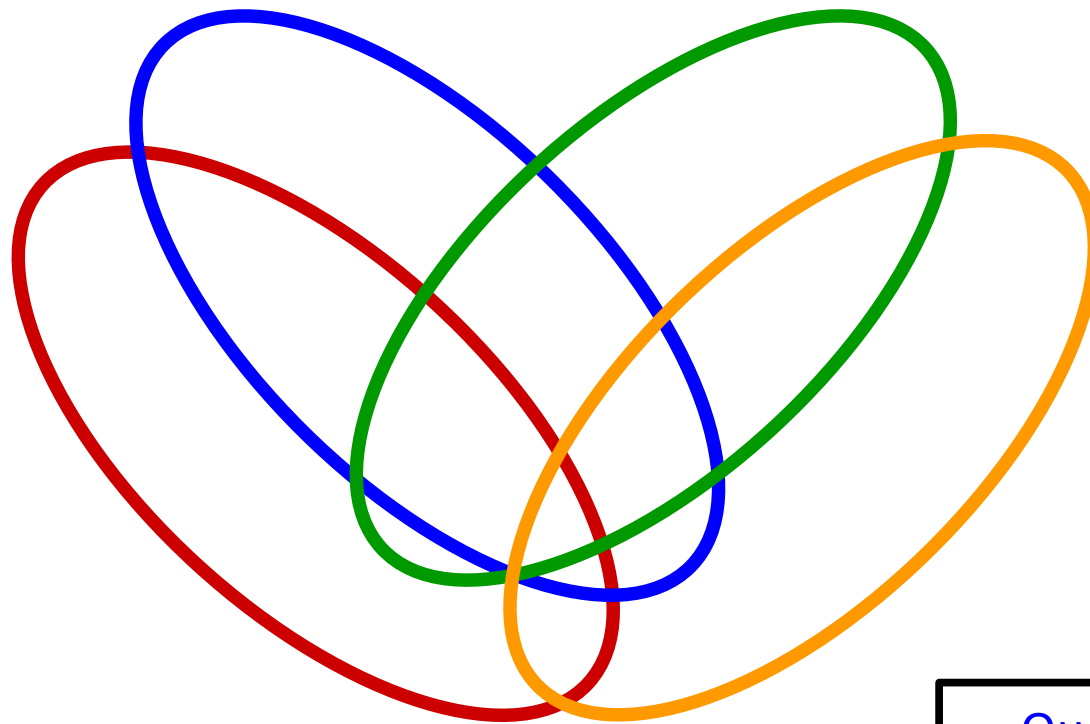


Venn Diagrams for Four Sets

B

C

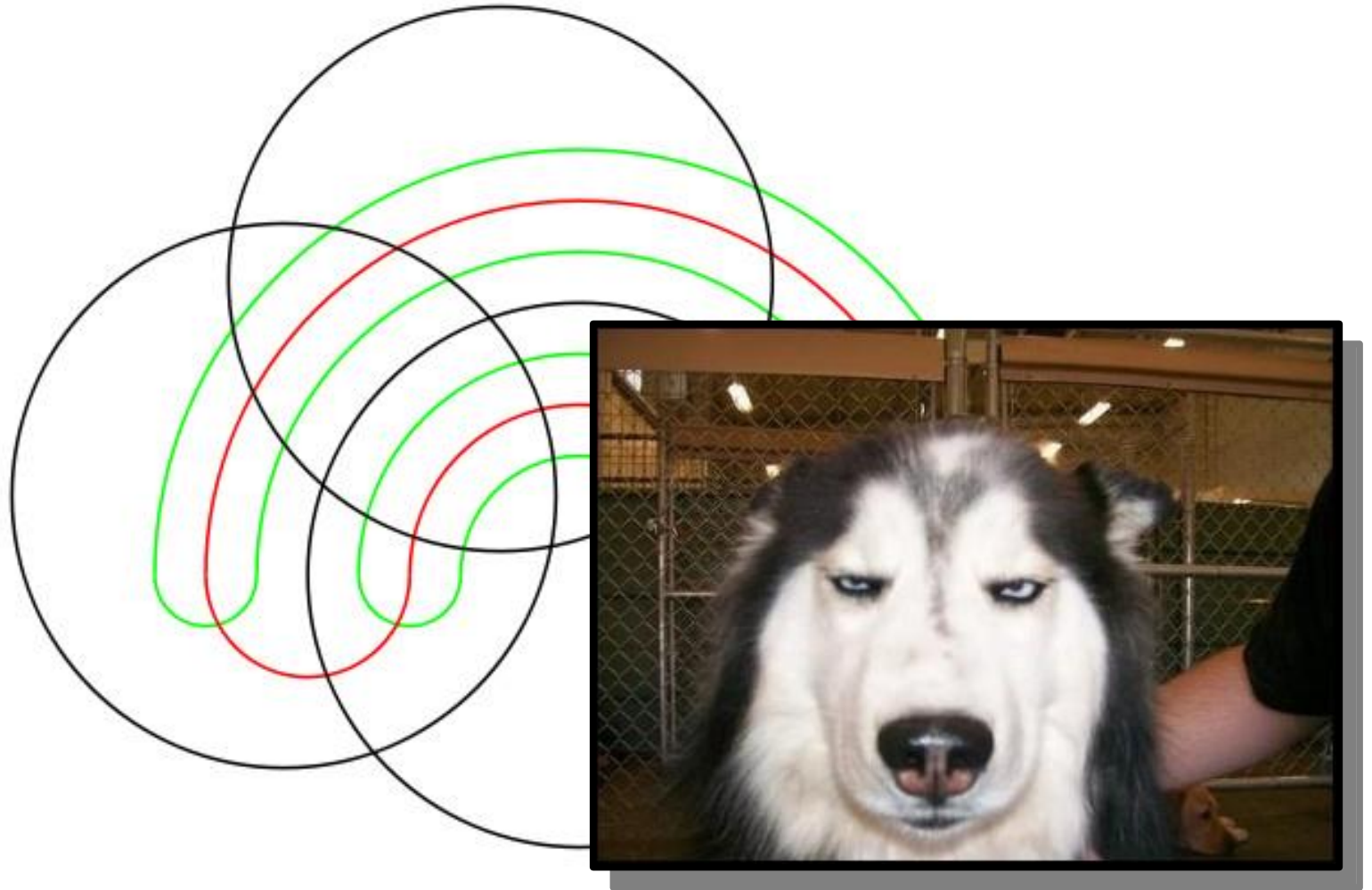
A



D

Question to ponder: why
don't we just draw four
circles?

Venn Diagrams for Five Sets



Venn Diagrams for Seven Sets

<http://moebio.com/research/sevensets/>

Subsets and Power Sets

Subsets

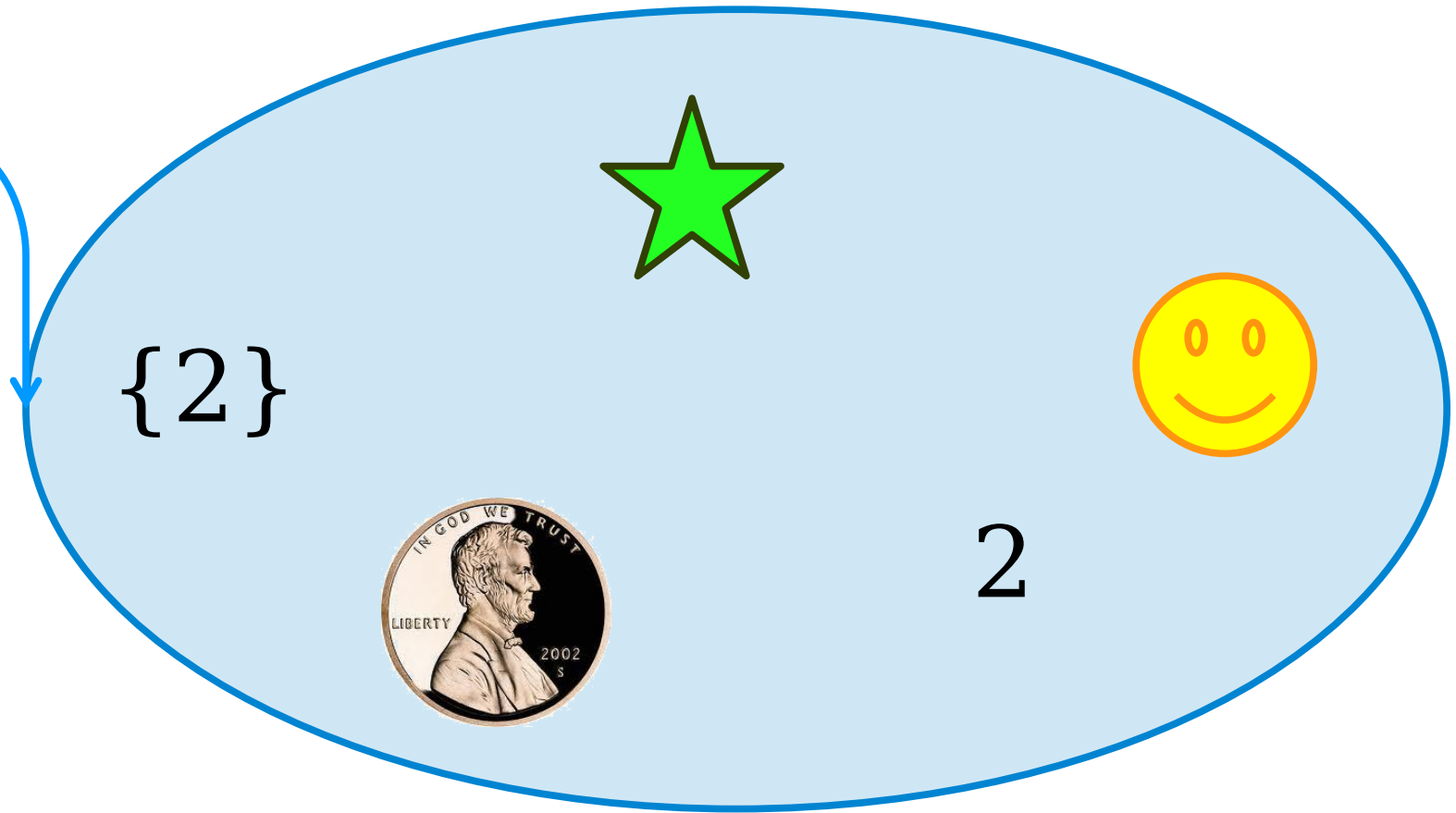
A set S is called a **subset** of a set T (denoted $S \subseteq T$) if all elements of S are also elements of T .

Examples:

- $\{ 1, 2, 3 \} \subseteq \{ 1, 2, 3, 4 \}$
- $\{ c, b \} \subseteq \{ a, b, c, d \}$
- $\{ \text{H, He, Li} \} \subseteq \{ \text{H, He, Li} \}$
- $\mathbb{N} \subseteq \mathbb{Z}$ (*every natural number is an integer*)
- $\mathbb{Z} \subseteq \mathbb{R}$ (*every integer is a real number*)

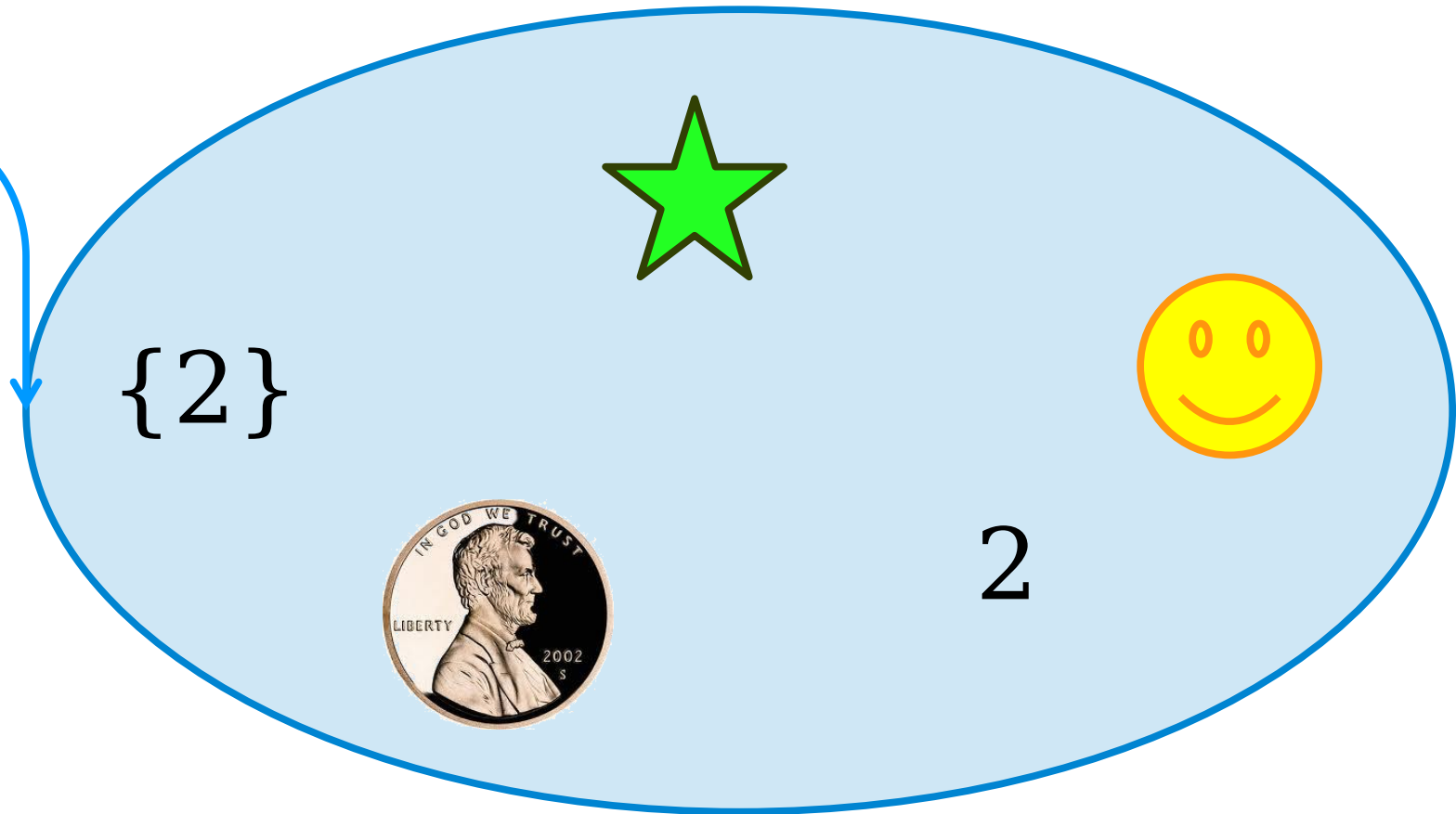
Subsets and Elements

Set S



Subsets and Elements

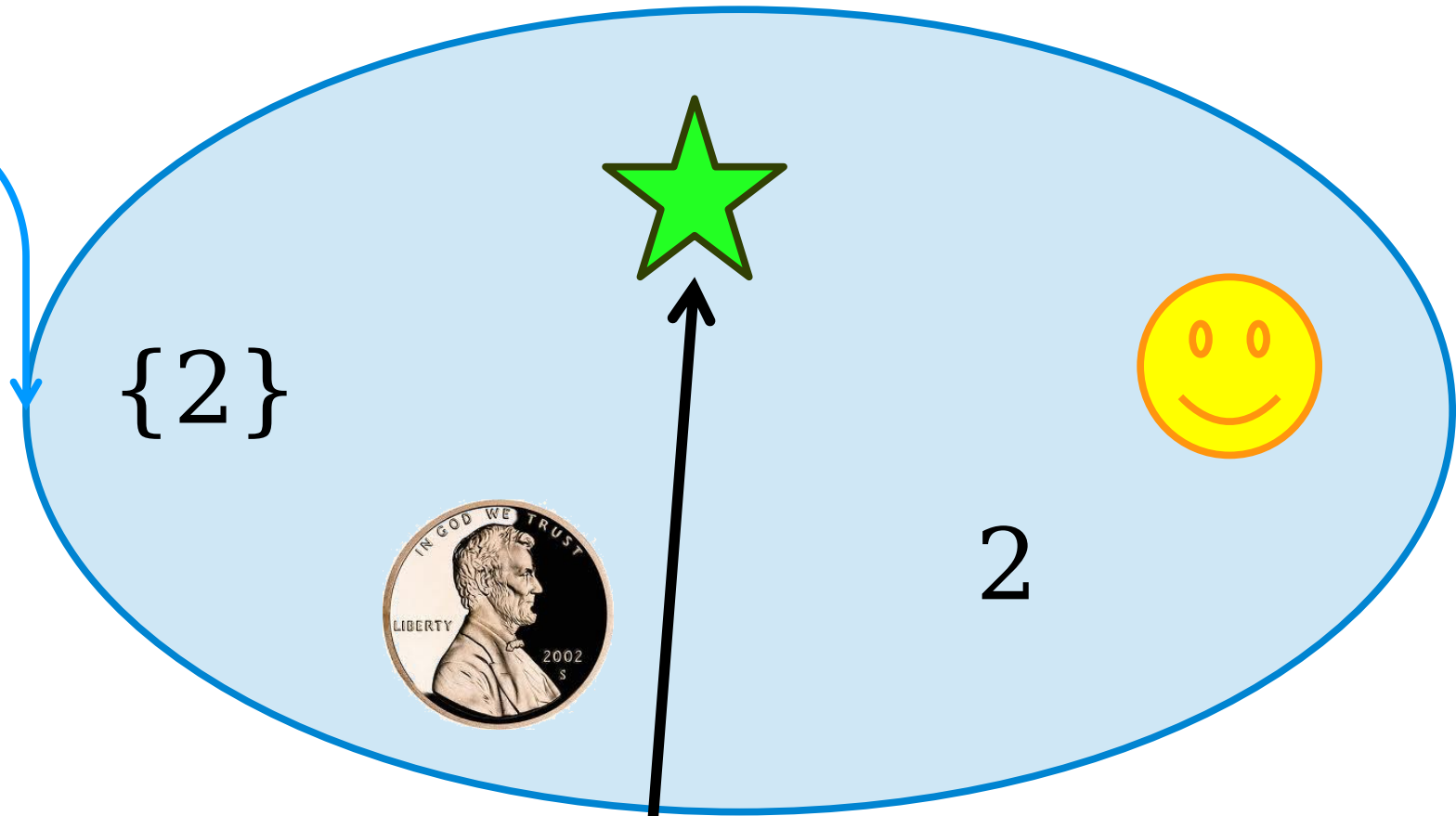
Set S



$$\star \in S$$

Subsets and Elements

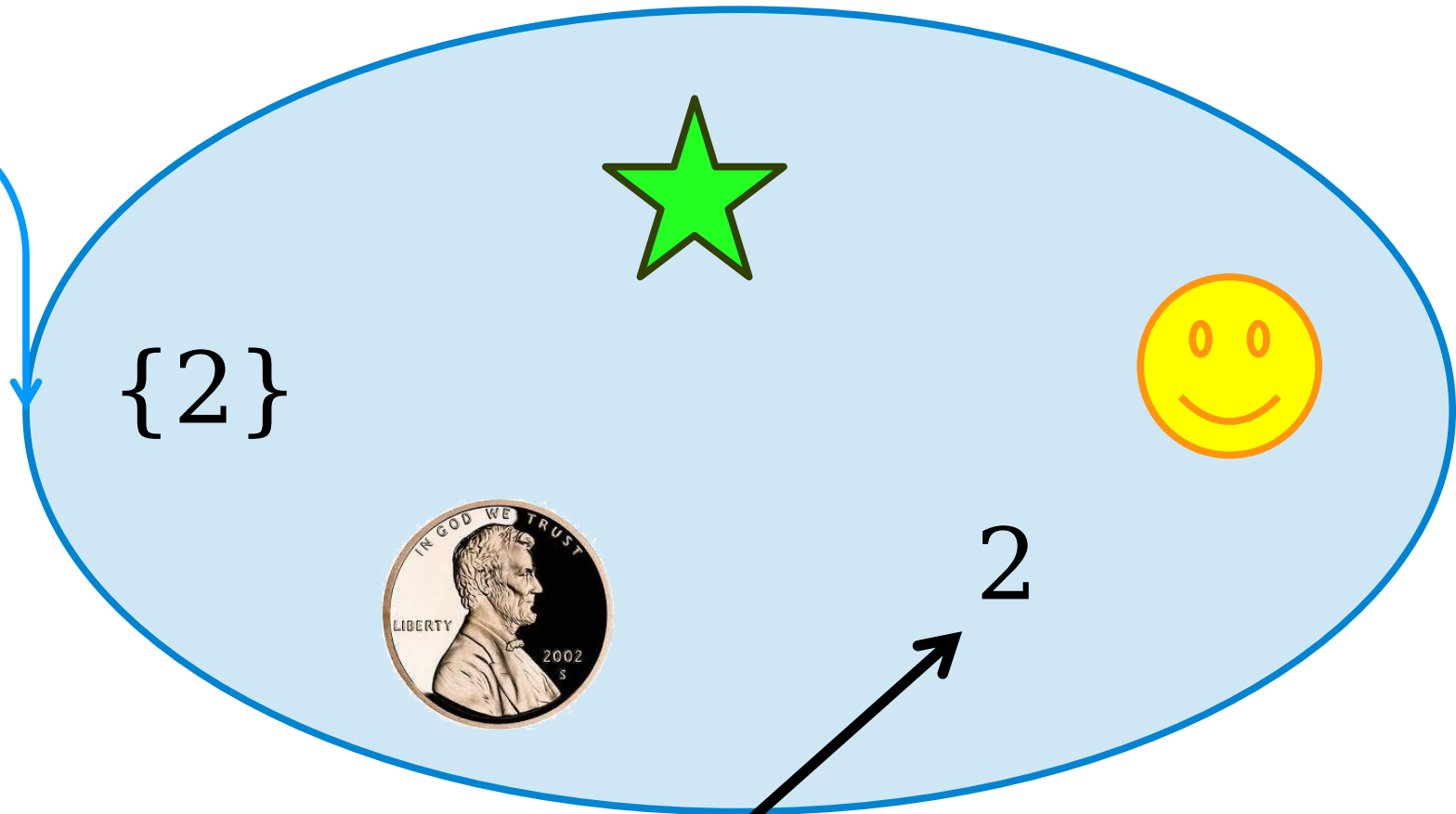
Set S



 $\in S$

Subsets and Elements

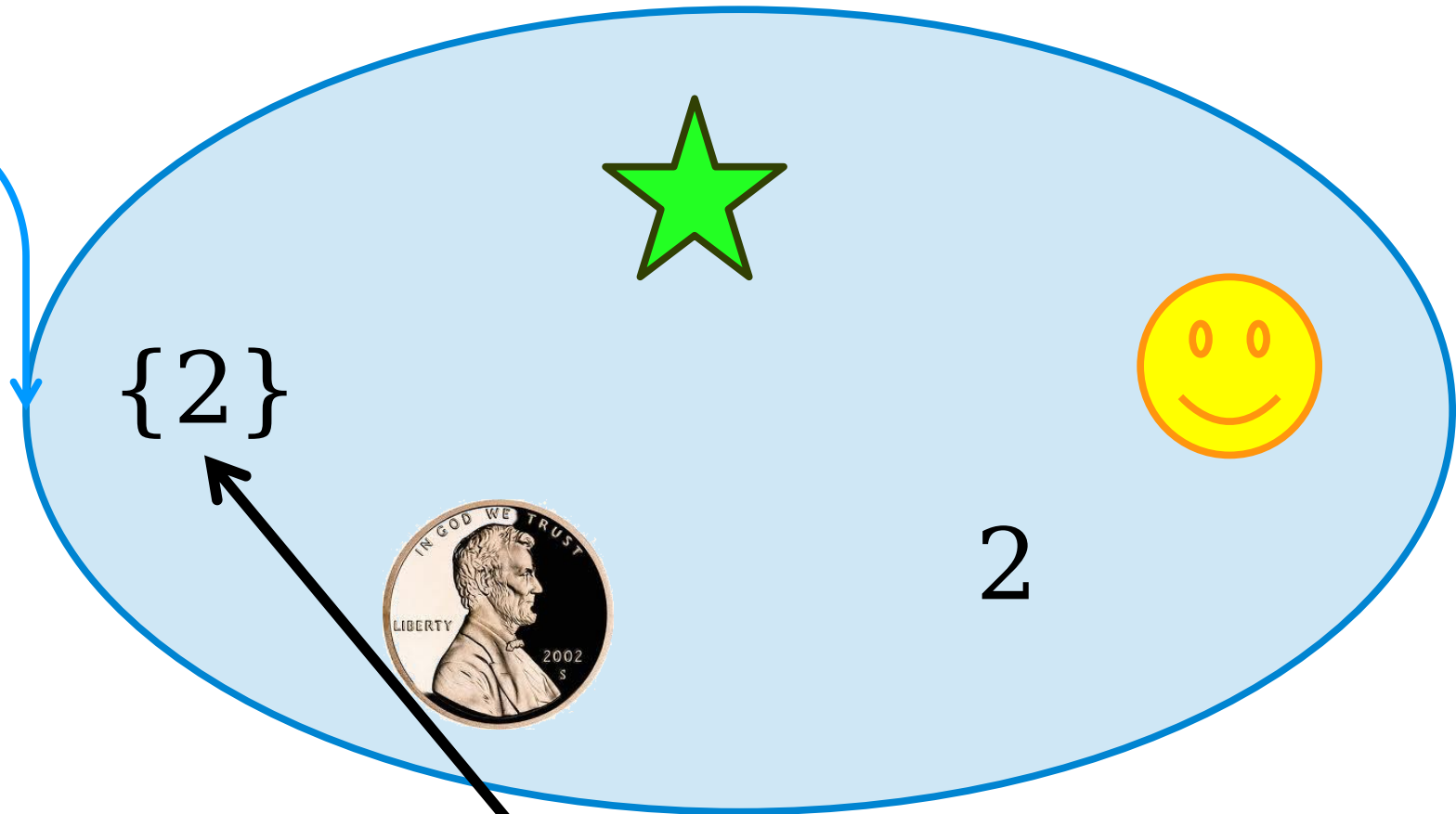
Set S



$$2 \in S$$

Subsets and Elements

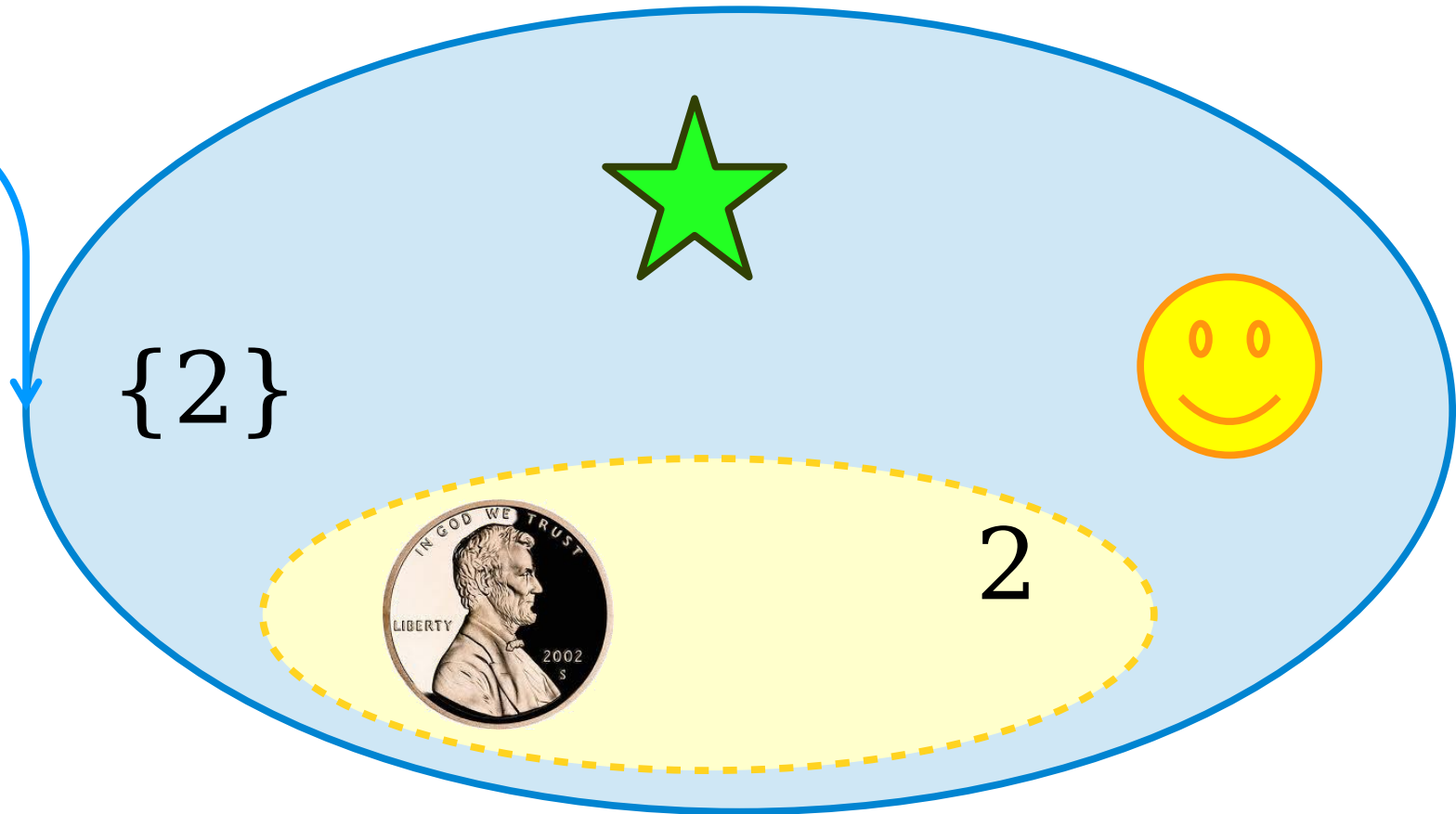
Set S



$$\{2\} \in S$$

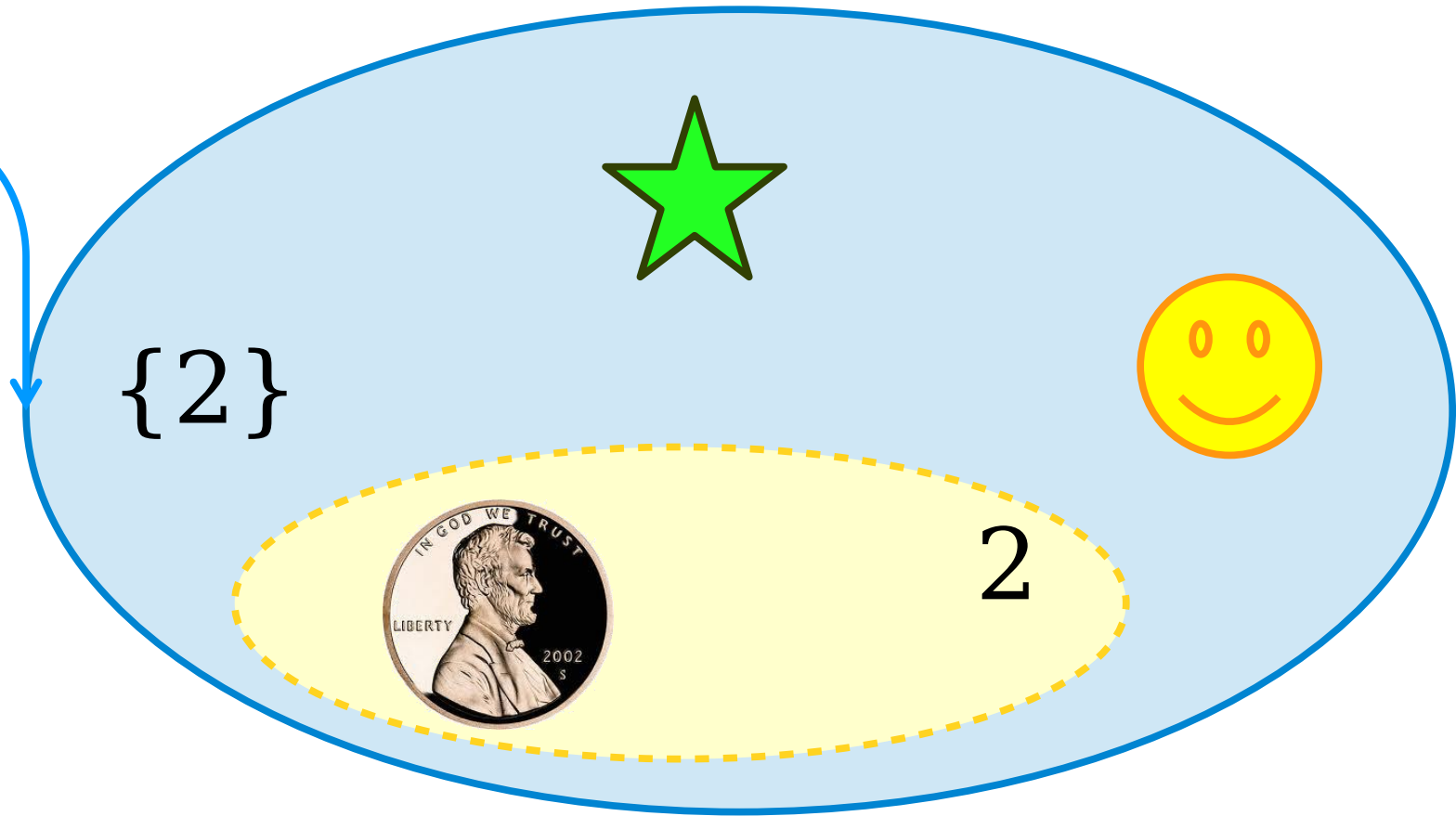
Subsets and Elements

Set S



Subsets and Elements

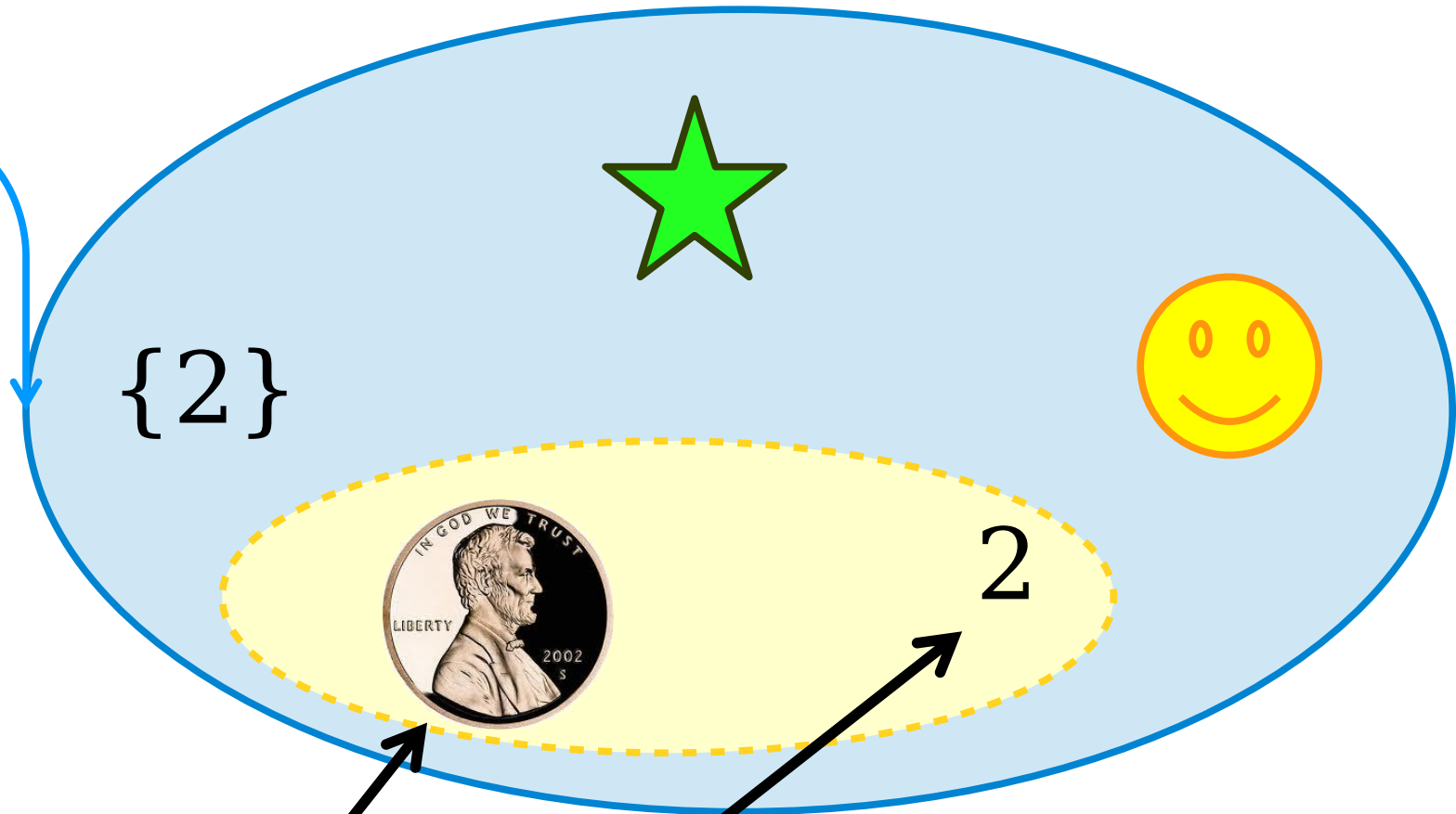
Set S



$$\left\{ \text{penny}, 2 \right\} \subseteq S$$

Subsets and Elements

Set S



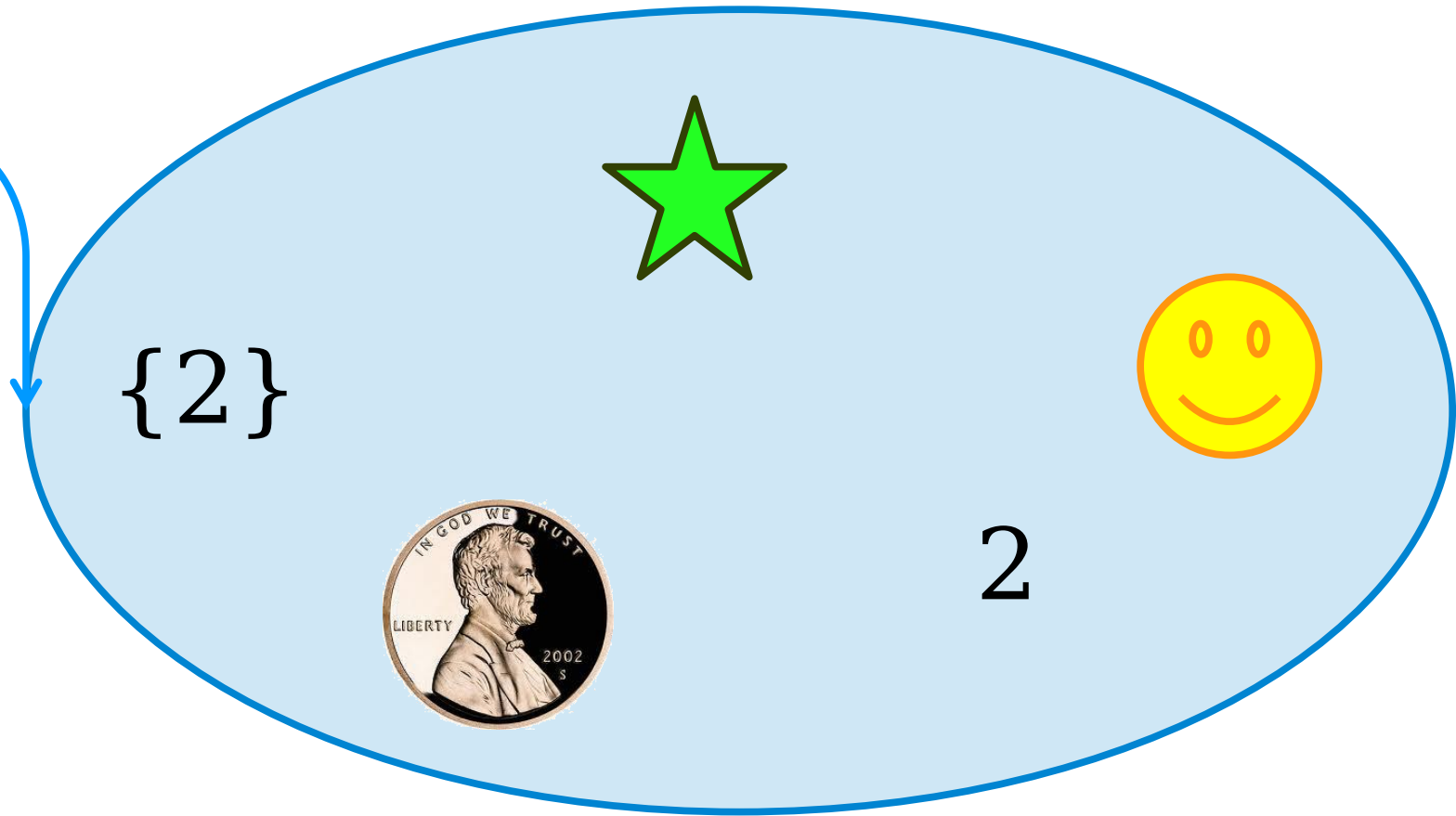
$\{2\}$

2

$$\{ \text{coin}, 2 \} \subseteq S$$

Subsets and Elements

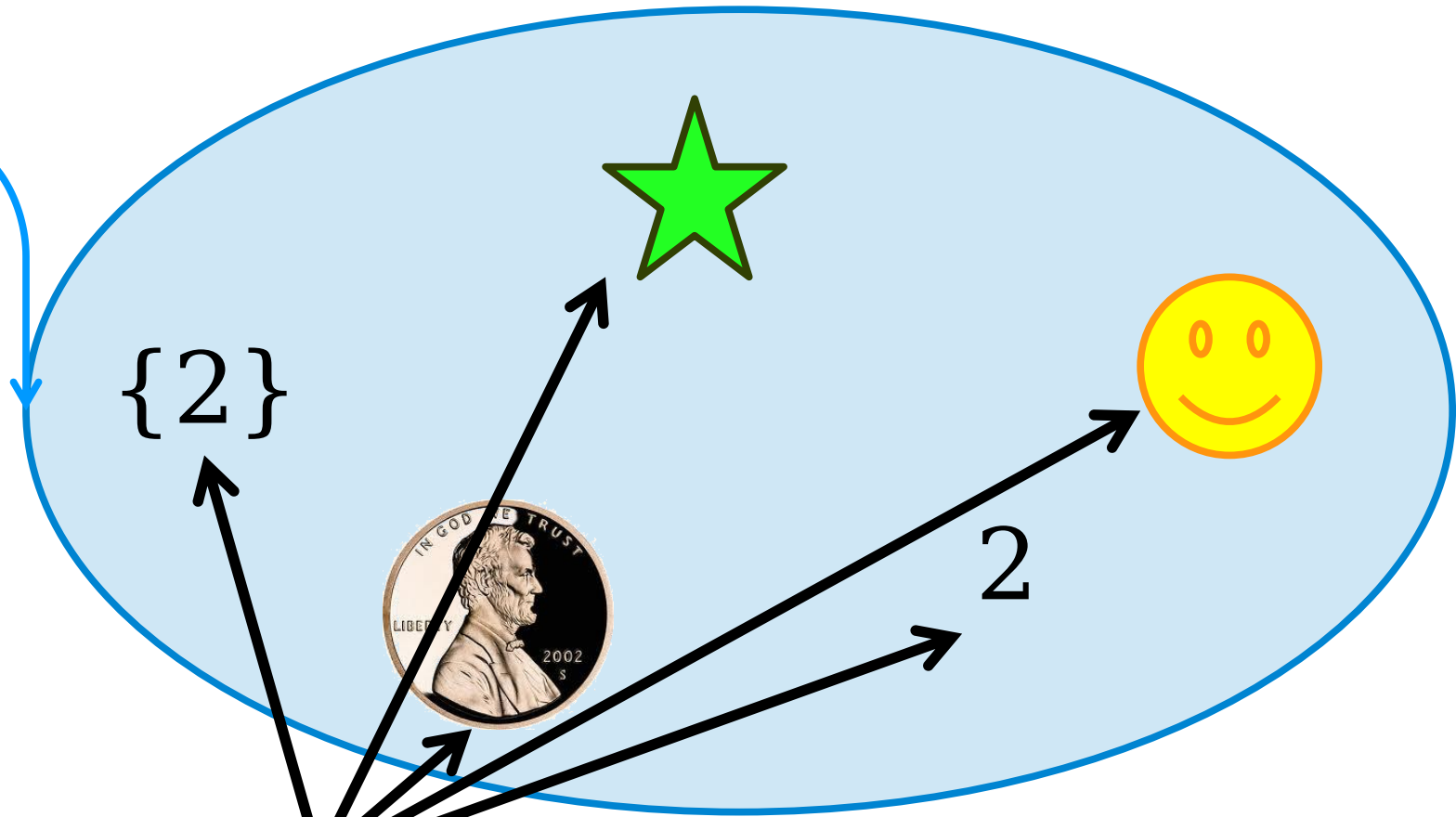
Set S



$$\left\{ \text{penny}, 2 \right\} \notin S$$

Subsets and Elements

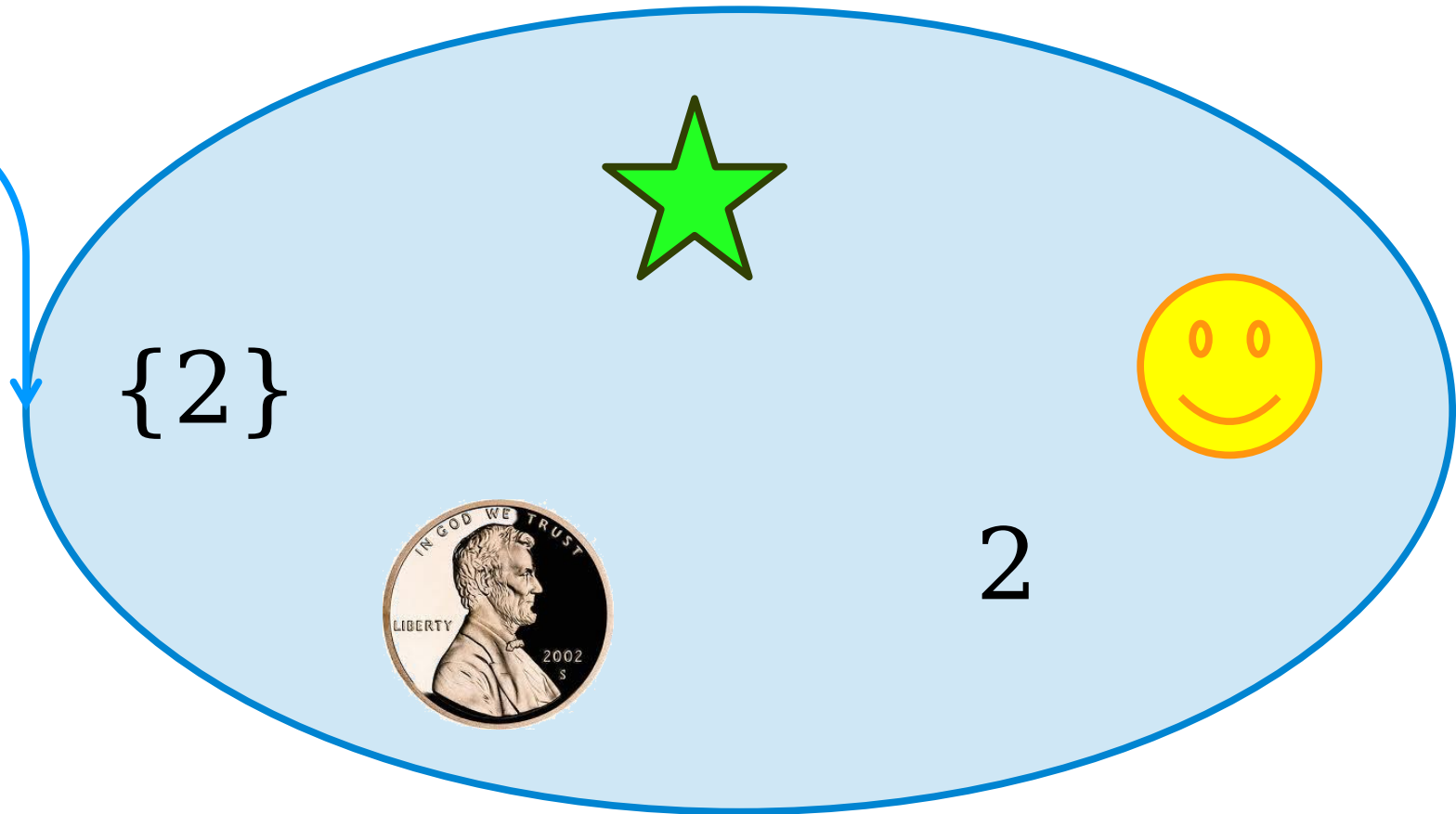
Set S



$\{ \text{penny}, 2 \} \notin S$

Subsets and Elements

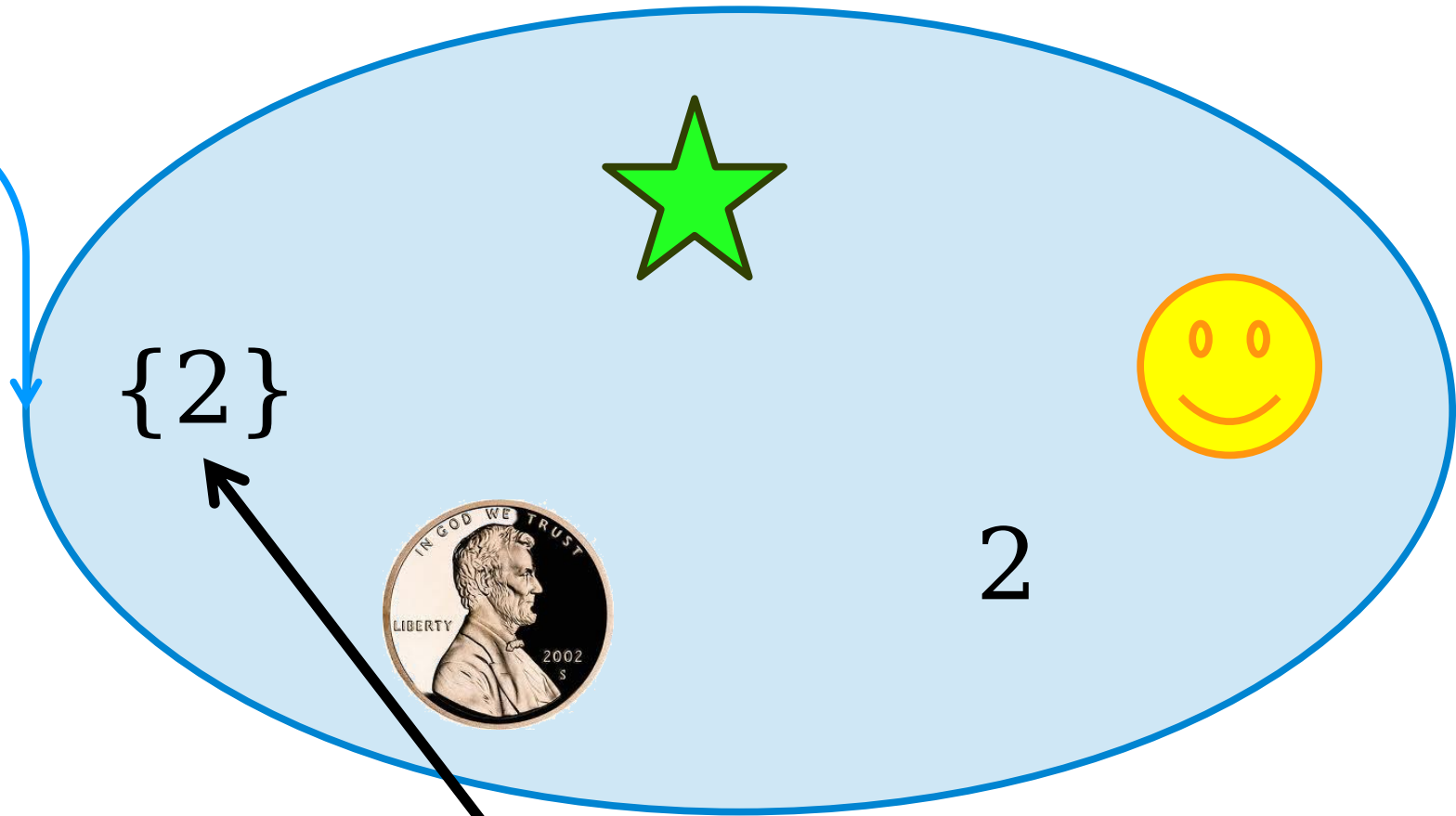
Set S



$$\{2\} \in S$$

Subsets and Elements

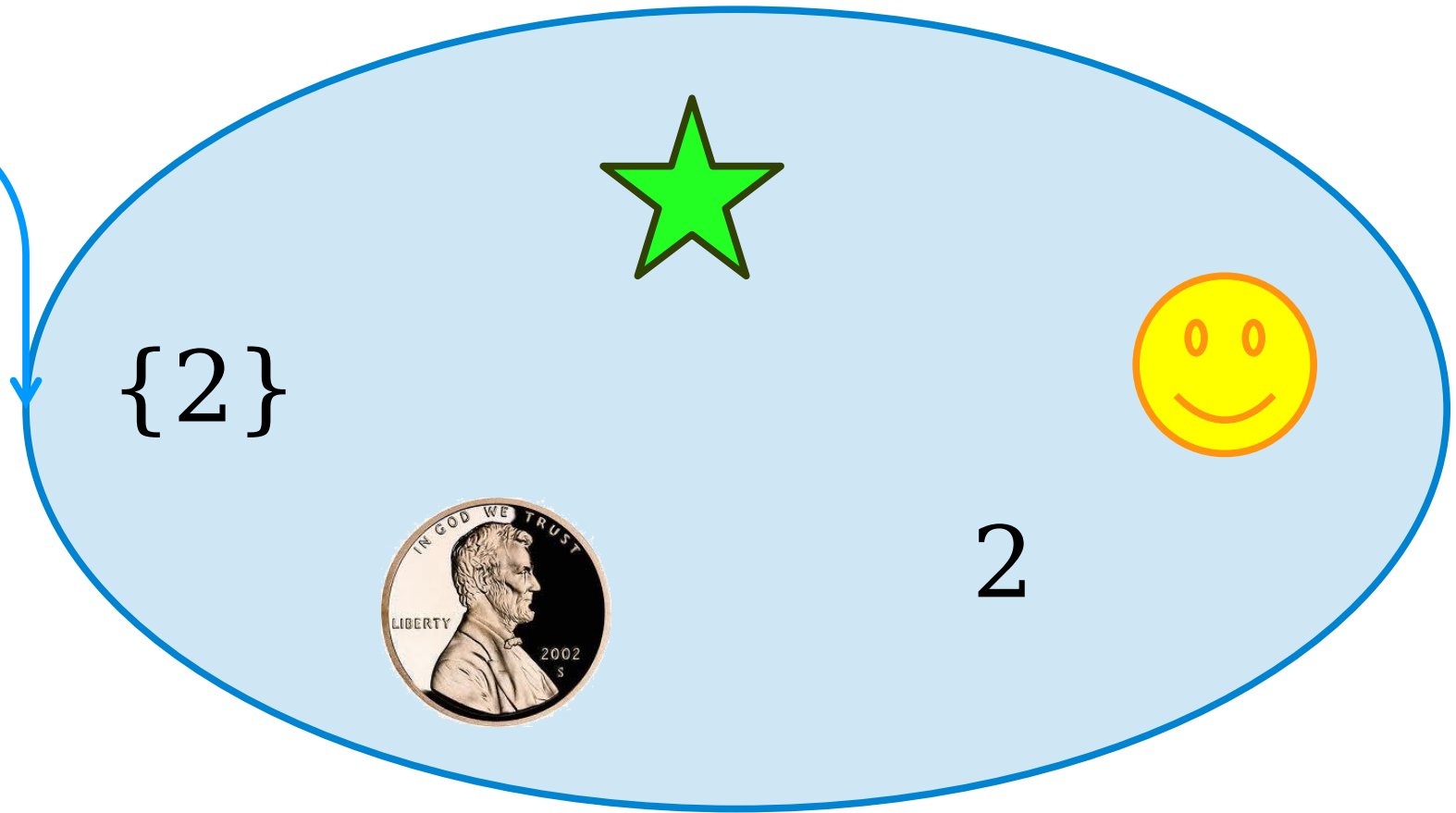
Set S



$$\{2\} \in S$$

Subsets and Elements

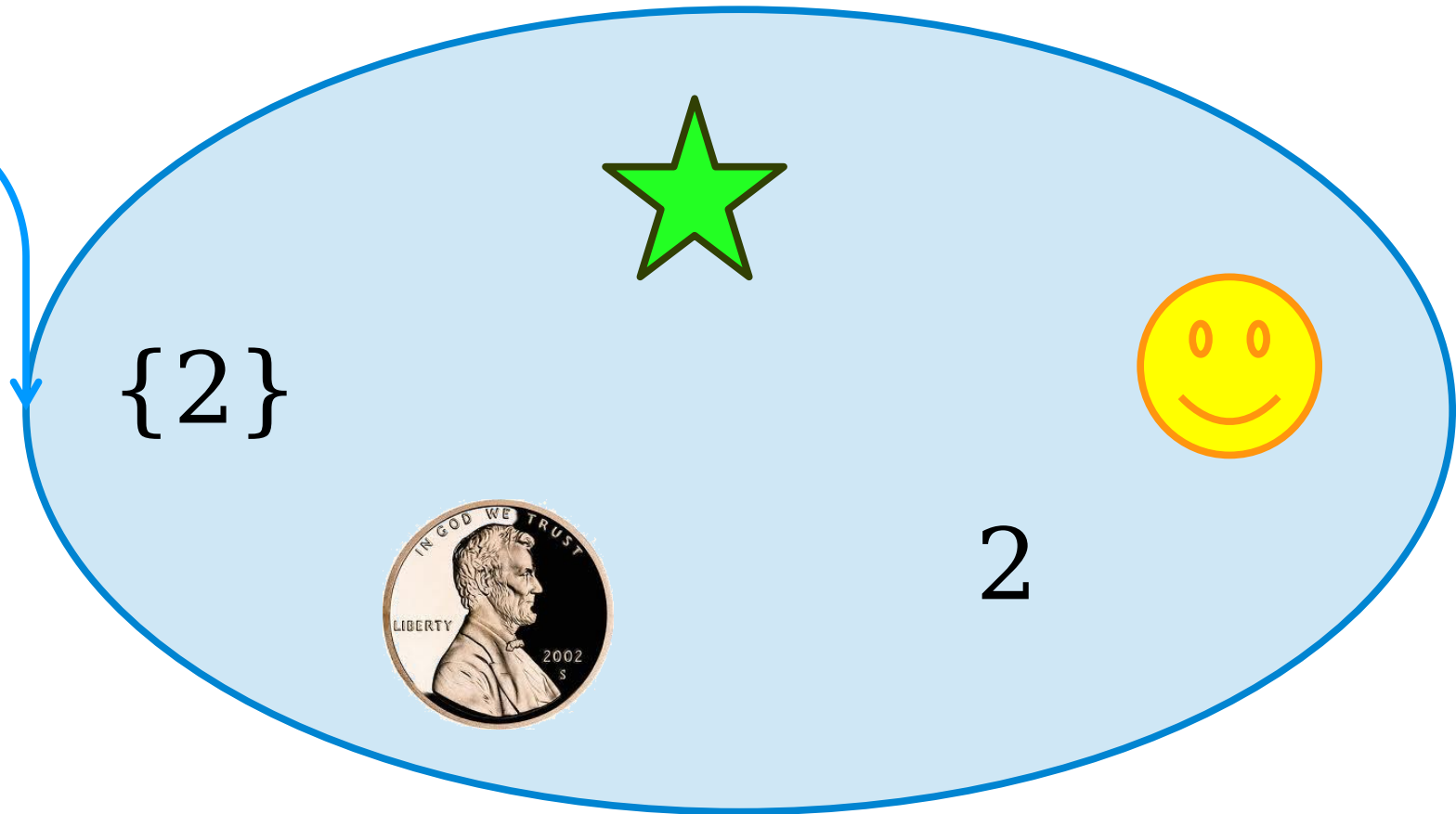
Set S



$$\{2\} \subseteq S$$

Subsets and Elements

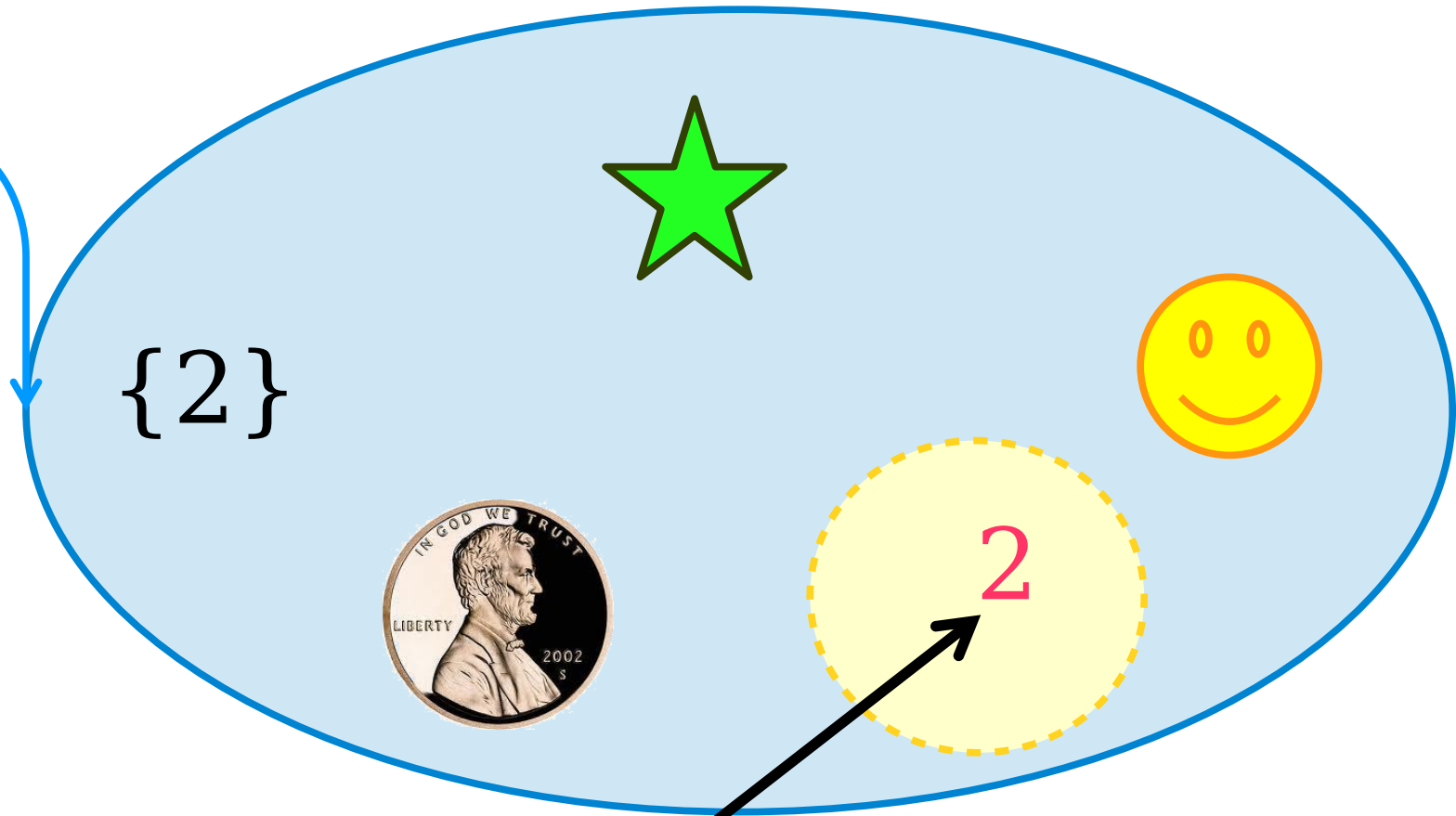
Set S



$$\{2\} \subseteq S$$

Subsets and Elements

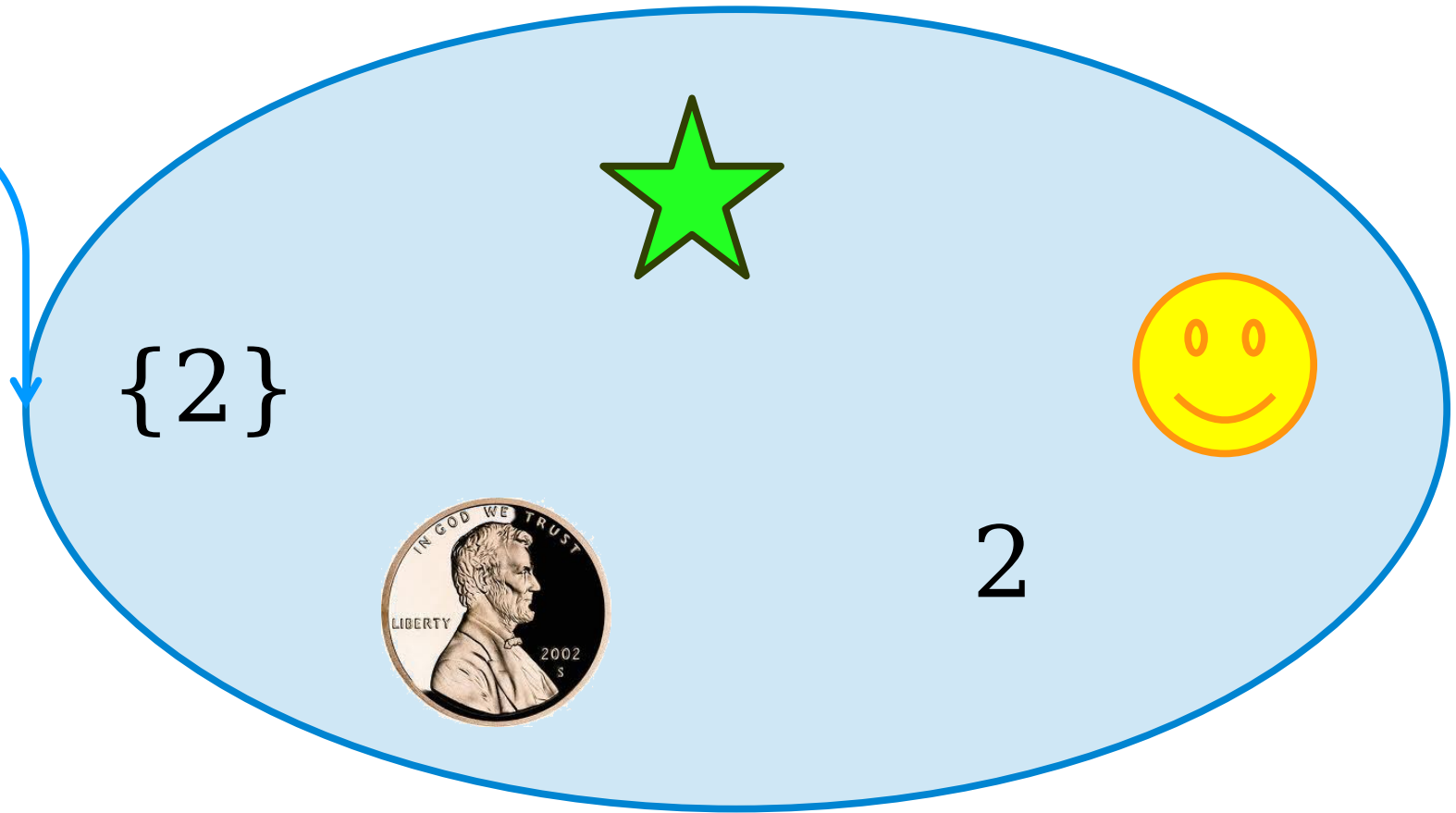
Set S



$$\{2\} \subseteq S$$

Subsets and Elements

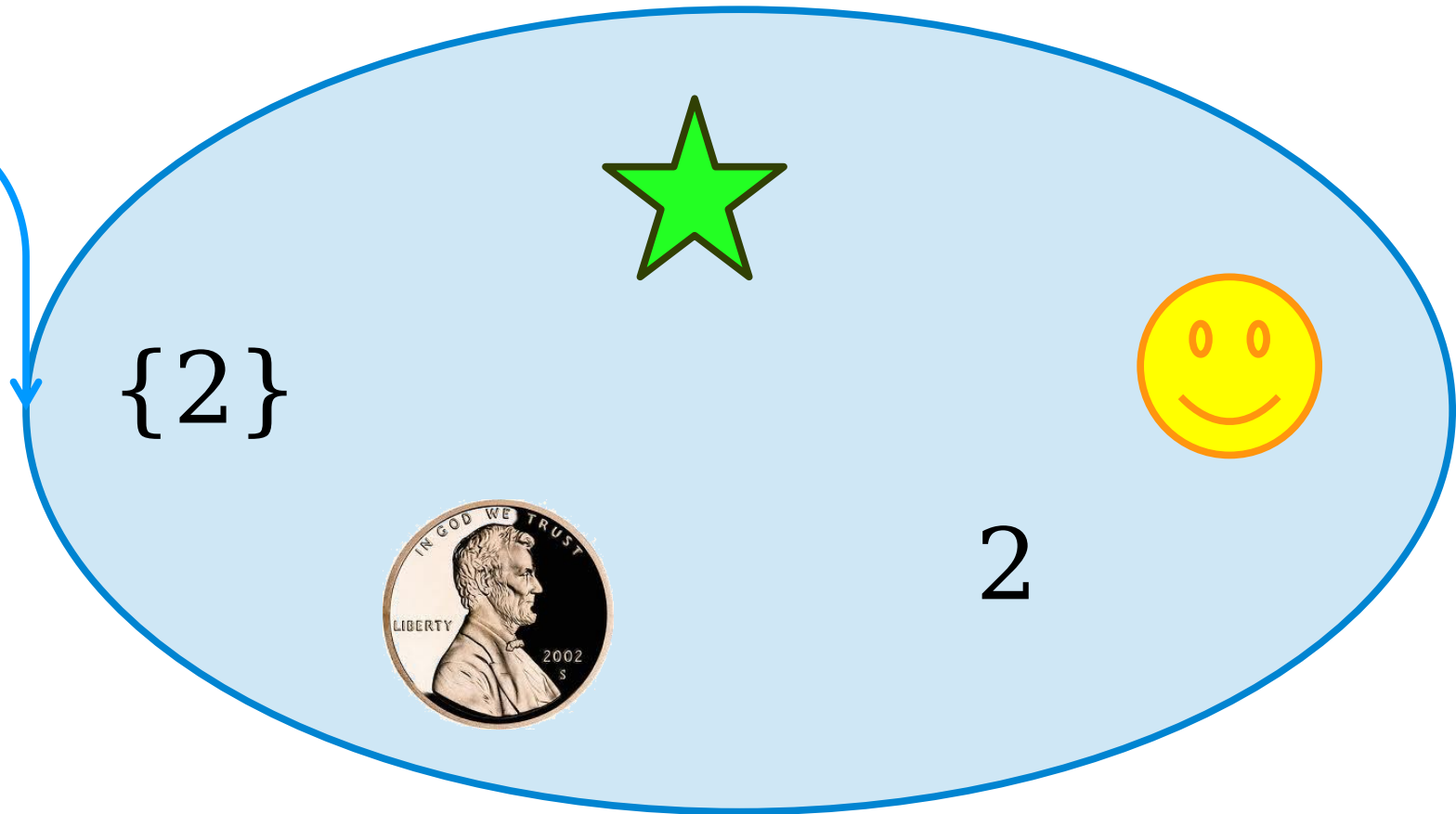
Set S



$2 \notin S$

Subsets and Elements

Set S



(Since 2 isn't a set.)

$$2 \notin S$$

Subsets and Elements

We say that $S \in T$ if, among the elements of T , one of them is *exactly* the object S .

We say that $S \subseteq T$ if S is a set and every element of S is also an element of T . (S has to be a set for the statement $S \subseteq T$ to be true.)

Although these concepts are similar, ***they are not the same!*** Not all elements of a set are subsets of that set and vice-versa.

We have a resource on the course website, the Guide to Elements and Subsets, that explores this in more depth.

What About the Empty Set?

A set S is called a **subset** of a set T (denoted $S \subseteq T$) if all elements of S are also elements of T .

Are there any sets T where $\emptyset \subseteq T$?

Equivalently, is there a set T where the following statement is true?

“All elements of \emptyset are also elements of T ”

Yes! In fact, this statement is true for *every* set T !

Vacuous Truth

A statement of the form

**“All objects of type P
are also of type Q ”**

is called ***vacuously true*** if there are no objects of type P .

Vacuously true statements are true *by definition*.
This is a convention used throughout mathematics.

Some examples:

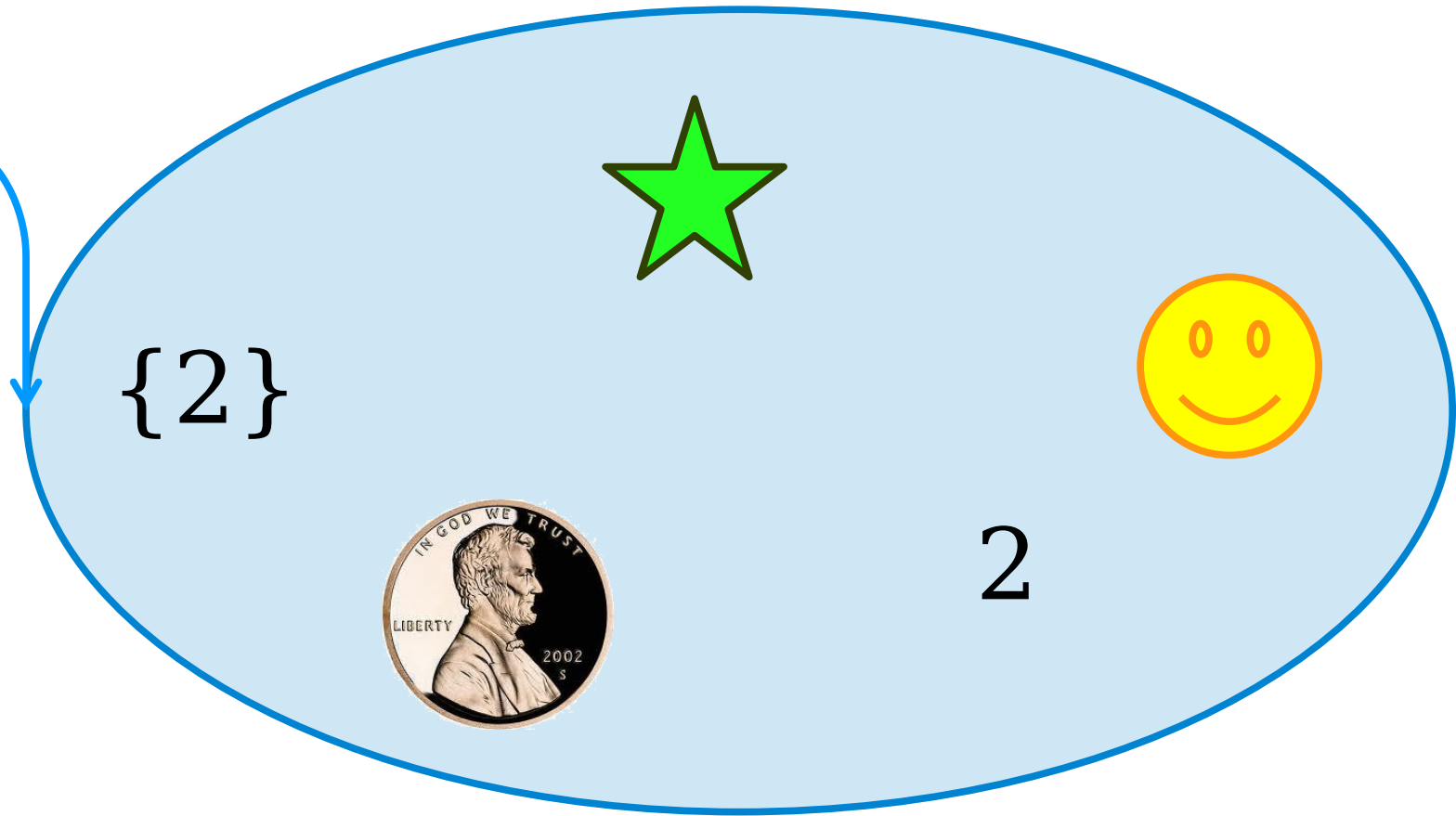
All unicorns are pink.

All unicorns are blue.

Every element of \emptyset is also an element of T .

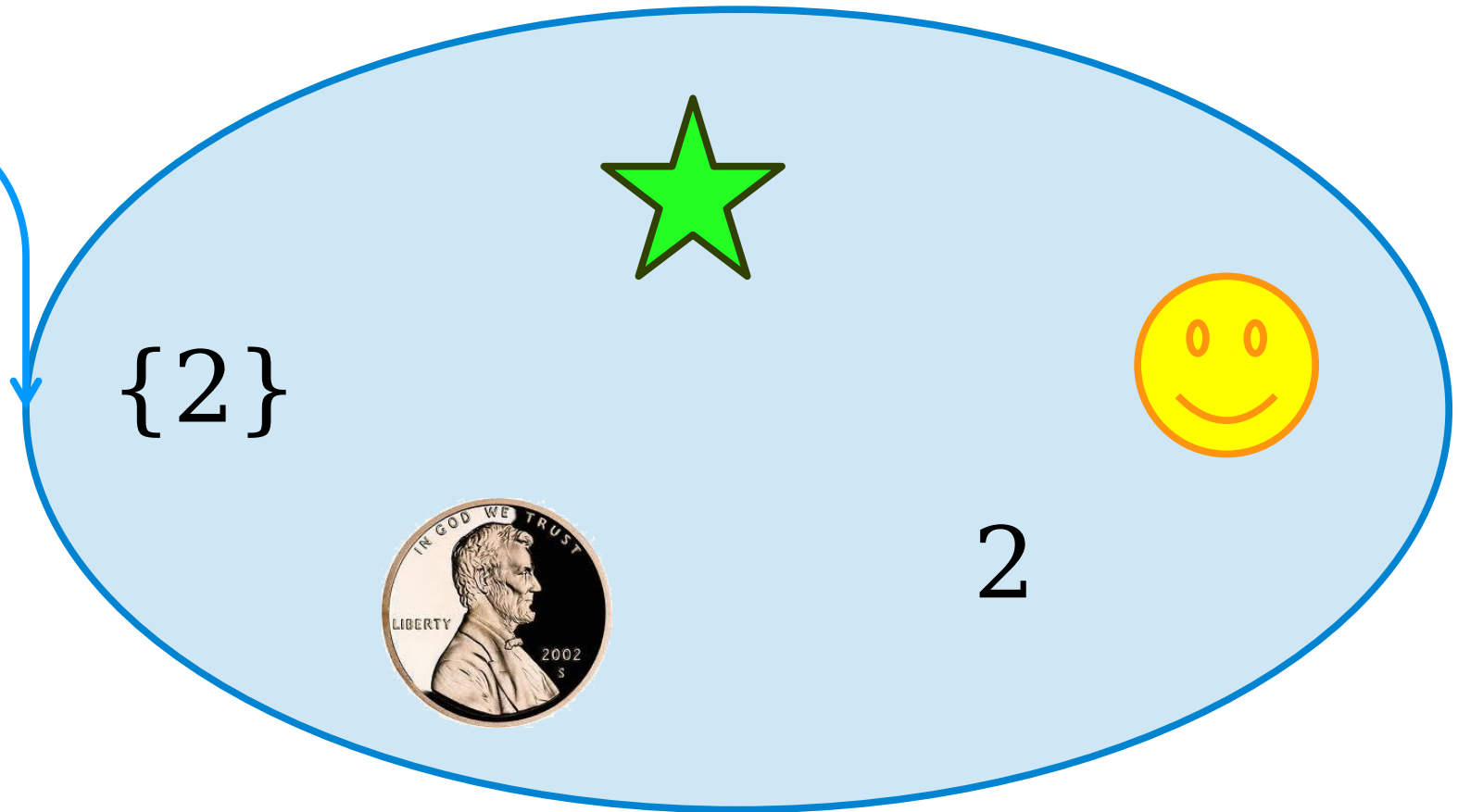
Subsets and Elements

Set S



Subsets and Elements

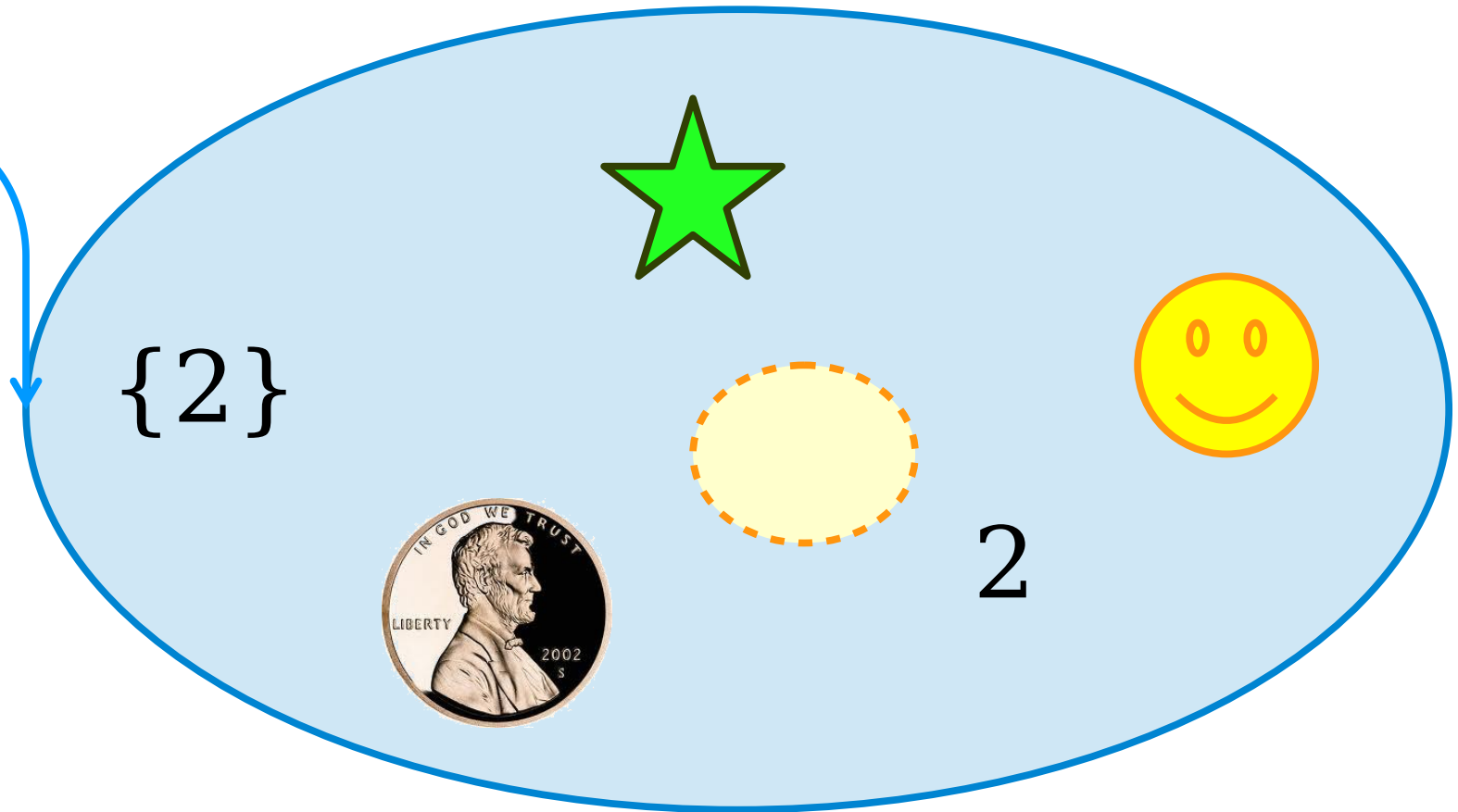
Set S



$$\emptyset \subseteq S$$

Subsets and Elements

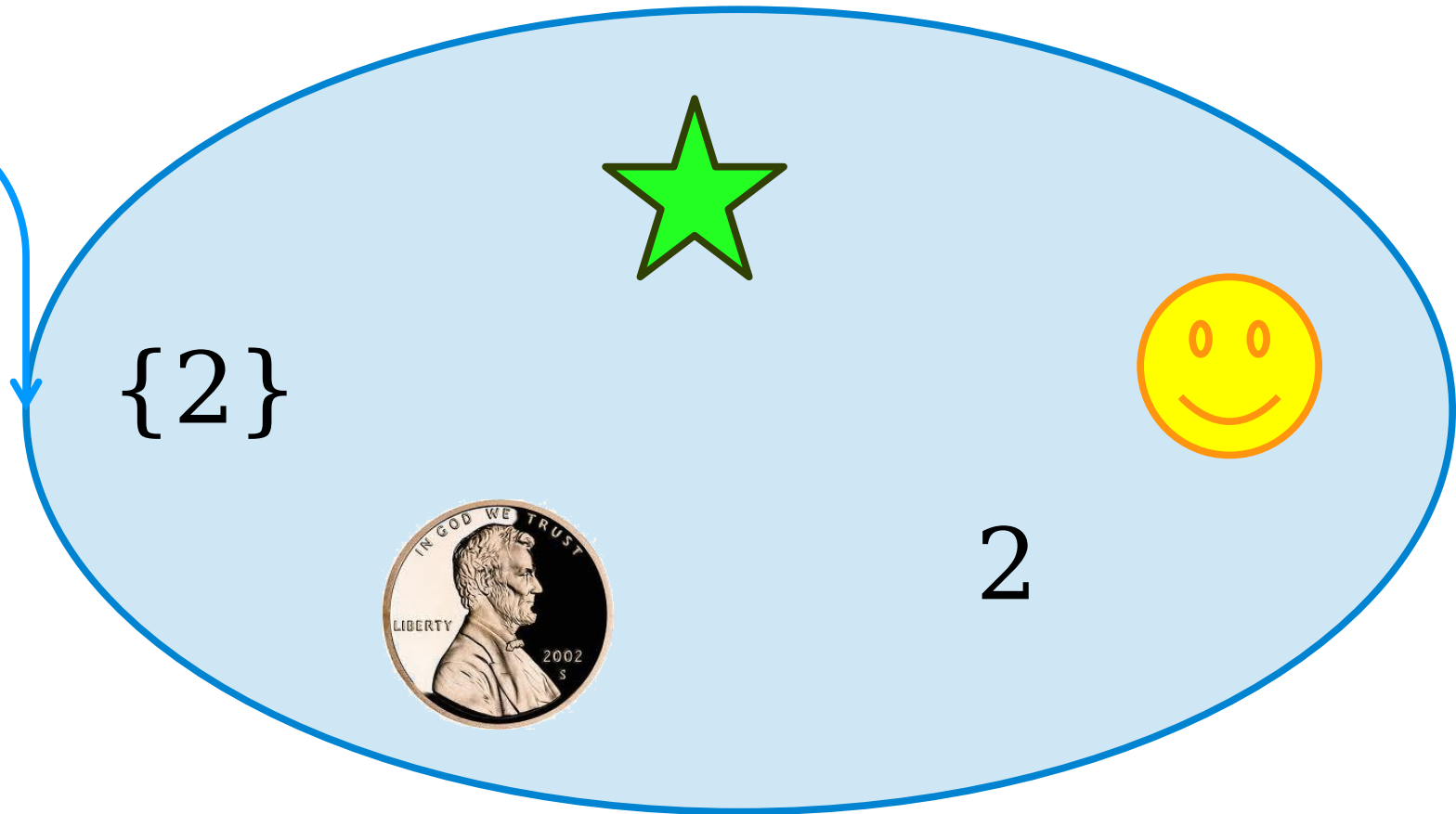
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Subsets and Elements

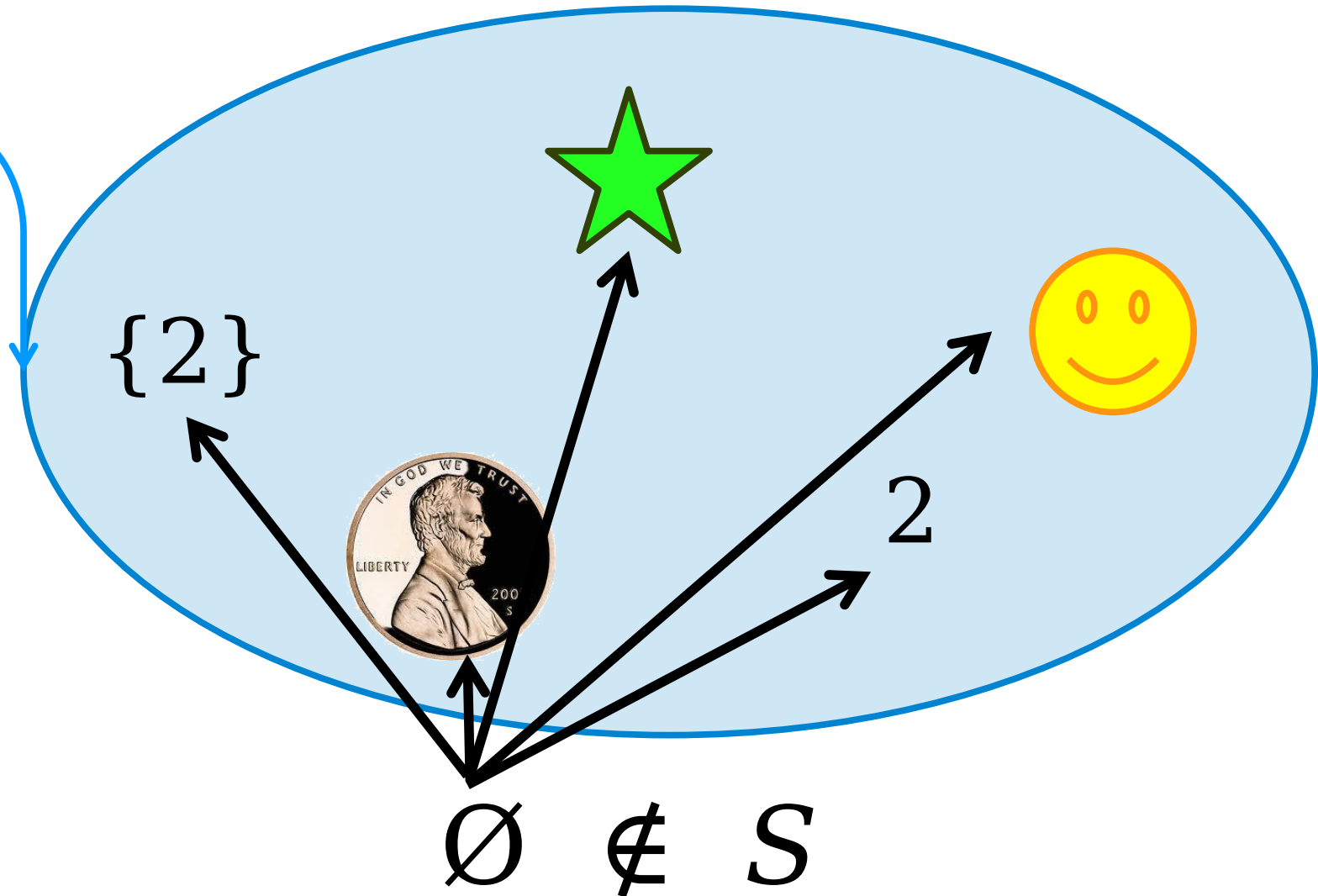
Set S



$\emptyset \notin S$

Subsets and Elements

Set S



$$S = \left\{ \text{Lincoln Penny}, \text{Lincoln Dime} \right\}$$

$$\wp(S) = \left\{ \emptyset, \left\{ \text{Lincoln Dime} \right\}, \left\{ \text{Lincoln Penny} \right\}, \left\{ \text{Lincoln Penny}, \text{Lincoln Dime} \right\} \right\}$$

This is the **power set** of S , the set of all subsets of S . We write the power set of S as $\wp(S)$.

Formally, $\wp(S) = \{ T \mid T \subseteq S \}$.

(Do you see why?)

What is $\wp(\emptyset)$?

Answer: $\{\emptyset\}$

Remember that $\emptyset \neq \{\emptyset\}$!

Cardinality

Cardinality

The *cardinality* of a set is the number of elements it contains.

If S is a set, we denote its cardinality by writing $|S|$.

Examples:

- $|\{38, 31\}| = 2$
- $|\{\{a, b\}, \{c, d, e, f, g\}, \{h\}\}| = 3$
- $|\{1, 2, 3, 3, 3, 3, 3\}| = 3$
- $|\{n \in \mathbb{N} \mid n < 137\}| = 137$

The Cardinality of \mathbb{N}

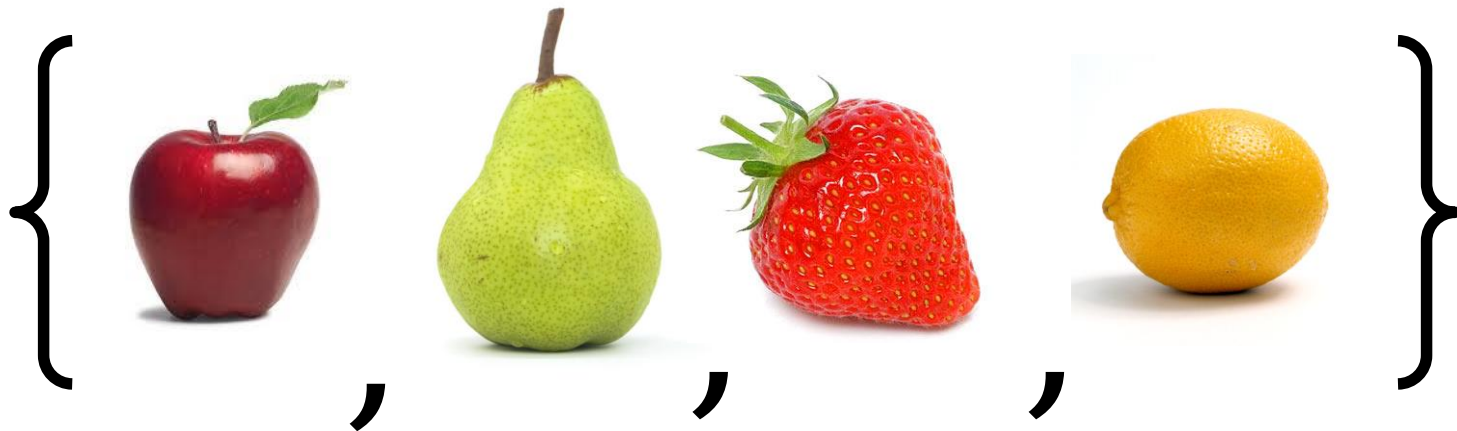
- What is $|\mathbb{N}|$?
- There are infinitely many natural numbers.
- $|\mathbb{N}|$ can't be a natural number, since it's infinitely large.
- We need to introduce a new term.
- Let's define $\aleph_0 = |\mathbb{N}|$.
- \aleph_0 is pronounced “aleph-zero,” “aleph-nought,” or “aleph-null.”

Consider the set

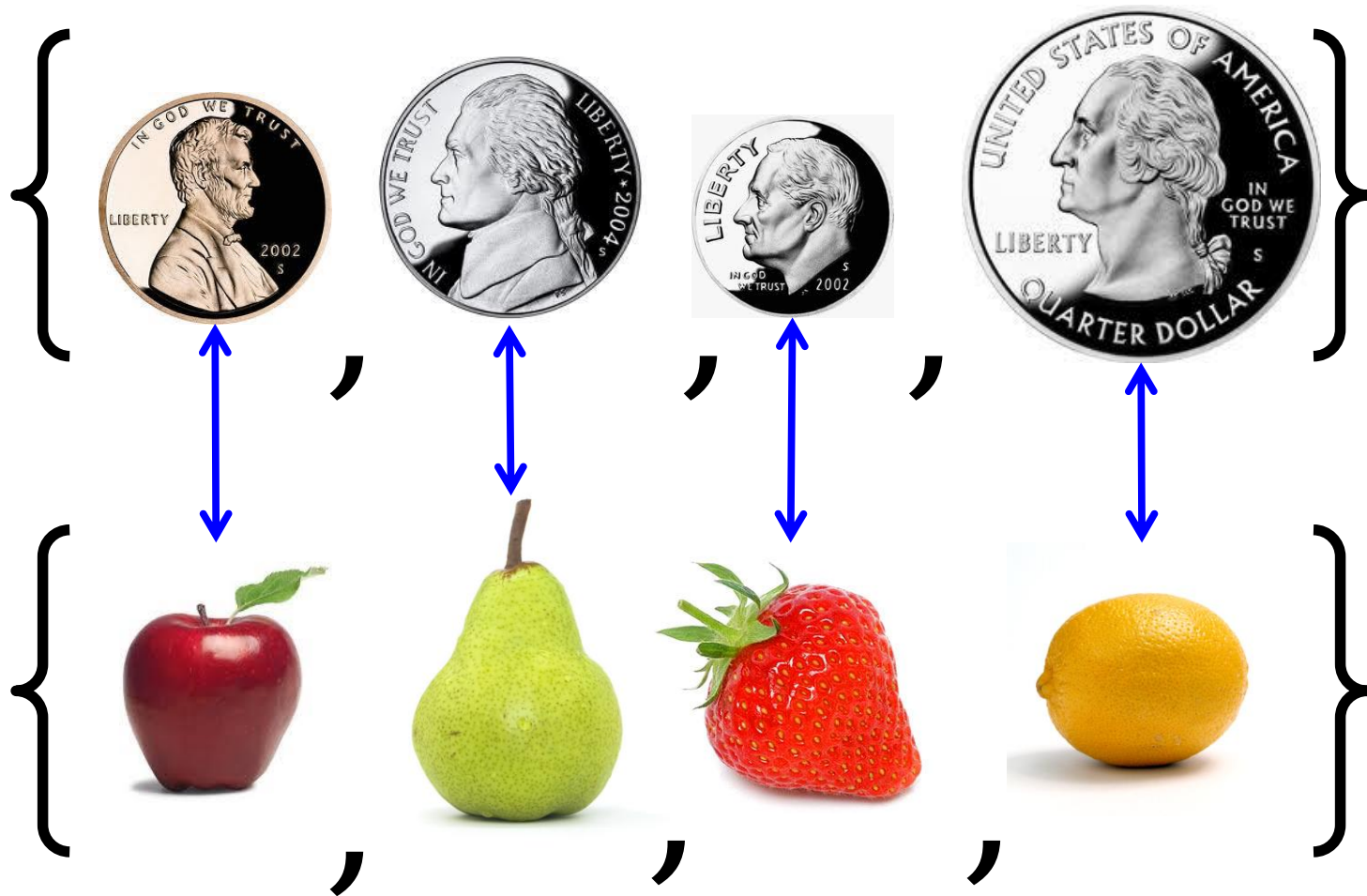
$$S = \{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

What is $|S|$?

How Big Are These Sets?



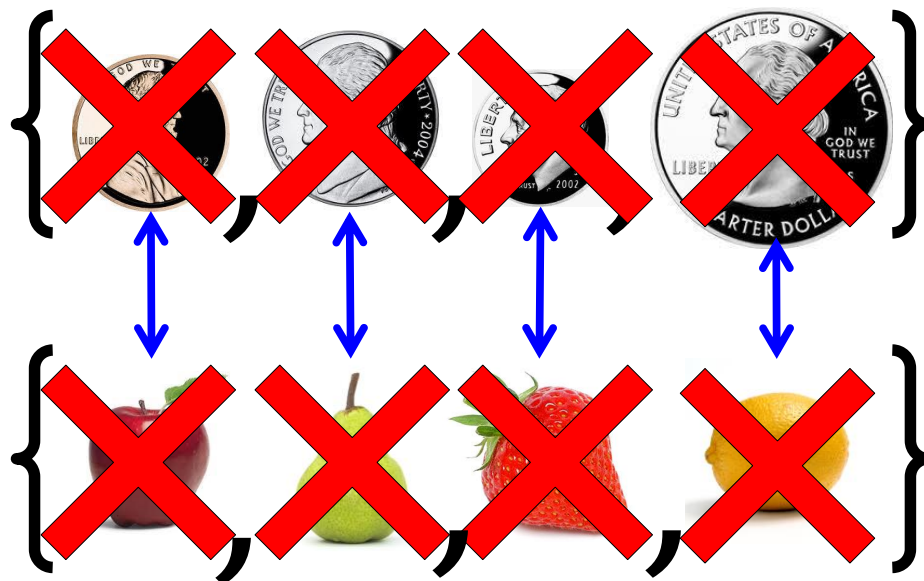
How Big Are These Sets?



Comparing Cardinalities

By definition, two sets have the same size if there is a way to pair their elements off without leaving any elements uncovered.

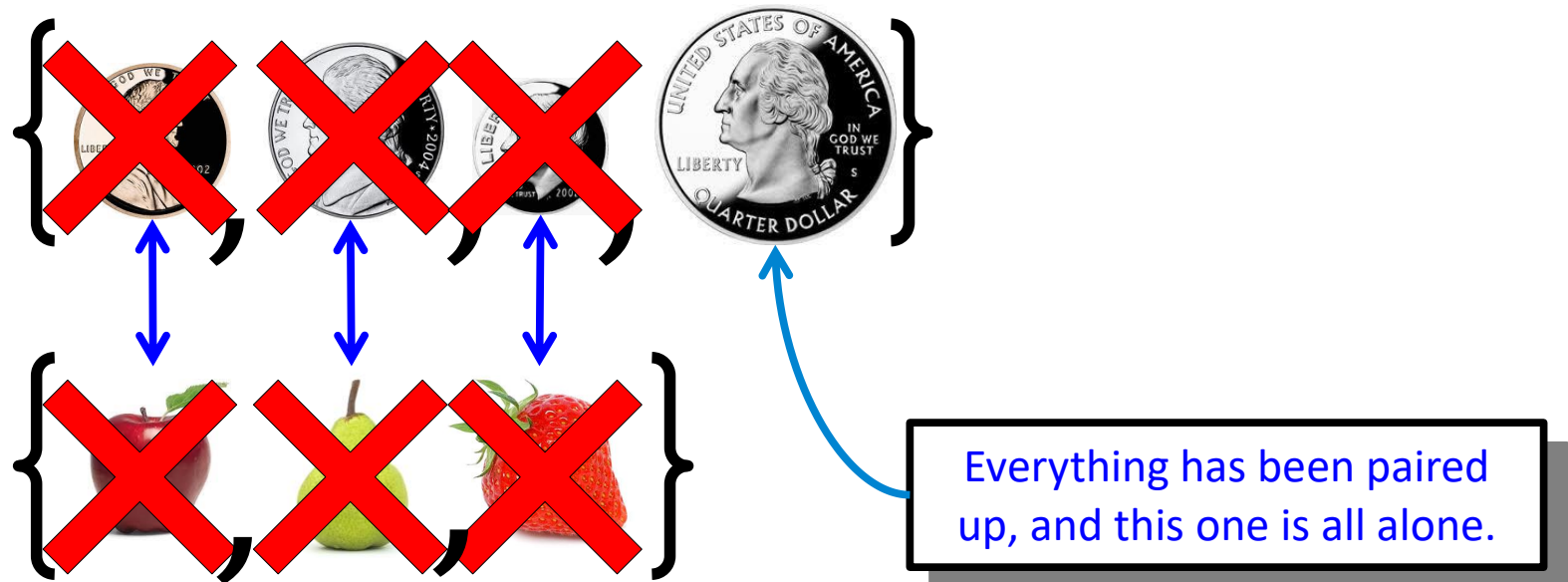
The intuition:



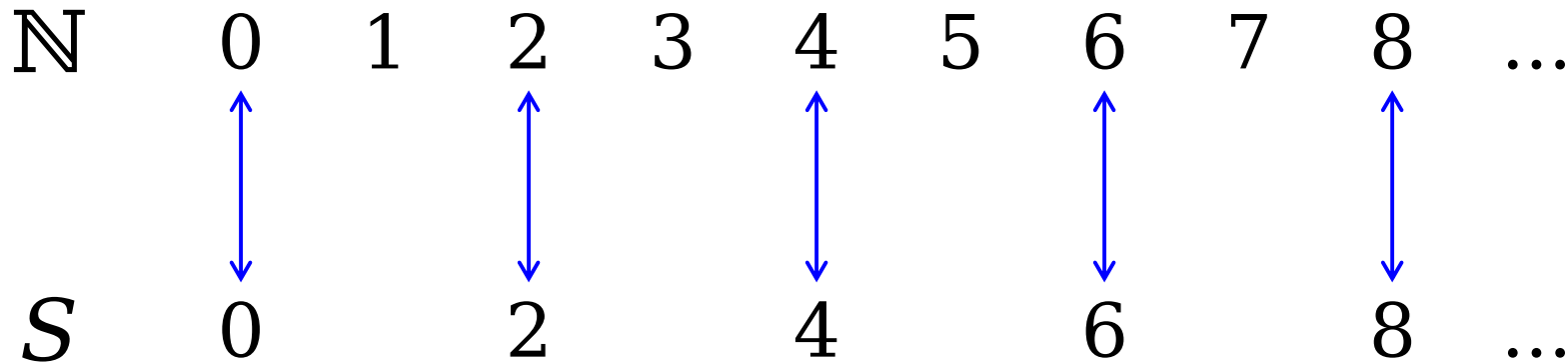
Comparing Cardinalities

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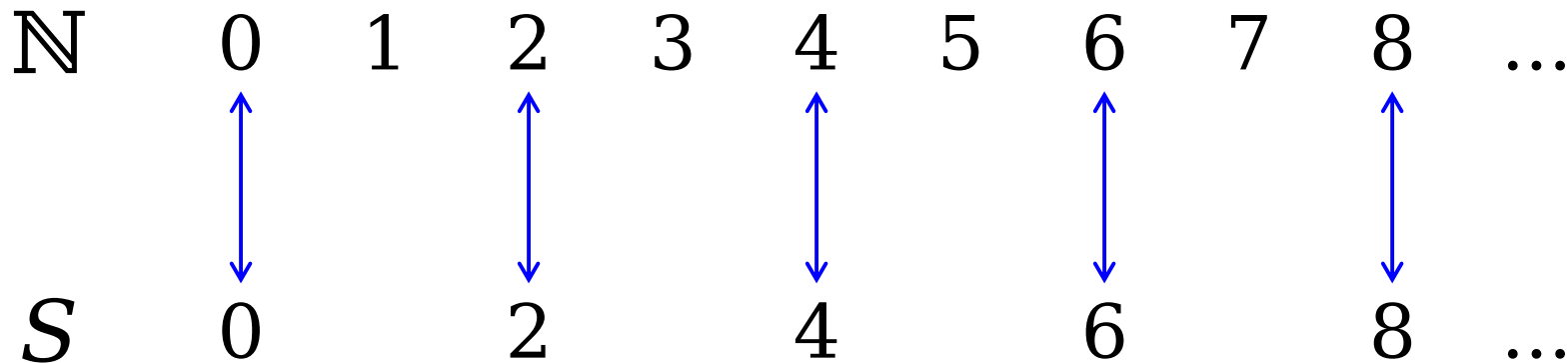
Infinite Cardinalities



$$S = \{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

Two sets have the same size if there is a way to pair their elements off without leaving any elements uncovered

Infinite Cardinalities



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Two sets have the same size if *there is a way* to pair their elements off without leaving any elements uncovered

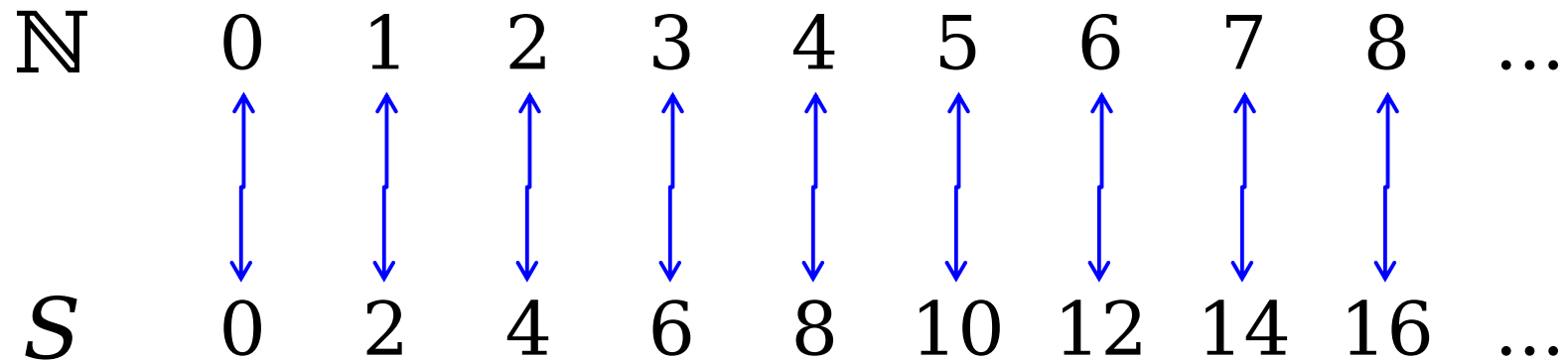
Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

S 0 2 4 6 8 ...

$$S = \{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

Infinite Cardinalities



$$n \leftrightarrow 2n$$

$$S = \{ n \mid n \in \mathbb{N} \text{ and } n \text{ is even} \}$$

$$|S| = |\mathbb{N}| = \aleph_0$$

Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

\mathbb{Z} ... -3 -2 -1 0 1 2 3 4 ...

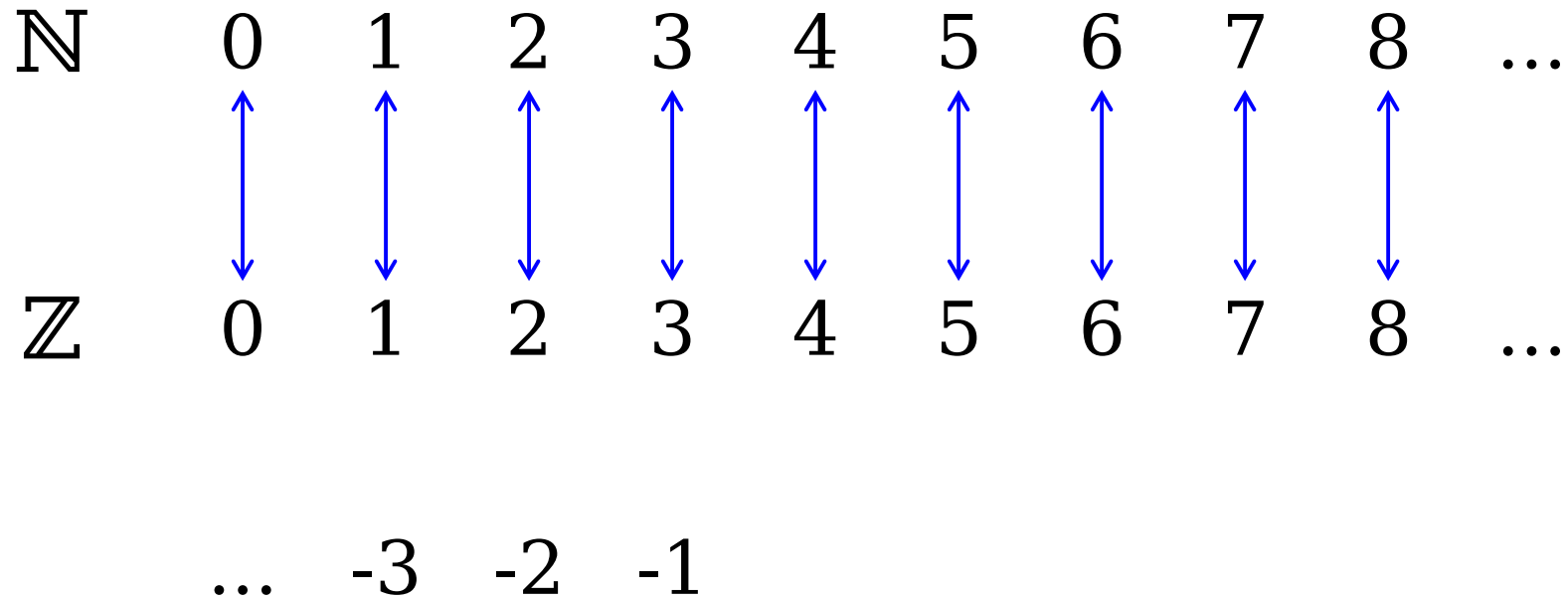
Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

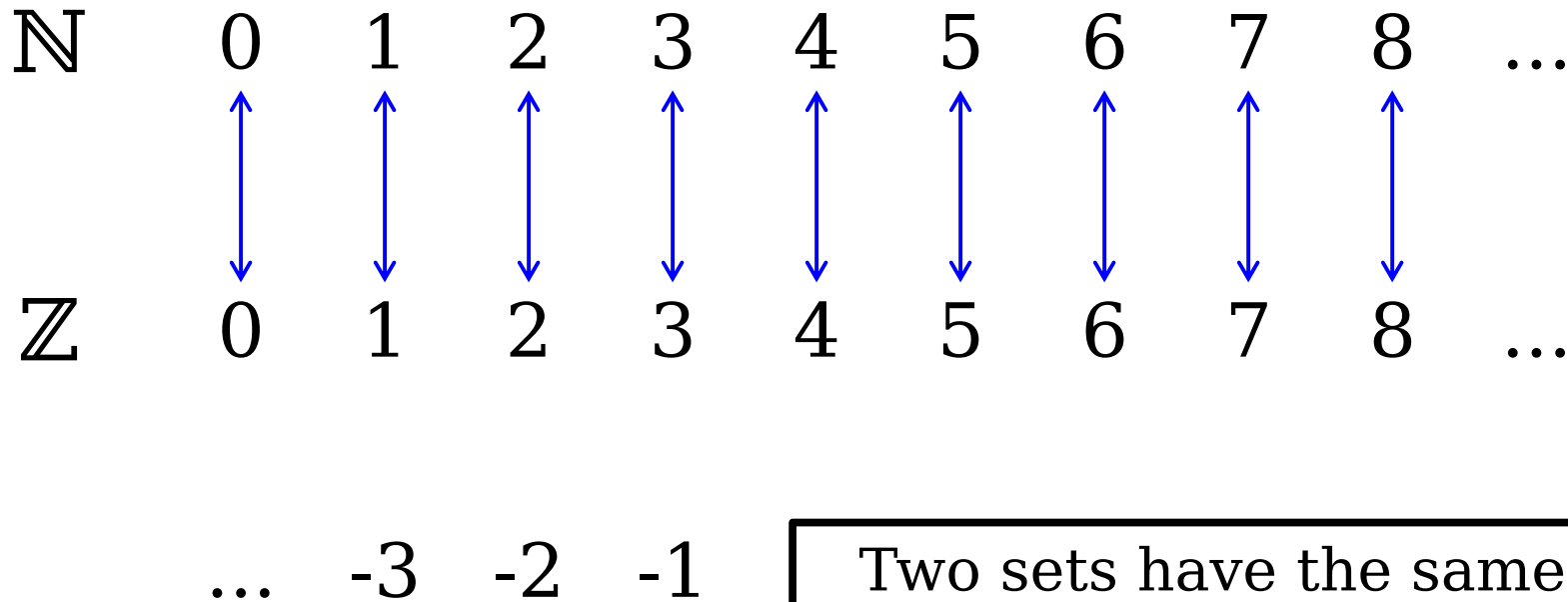
\mathbb{Z} 0 1 2 3 4 ...

... -3 -2 -1

Infinite Cardinalities



Infinite Cardinalities



Two sets have the same size if *there is a way* to pair their elements off without leaving any elements uncovered

Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

\mathbb{Z} ... -3 -2 -1 0 1 2 3 4 ...

Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

\mathbb{Z}

... -3 -2 -1 0 1 2 3 4 ...

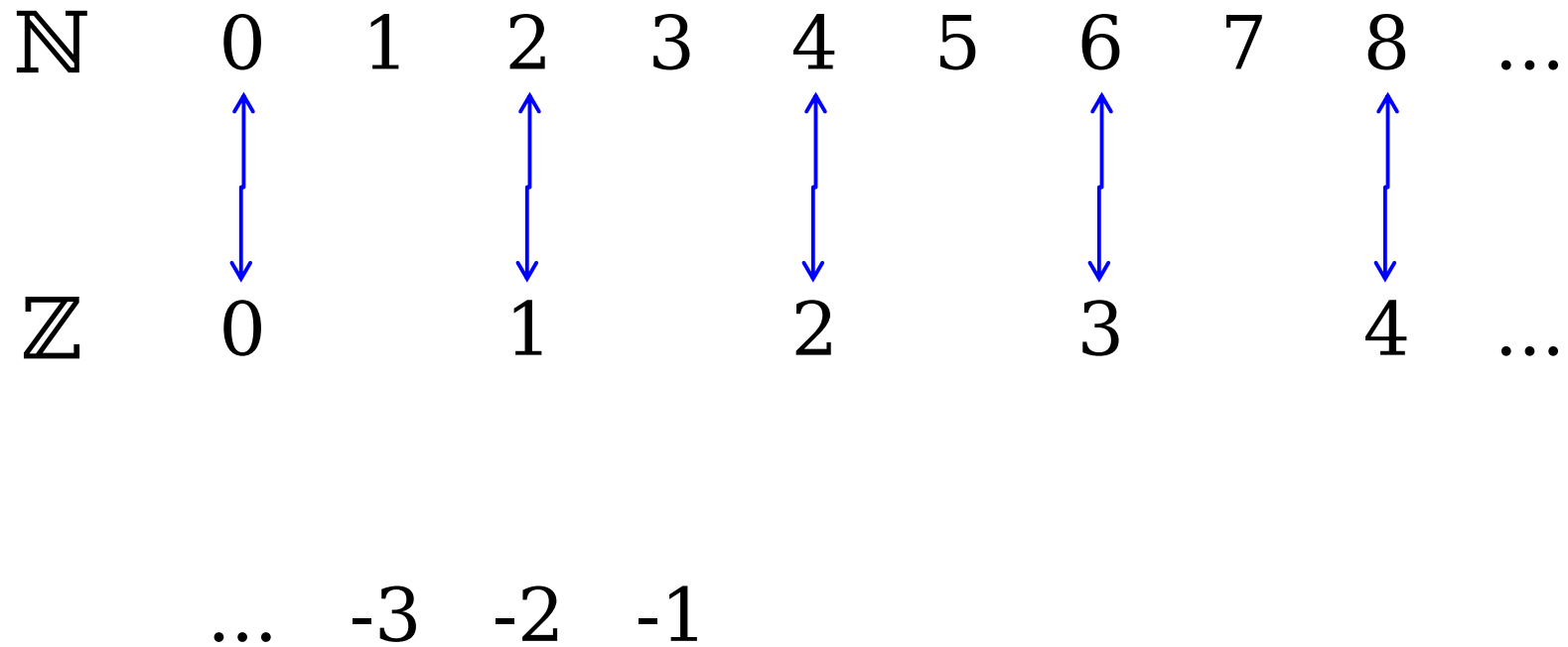
Infinite Cardinalities

\mathbb{N} 0 1 2 3 4 5 6 7 8 ...

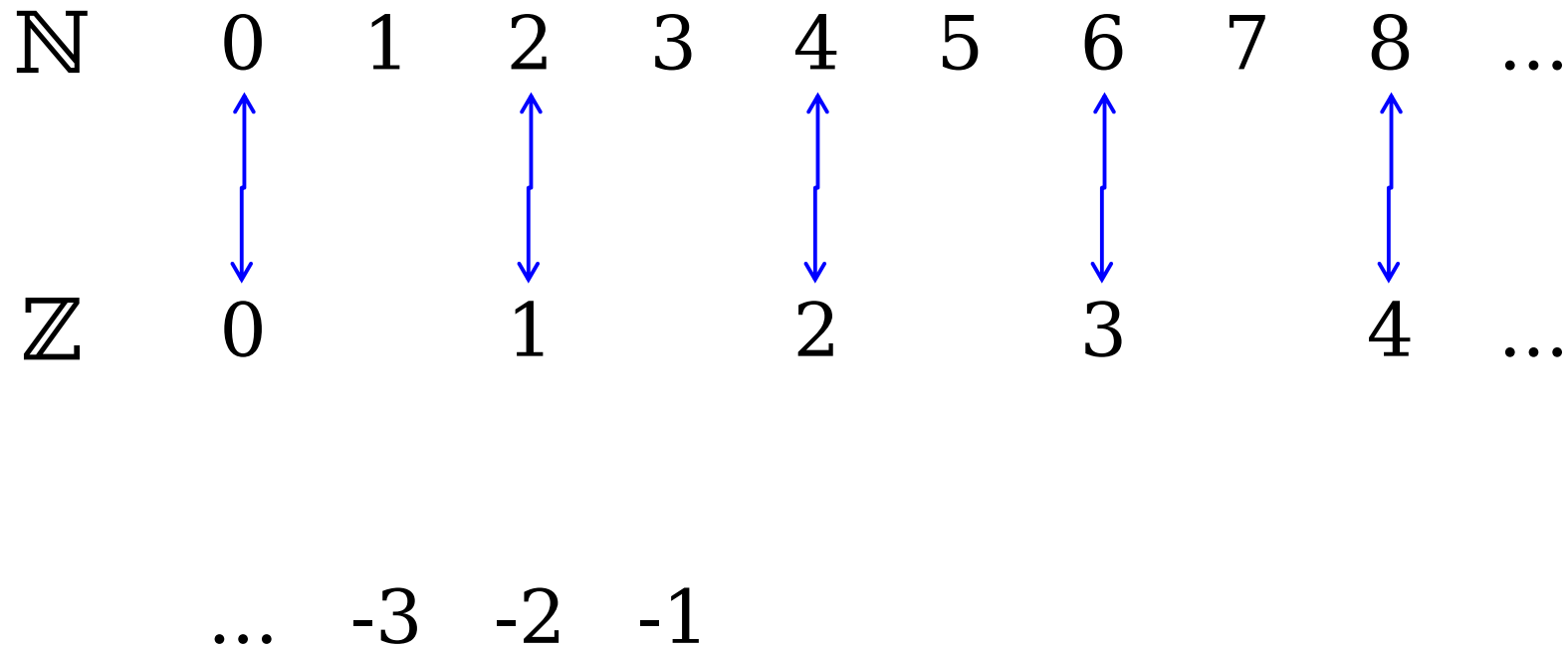
\mathbb{Z} 0 1 2 3 4 ...

... -3 -2 -1

Infinite Cardinalities

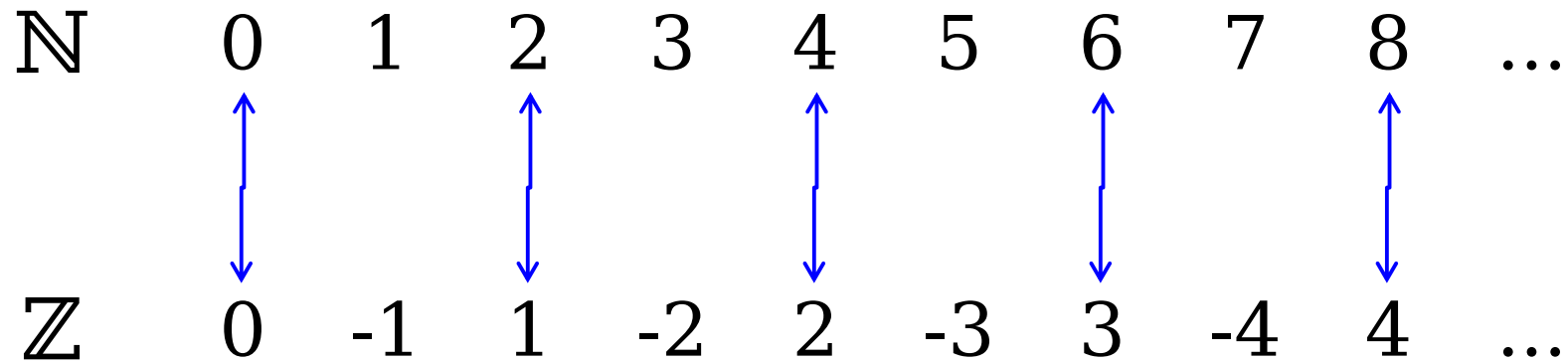


Infinite Cardinalities



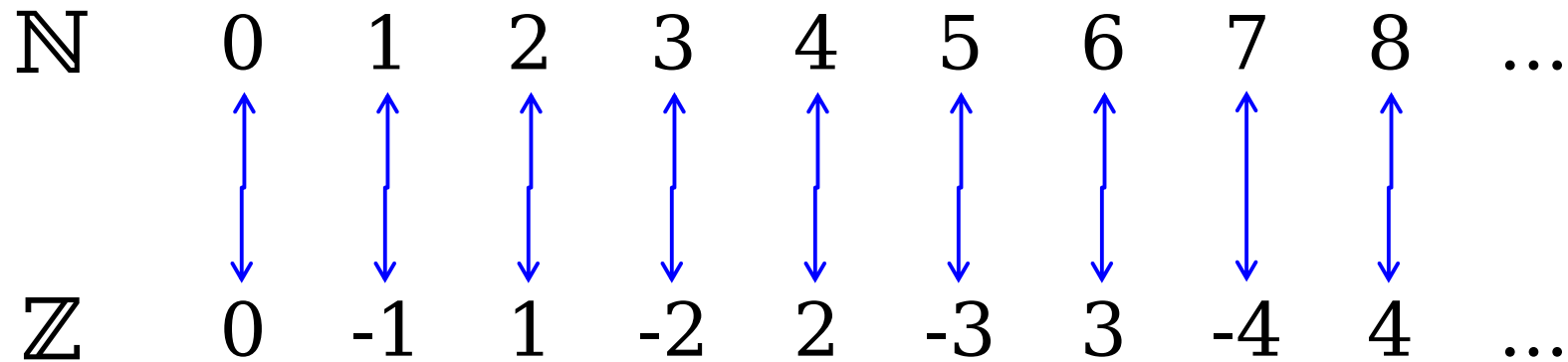
Pair nonnegative integers with even natural numbers.

Infinite Cardinalities



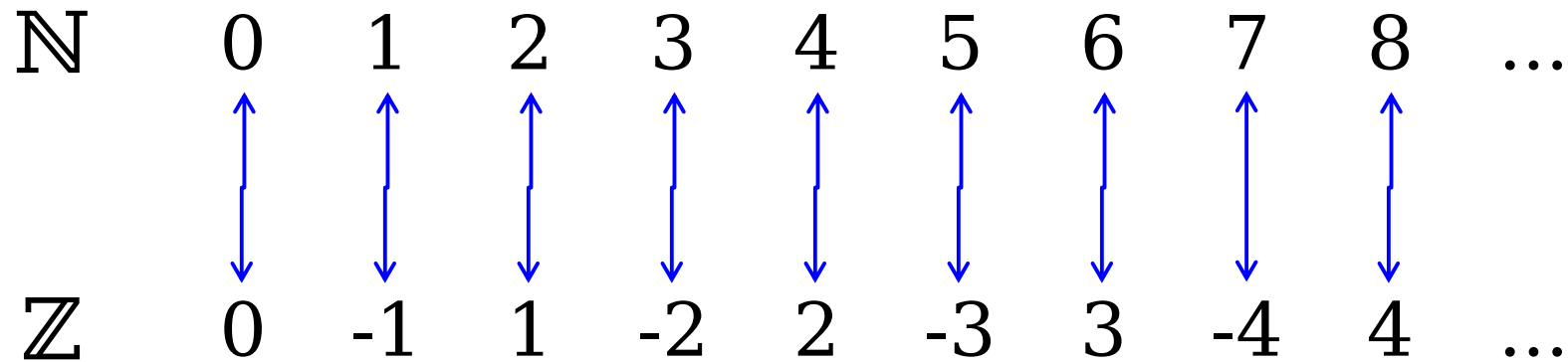
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Infinite Cardinalities



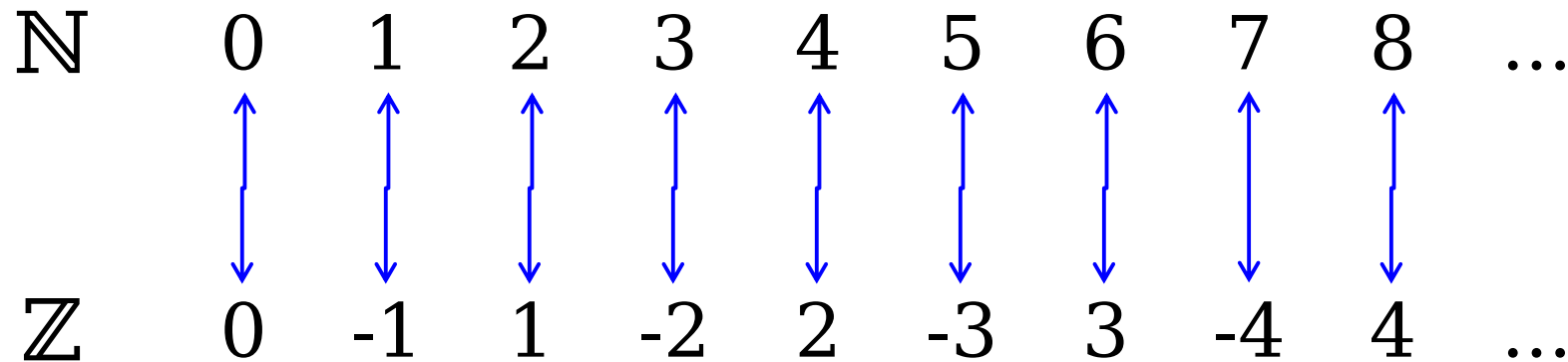
Pair nonnegative integers with even natural numbers.

Infinite Cardinalities



Pair nonnegative integers with even natural numbers.
Pair negative integers with odd natural numbers.

Infinite Cardinalities



$$|\mathbb{N}| = |\mathbb{Z}| = \aleph_0$$

Pair nonnegative integers with even natural numbers.
Pair negative integers with odd natural numbers.

Important Question:

Do all infinite sets have
the same cardinality?

$$S = \left\{ \text{Lincoln Penny}, \text{Lincoln Nickel} \right\}$$

$$\wp(S) = \left\{ \emptyset, \left\{ \text{Lincoln Nickel} \right\}, \left\{ \text{Lincoln Penny} \right\}, \left\{ \text{Lincoln Penny}, \text{Lincoln Nickel} \right\} \right\}$$

$$|S| < |\wp(S)|$$

$$S = \left\{ \text{Lincoln Penny}, \text{Lincoln Dime}, \text{Button} \right\}$$

$$\wp(S) = \left\{ \emptyset, \left\{ \text{Lincoln Penny} \right\}, \left\{ \text{Lincoln Dime} \right\}, \left\{ \text{Button} \right\}, \left\{ \text{Lincoln Penny}, \text{Lincoln Dime} \right\}, \left\{ \text{Lincoln Penny}, \text{Button} \right\}, \left\{ \text{Lincoln Dime}, \text{Button} \right\}, \left\{ \text{Lincoln Penny}, \text{Lincoln Dime}, \text{Button} \right\} \right\}$$

$$|S| < |\wp(S)|$$

$$S = \{a, b, c, d\}$$

$$\begin{aligned} \wp(S) = \{ & \\ & \emptyset, \\ & \{a\}, \{b\}, \{c\}, \{d\}, \\ & \{a, b\}, \{a, c\}, \{a, d\}, \{b, c\}, \{b, d\}, \{c, d\} \\ & \{a, b, c\}, \{a, b, d\}, \{a, c, d\}, \{b, c, d\}, \\ & \{a, b, c, d\} \\ & \} \end{aligned}$$

$$|S| < |\wp(S)|$$

If $|S|$ is infinite, what is the relation between $|S|$ and $|\wp(S)|$?

Does $|S| = |\wp(S)|$?

If $|S| = |\wp(S)|$, we can pair up the elements of S and the elements of $\wp(S)$ without leaving anything out.

If $|S| = |\wp(S)|$, we can pair up the elements of S and **the elements of $\wp(S)$** without leaving anything out.

If $|S| = |\wp(S)|$, we can pair up the elements of S and **the subsets of S** without leaving anything out.

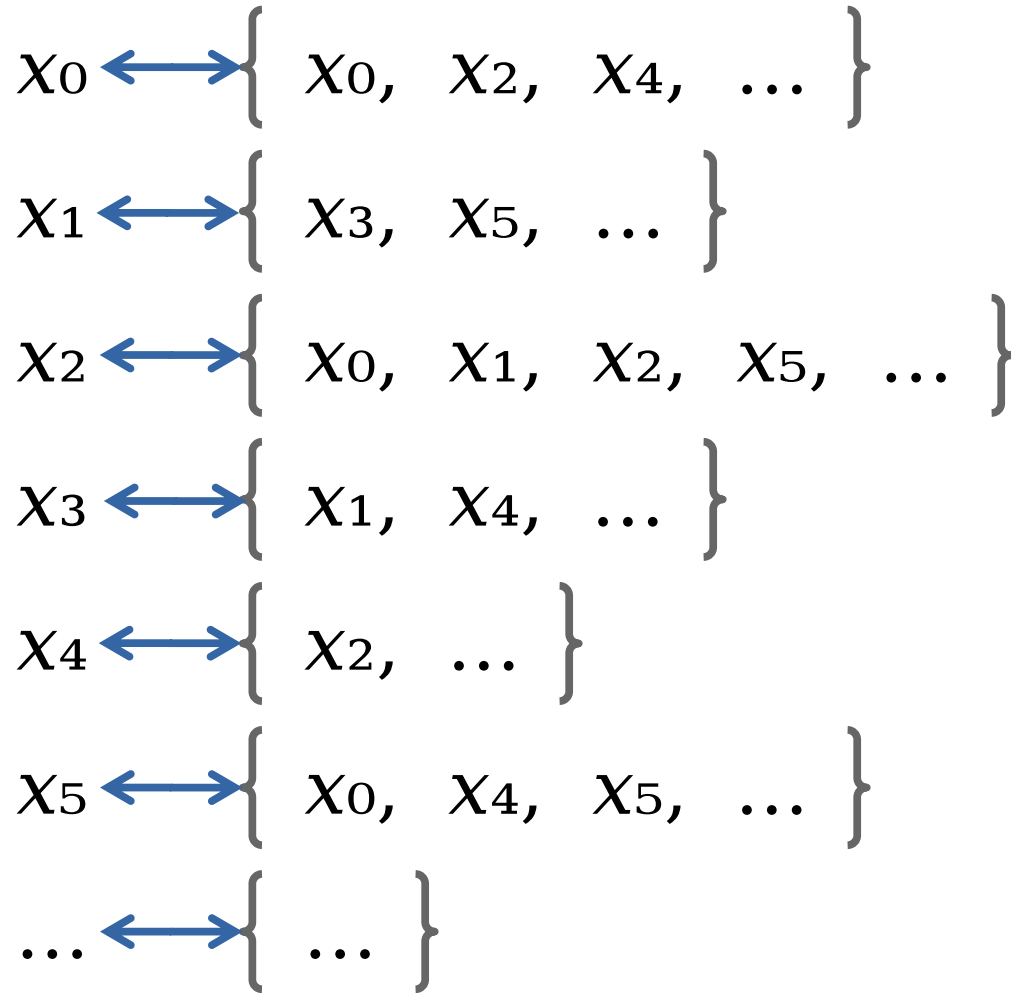
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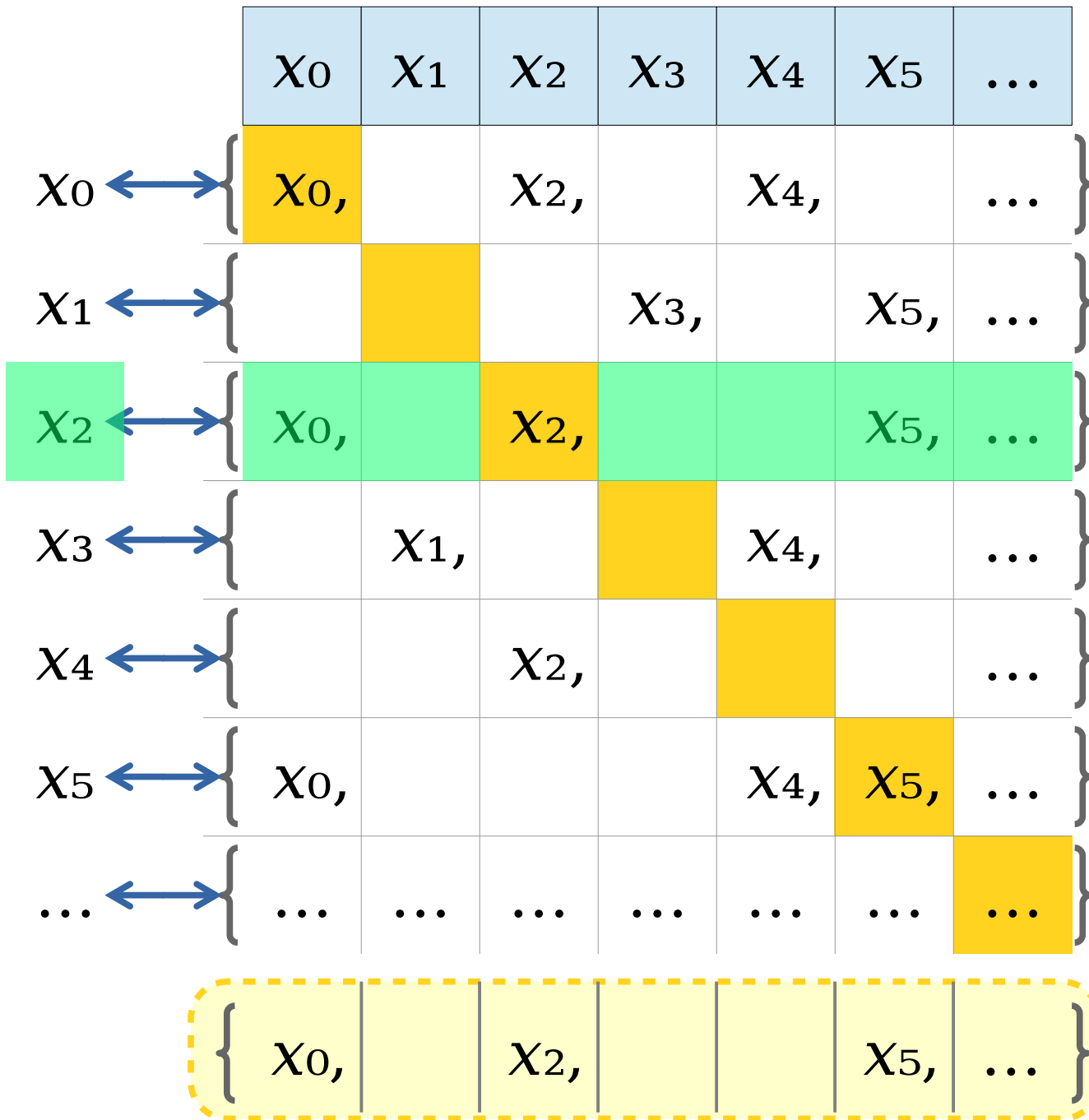
If $|S| = |\wp(S)|$, we can pair up the elements of S and the subsets of S without leaving anything out.

What would that look like?

$$\begin{array}{l}
x_0 \longleftrightarrow \{ x_0, x_2, x_4, \dots \} \\
x_1 \longleftrightarrow \{ x_3, x_5, \dots \} \\
x_2 \longleftrightarrow \{ x_0, x_1, x_2, x_5, \dots \} \\
x_3 \longleftrightarrow \{ x_1, x_4, \dots \} \\
x_4 \longleftrightarrow \{ x_2, \dots \} \\
x_5 \longleftrightarrow \{ x_0, x_4, x_5, \dots \} \\
\dots \longleftrightarrow \{ \dots \}
\end{array}$$

x_0	x_1	x_2	x_3	x_4	x_5	\dots
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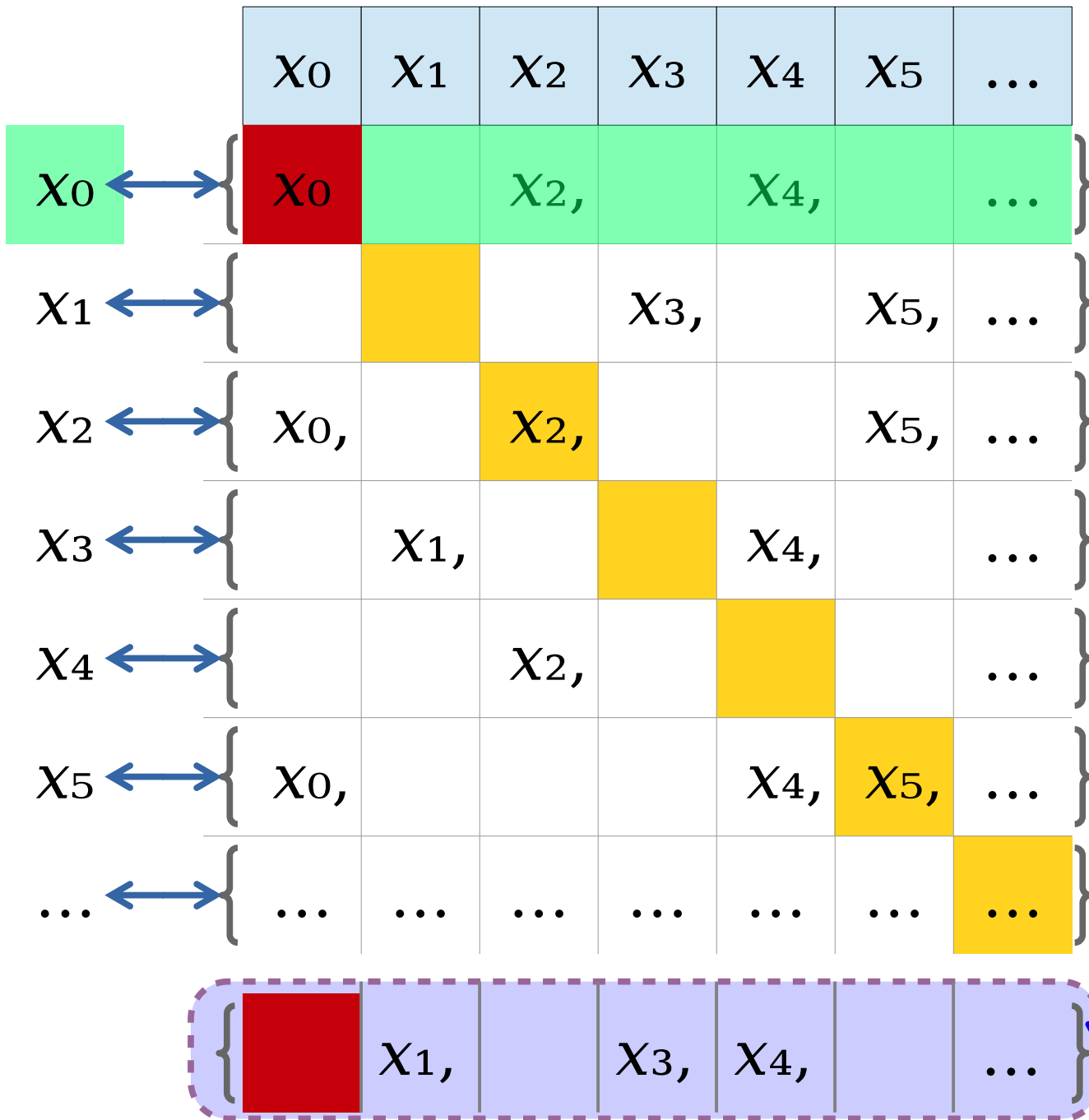


Which element is paired with this set?

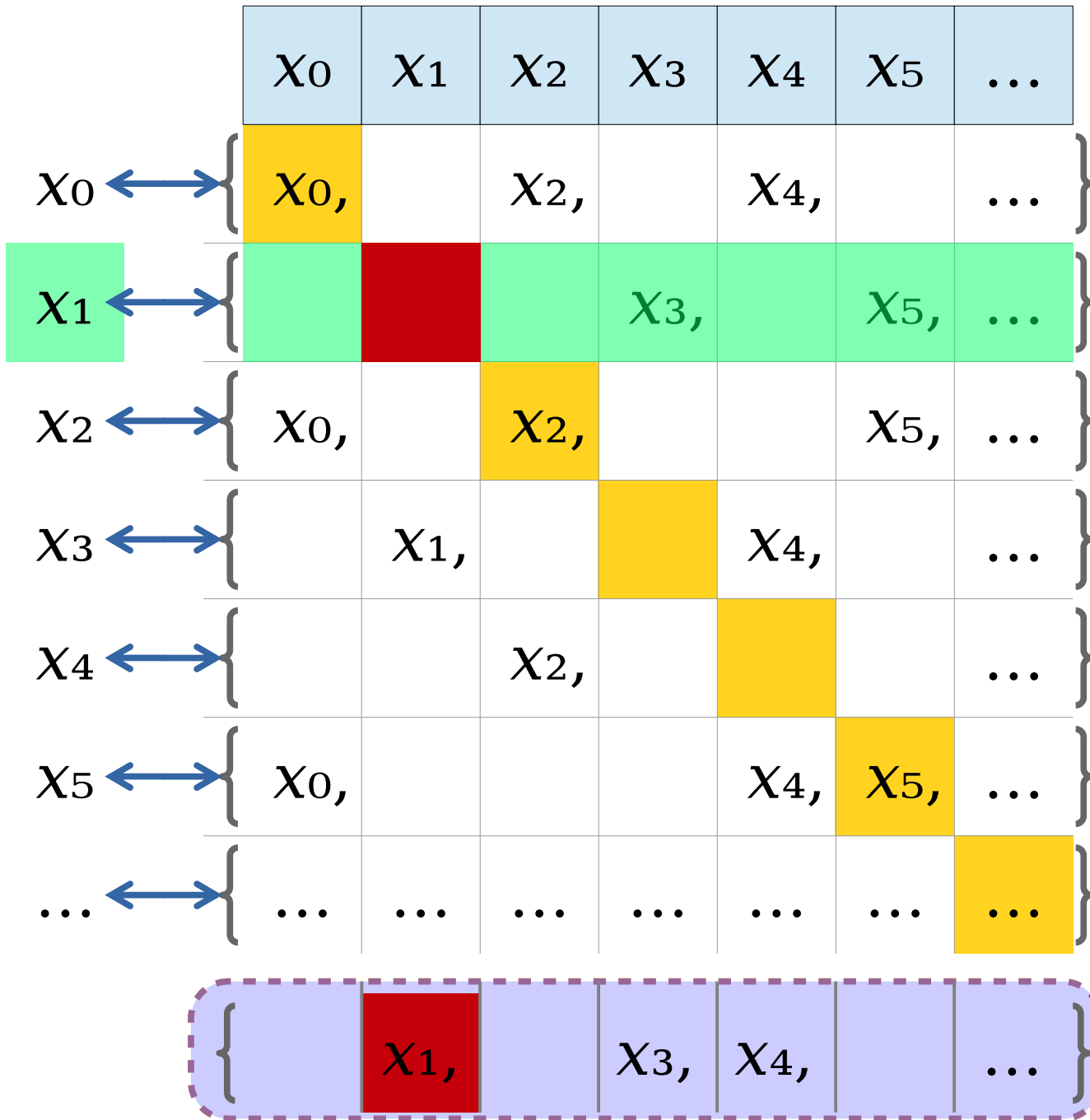
	x_0	x_1	x_2	x_3	x_4	x_5	...
x_0 ↔	$x_0,$		$x_2,$		$x_4,$...
x_1 ↔				$x_3,$		$x_5,$...
x_2 ↔	$x_0,$		$x_2,$			$x_5,$...
x_3 ↔		$x_1,$			$x_4,$...
x_4 ↔			$x_2,$...
x_5 ↔	$x_0,$				$x_4,$	$x_5,$...
...

“Flip” this set. Swap what’s included and what’s excluded.

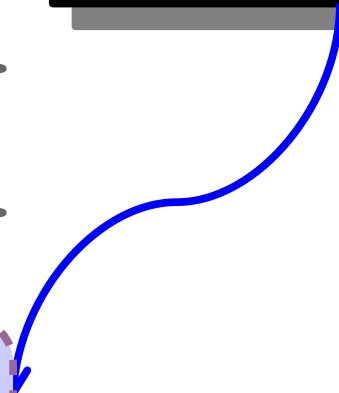
{ $x_1,$ $x_3,$ $x_4,$... }

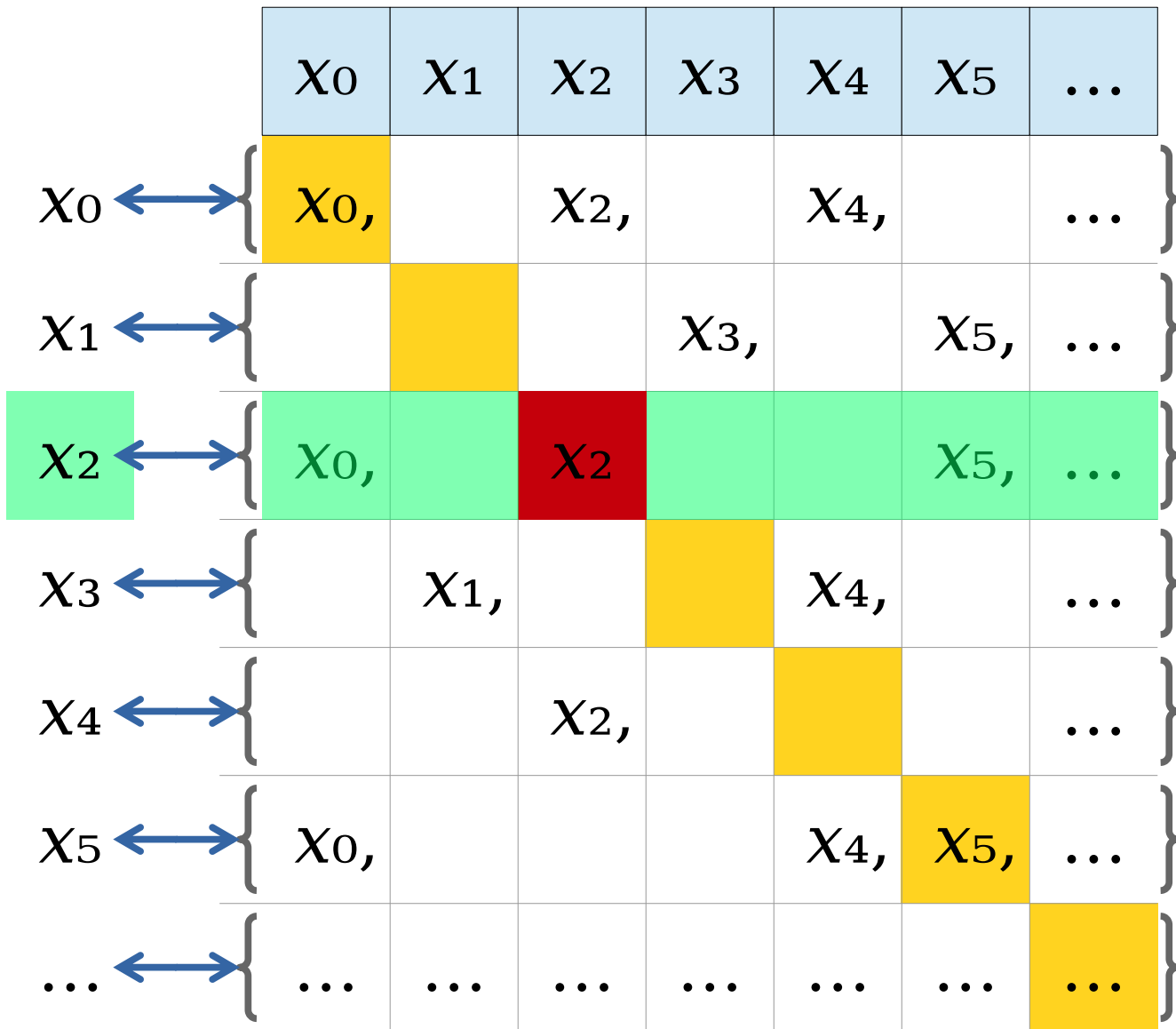


Which element is paired with this set?

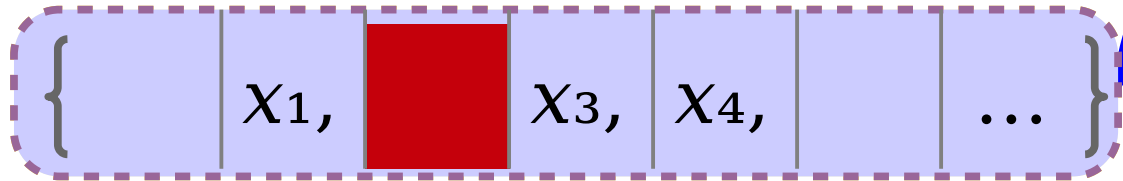


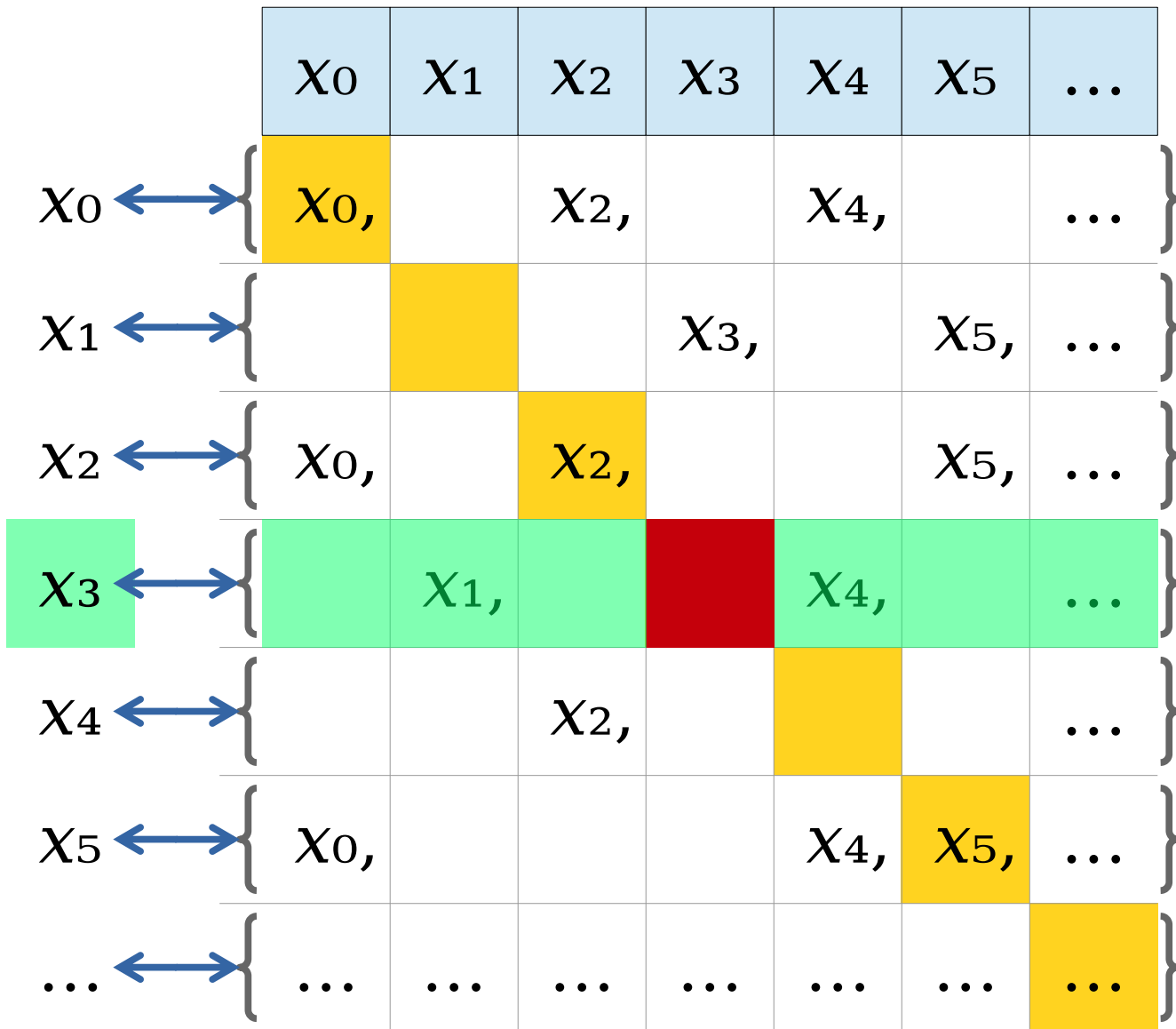
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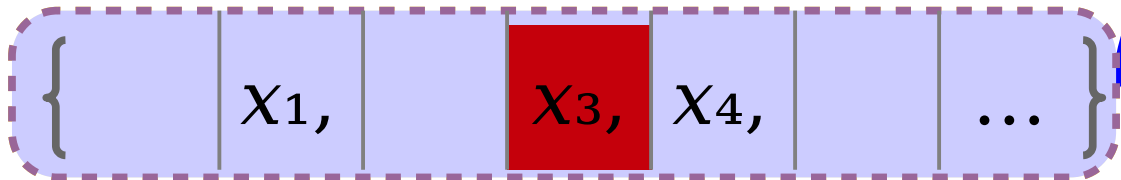


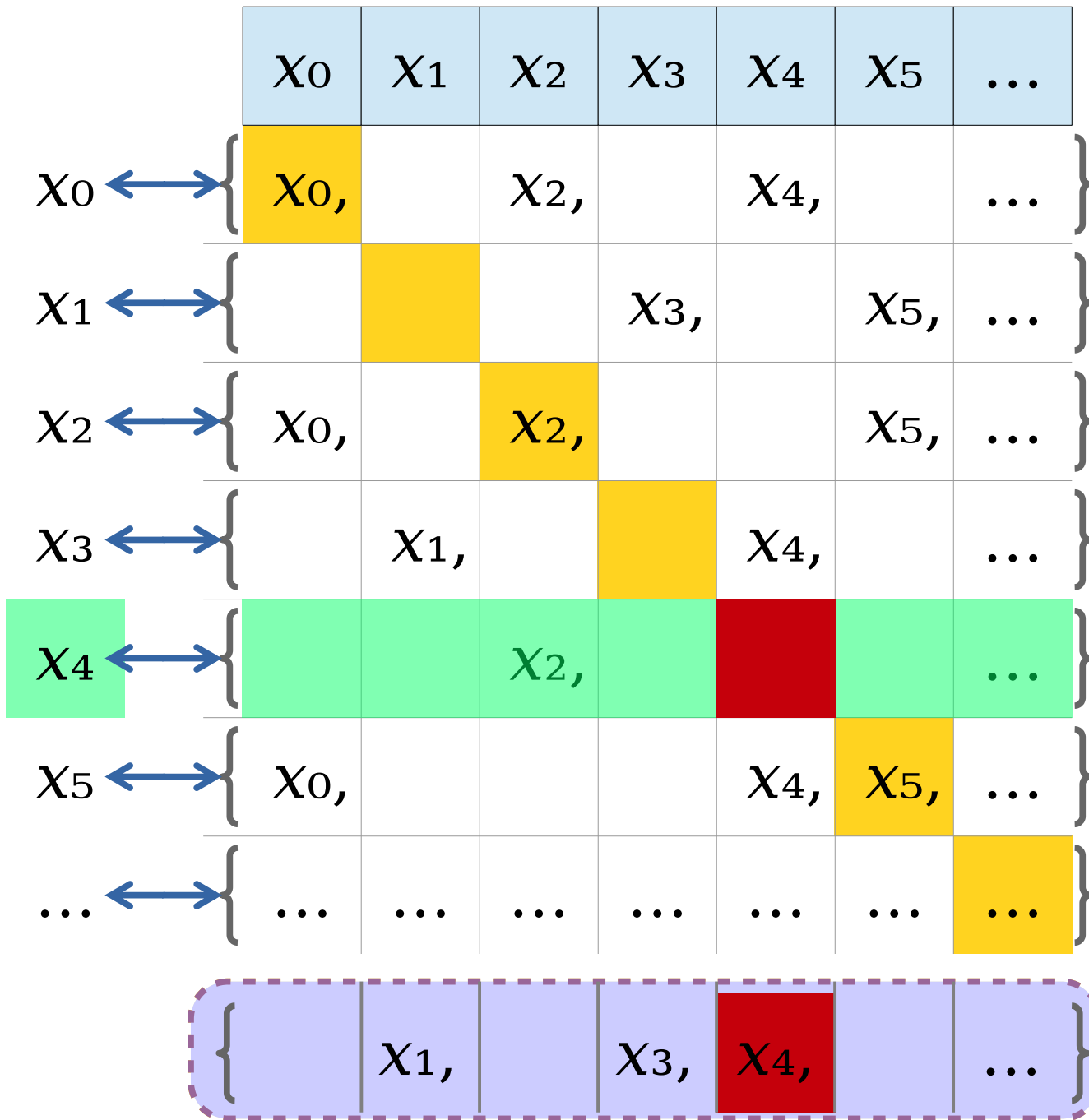
Which element is paired with this set?



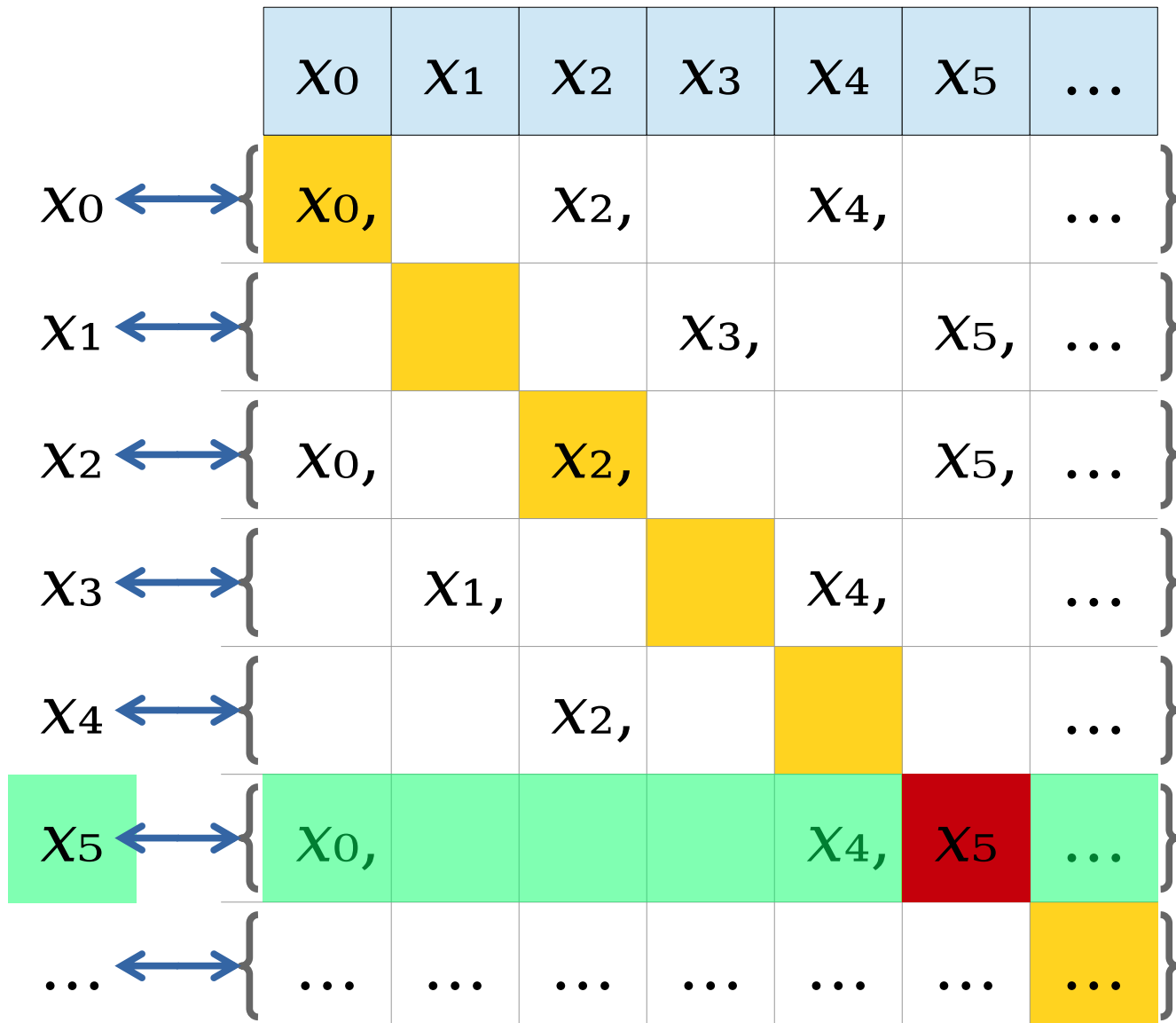


Which element is paired with this set?

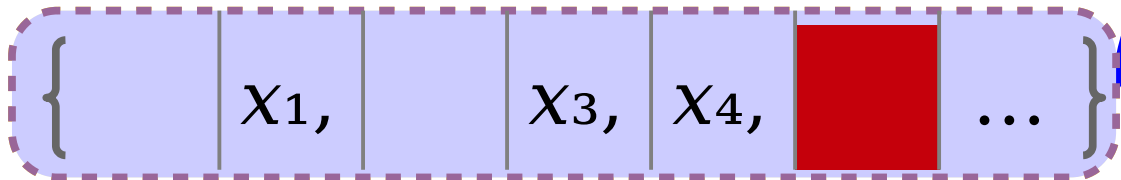


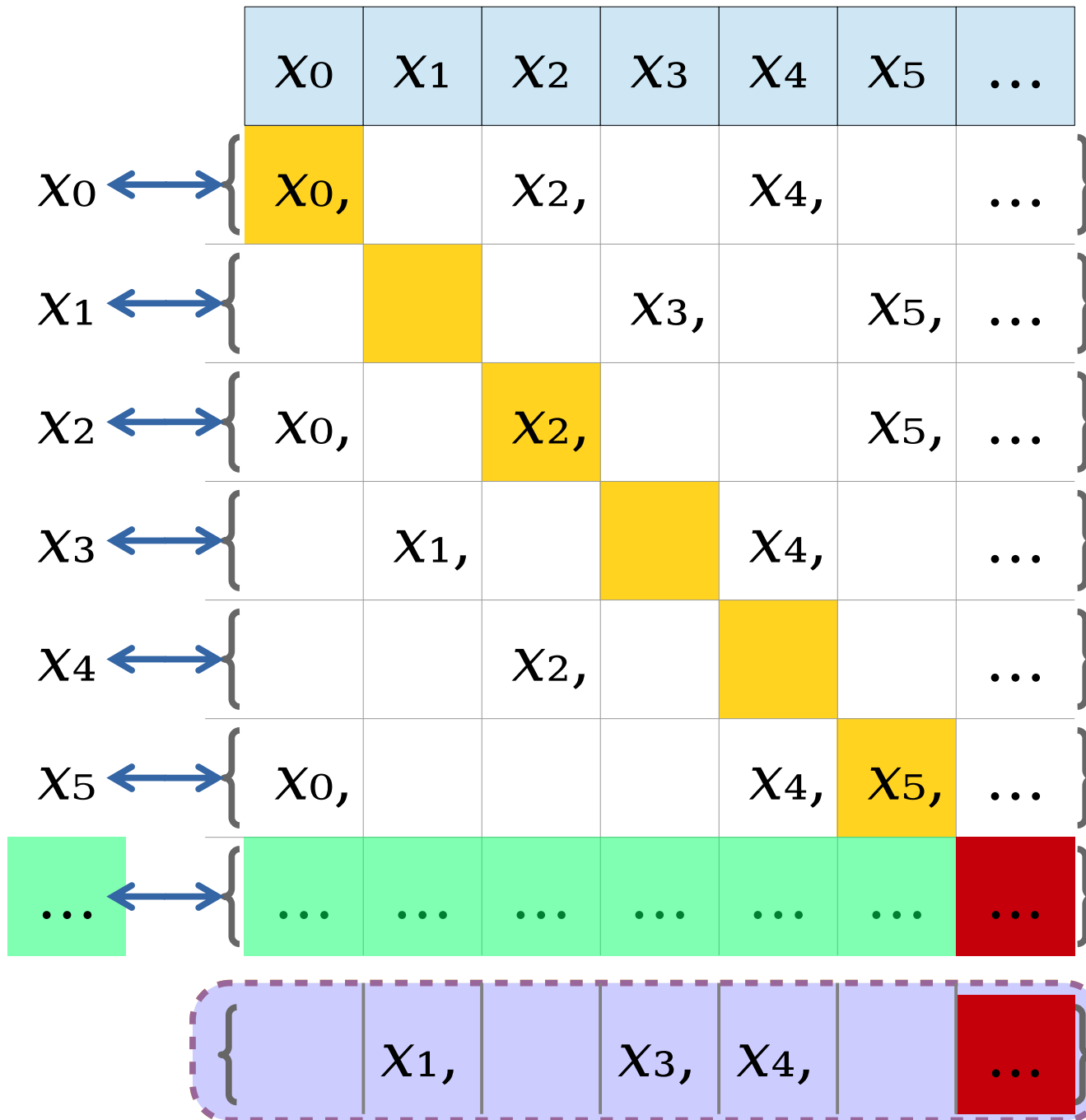


Which element is paired with this set?



Which element is paired with this set?





Which element is paired with this set?

The Diagonalization Proof

No matter how we pair up elements of S and subsets of S , the complemented diagonal won't appear in the table.

In row n , the n th element must be wrong.

No matter how we pair up elements of S and subsets of S , there is *always* at least one subset left over.

This result is ***Cantor's theorem***: Every set is strictly smaller than its power set:

If S is a set, then $|S| < |\wp(S)|$.

Infinite Cardinalities

By Cantor's Theorem:

$$|\mathbb{N}| < |\wp(\mathbb{N})|$$

$$|\wp(\mathbb{N})| < |\wp(\wp(\mathbb{N}))|$$

$$|\wp(\wp(\mathbb{N}))| < |\wp(\wp(\wp(\mathbb{N})))|$$

$$|\wp(\wp(\wp(\mathbb{N})))| < |\wp(\wp(\wp(\wp(\mathbb{N}))))|$$

...

Not all infinite sets have the same size!

There is no biggest infinity!

There are infinitely many infinities!

What does this have to do
with computation?

“The set of all computer programs”

“The set of all problems to solve”

Where We're Going

A ***string*** is a sequence of characters.

We're going to prove the following results:

- There are ***at most*** as many programs as there are strings.
- There are ***at least*** as many problems as there are sets of strings.

This leads to some *incredible* results – we'll see why in a minute!

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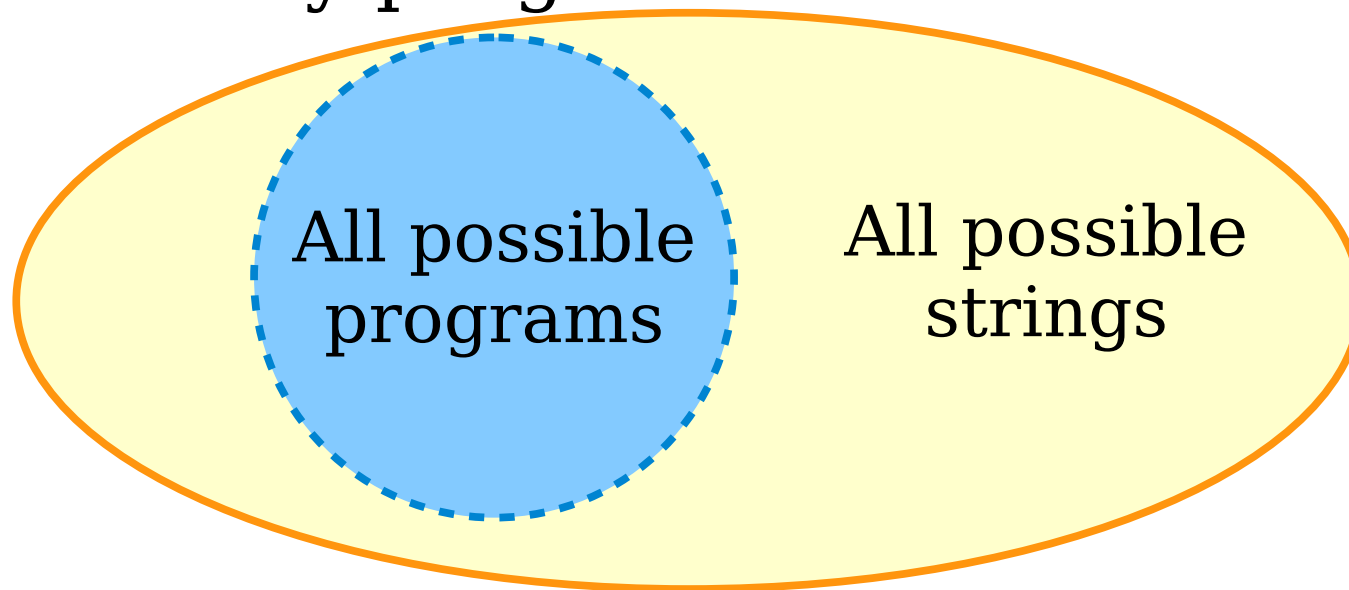
- There are ***at most*** as many programs as there are strings.
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Strings and Programs

The source code of a computer program is just a (long, structured, well-commented) string of text.

All programs are strings, but not all strings are necessarily programs.



$$|\mathbf{Programs}| \leq |\mathbf{Strings}|$$

Where We're Going

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A ***string*** is a sequence of characters.

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- There are *at least* as many problems as there are sets of strings.

This leads to some *incredible* results – we'll see why in a minute!

Strings and Problems

- There is a connection between the number of sets of strings and the number of problems to solve.
- Let S be any set of strings. This set S gives rise to a problem to solve:

Given a string w , determine whether $w \in S$.

Strings and Problems

Given a string w , determine whether $w \in S$.

Suppose that S is the set

$$S = \{ "a", "b", "c", \dots, "z" \}$$

From this set S , we get this problem:

Given a string w , determine whether w is a single lower-case English letter.

Strings and Problems

Given a string w , determine whether $w \in S$.

Suppose that S is the set

$$S = \{ "0", "1", "2", \dots, "9", "10", "11", \dots \}$$

From this set S , we get this problem:

Given a string w , determine whether w represents a natural number.

Strings and Problems

Given a string w , determine whether $w \in S$.

Suppose that S is the set

$$S = \{ p \mid p \text{ is a legal C++ program} \}$$

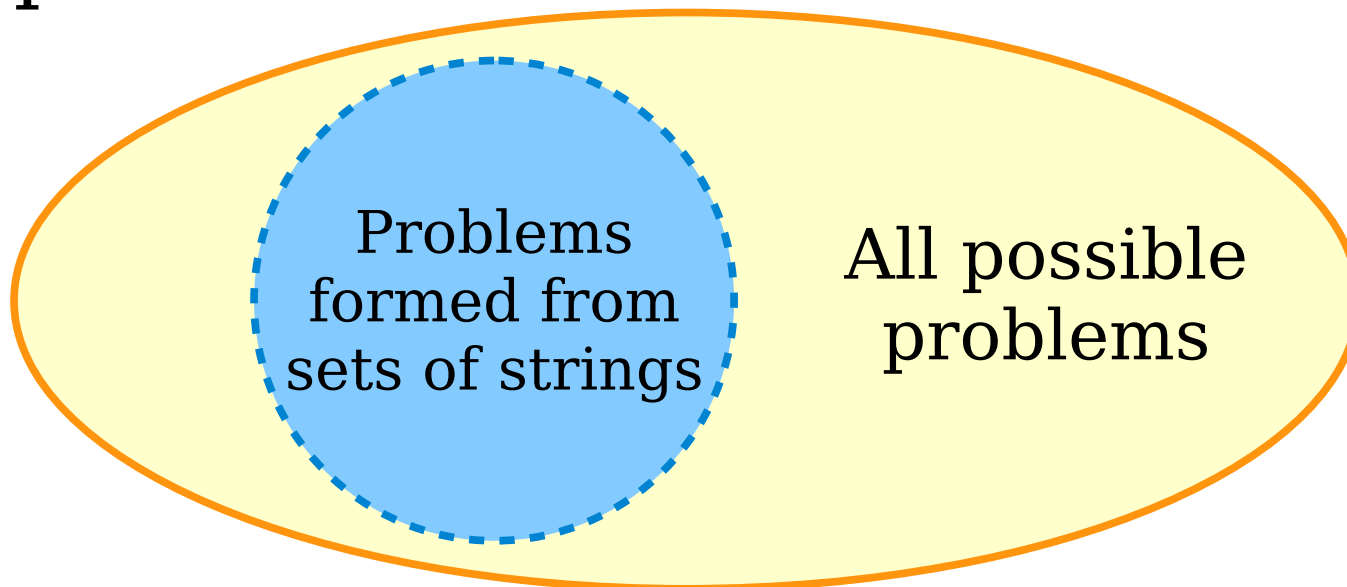
From this set S , we get this problem:

Given a string w , determine whether w is a legal C++ program.

Strings and Problems

Every set of strings gives rise to a unique problem to solve.

Other problems exist as well.



$$|\mathbf{Sets\ of\ Strings}| \leq |\mathbf{Problems}|$$

Where We're Going

A ***string*** is a sequence of characters.

We're going to prove the following results:

- There are ***at most*** as many programs as there are strings. ✓
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Where We're Going

A *string* is a sequence of characters.

We're going to prove the following results:

- There are *at most* as many programs as there are strings. ✓
- There are *at least* as many problems as there are sets of strings. ✓

This leads to some *incredible* results – we'll see why ~~in a minute!~~ *right now!*

Every computer program is a string.

So, the number of programs is at most the number of strings.

From Cantor's Theorem, we know that there are more sets of strings than strings.

There are at least as many problems as there are sets of strings.

$$|\mathbf{Programs}| \leq |\mathbf{Strings}| < |\wp(\mathbf{Strings})| \leq |\mathbf{Problems}|$$

Every computer program is a string.

So, the number of programs is at most the number of strings.

From Cantor's Theorem, we know that there are more sets of strings than strings.

There are at least as many problems as there are sets of strings.

|Programs| < |Problems|

There are more problems to solve than there are programs to solve them.

|Programs| < |Problems|

It Gets Worse

Using more advanced set theory, we can show that there are *infinitely more* problems than solutions.

In fact, if you pick a totally random problem, the probability that you can solve it is *zero*.

More troubling fact: We've just shown that *some* problems are impossible to solve with computers, but we don't know *which* problems those are!

We need to develop a more nuanced understanding of computation.

Where We're Going

What makes a problem impossible to solve with computers?

- Is there a deep reason why certain problems can't be solved with computers, or is it completely arbitrary?
- How do you know when you're looking at an impossible problem?
- Are these real-world problems, or are they highly contrived?

How do we know that we're right?

- How can we back up our pictures with rigorous proofs?
- How do we build a mathematical framework for studying computation?

Recap

Introduction to Set Theory

- A set S is an unordered collection of unique objects.
 - They have a cardinality, $|S|$, that can be finite or infinite.
 - The empty set, \emptyset , is the set of cardinality 0.
- We can use element of ($x \in S$) to describe membership in the set S .
- We can use the following operations to make new sets:
 - Union: $(A \cup B)$
 - Intersection: $(A \cap B)$
 - Difference: $(A \setminus B)$
 - Symmetric Difference: $(A \Delta B)$
 - Power Set: $\wp(S)$
- The cardinality of S is less than the cardinality of the power set of S .
 - $|S| < |\wp(S)|$
- To compare two sets, we can use the subset relation. ($A \subseteq B$)
- There are more problems than programs.

Next Time

Mathematical Proof

- What is a mathematical proof?
- How can we prove things with certainty?